

libbfd

The Binary File Descriptor Library

First Edition—BFD version < 3.0 % Since no product is stable before version 3.0 :-)

Original Document Created: April 1991

Steve Chamberlain
Cygnus Support

Free Software Foundation
sac@www.gnu.org
BFD, 1.5
T_EXinfo 2007-06-20.13

Copyright © 1991, 2001, 2003, 2006, 2008 Free Software Foundation, Inc.

Permission is granted to copy, distribute and/or modify this document under the terms of the GNU Free Documentation License, Version 1.3 or any later version published by the Free Software Foundation; with no Invariant Sections, with no Front-Cover Texts, and with no Back-Cover Texts. A copy of the license is included in the section entitled “GNU Free Documentation License”.

Table of Contents

1	Introduction	1
1.1	History	1
1.2	How To Use BFD	1
1.3	What BFD Version 2 Can Do	2
1.3.1	Information Loss	2
1.3.2	The BFD canonical object-file format	3
2	BFD Front End	5
2.1	typedef bfd	5
2.2	Error reporting	11
2.2.1	Type bfd_error_type	11
2.2.1.1	bfd_get_error	11
2.2.1.2	bfd_set_error	12
2.2.1.3	bfd_errmsg	12
2.2.1.4	bfd_perror	12
2.2.2	BFD error handler	12
2.2.2.1	bfd_set_error_handler	12
2.2.2.2	bfd_set_error_program_name	12
2.2.2.3	bfd_get_error_handler	13
2.3	Miscellaneous	13
2.3.1	Miscellaneous functions	13
2.3.1.1	bfd_get_reloc_upper_bound	13
2.3.1.2	bfd_canonicalize_reloc	13
2.3.1.3	bfd_set_reloc	13
2.3.1.4	bfd_set_file_flags	13
2.3.1.5	bfd_get_arch_size	14
2.3.1.6	bfd_get_sign_extend_vma	14
2.3.1.7	bfd_set_start_address	14
2.3.1.8	bfd_get_gp_size	14
2.3.1.9	bfd_set_gp_size	15
2.3.1.10	bfd_scan_vma	15
2.3.1.11	bfd_copy_private_header_data	15
2.3.1.12	bfd_copy_private_bfd_data	15
2.3.1.13	bfd_merge_private_bfd_data	16
2.3.1.14	bfd_set_private_flags	16
2.3.1.15	Other functions	16
2.3.1.16	bfd_alt_mach_code	18
2.3.1.17	bfd_preserve_save	18
2.3.1.18	bfd_preserve_restore	19
2.3.1.19	bfd_preserve_finish	19
2.3.1.20	bfd_emul_get_maxpagesize	19
2.3.1.21	bfd_emul_set_maxpagesize	19
2.3.1.22	bfd_emul_get_commonpagesize	19

2.3.1.23	bfd_emul_set_commonpagesize.....	20
2.3.1.24	bfd_demangle.....	20
2.3.1.25	struct bfd_iovec.....	20
2.3.1.26	bfd_get_mtime.....	21
2.3.1.27	bfd_get_size.....	21
2.3.1.28	bfd_mmap.....	21
2.4	Memory Usage.....	21
2.5	Initialization.....	22
2.5.1	Initialization functions.....	22
2.5.1.1	bfd_init.....	22
2.6	Sections.....	22
2.6.1	Section input.....	22
2.6.2	Section output.....	22
2.6.3	Link orders.....	23
2.6.4	typedef asection.....	23
2.6.5	Section prototypes.....	35
2.6.5.1	bfd_section_list_clear.....	35
2.6.5.2	bfd_get_section_by_name.....	35
2.6.5.3	bfd_get_section_by_name_if.....	35
2.6.5.4	bfd_get_unique_section_name.....	36
2.6.5.5	bfd_make_section_old_way.....	36
2.6.5.6	bfd_make_section_anyway_with_flags.....	36
2.6.5.7	bfd_make_section_anyway.....	37
2.6.5.8	bfd_make_section_with_flags.....	37
2.6.5.9	bfd_make_section.....	37
2.6.5.10	bfd_set_section_flags.....	37
2.6.5.11	bfd_rename_section.....	37
2.6.5.12	bfd_map_over_sections.....	38
2.6.5.13	bfd_sections_find_if.....	38
2.6.5.14	bfd_set_section_size.....	38
2.6.5.15	bfd_set_section_contents.....	38
2.6.5.16	bfd_get_section_contents.....	39
2.6.5.17	bfd_malloc_and_get_section.....	39
2.6.5.18	bfd_copy_private_section_data.....	39
2.6.5.19	bfd_generic_is_group_section.....	39
2.6.5.20	bfd_generic_discard_group.....	40
2.7	Symbols.....	40
2.7.1	Reading symbols.....	40
2.7.2	Writing symbols.....	41
2.7.3	Mini Symbols.....	42
2.7.4	typedef asymbol.....	42
2.7.5	Symbol handling functions.....	45
2.7.5.1	bfd_get_symtab_upper_bound.....	45
2.7.5.2	bfd_is_local_label.....	45
2.7.5.3	bfd_is_local_label_name.....	45
2.7.5.4	bfd_is_target_special_symbol.....	46
2.7.5.5	bfd_canonicalize_symtab.....	46
2.7.5.6	bfd_set_symtab.....	46

2.7.5.7	<code>bfd_print_symbol_vandf</code>	46
2.7.5.8	<code>bfd_make_empty_symbol</code>	46
2.7.5.9	<code>_bfd_generic_make_empty_symbol</code>	47
2.7.5.10	<code>bfd_make_debug_symbol</code>	47
2.7.5.11	<code>bfd_decode_symclass</code>	47
2.7.5.12	<code>bfd_is_undefined_symclass</code>	47
2.7.5.13	<code>bfd_symbol_info</code>	47
2.7.5.14	<code>bfd_copy_private_symbol_data</code>	48
2.8	Archives	48
2.8.1	Archive functions	48
2.8.1.1	<code>bfd_get_next_mapent</code>	49
2.8.1.2	<code>bfd_set_archive_head</code>	49
2.8.1.3	<code>bfd_openr_next_archived_file</code>	49
2.9	File formats	49
2.9.1	File format functions	49
2.9.1.1	<code>bfd_check_format</code>	50
2.9.1.2	<code>bfd_check_format_matches</code>	50
2.9.1.3	<code>bfd_set_format</code>	50
2.9.1.4	<code>bfd_format_string</code>	50
2.10	Relocations	51
2.10.1	<code>typedef arelent</code>	51
2.10.1.1	<code>enum complain_overflow</code>	54
2.10.1.2	<code>reloc_howto_type</code>	54
2.10.1.3	The HOWTO Macro	56
2.10.1.4	<code>bfd_get_reloc_size</code>	57
2.10.1.5	<code>arelent_chain</code>	57
2.10.1.6	<code>bfd_check_overflow</code>	58
2.10.1.7	<code>bfd_perform_relocation</code>	58
2.10.1.8	<code>bfd_install_relocation</code>	58
2.10.2	The howto manager	59
2.10.2.1	<code>bfd_reloc_code_type</code>	59
2.10.2.2	<code>bfd_reloc_type_lookup</code>	105
2.10.2.3	<code>bfd_default_reloc_type_lookup</code>	105
2.10.2.4	<code>bfd_get_reloc_code_name</code>	105
2.10.2.5	<code>bfd_generic_relax_section</code>	105
2.10.2.6	<code>bfd_generic_gc_sections</code>	105
2.10.2.7	<code>bfd_generic_merge_sections</code>	106
2.10.2.8	<code>bfd_generic_get_relocated_section_contents</code>	106
2.11	Core files	106
2.11.1	Core file functions	106
2.11.1.1	<code>bfd_core_file_failing_command</code>	106
2.11.1.2	<code>bfd_core_file_failing_signal</code>	106
2.11.1.3	<code>bfd_core_file_pid</code>	107
2.11.1.4	<code>core_file_matches_executable_p</code>	107
2.11.1.5	<code>generic_core_file_matches_executable_p</code>	107
2.12	Targets	107
2.12.1	<code>bfd_target</code>	108

2.12.1.1	bfd_set_default_target.....	117
2.12.1.2	bfd_find_target.....	117
2.12.1.3	bfd_get_target_info.....	118
2.12.1.4	bfd_target_list.....	118
2.12.1.5	bfd_seach_for_target.....	118
2.13	Architectures.....	118
2.13.1	bfd_architecture.....	119
2.13.2	bfd_arch_info.....	127
2.13.2.1	bfd_printable_name.....	127
2.13.2.2	bfd_scan_arch.....	128
2.13.2.3	bfd_arch_list.....	128
2.13.2.4	bfd_arch_get_compatible.....	128
2.13.2.5	bfd_default_arch_struct.....	128
2.13.2.6	bfd_set_arch_info.....	128
2.13.2.7	bfd_default_set_arch_mach.....	128
2.13.2.8	bfd_get_arch.....	129
2.13.2.9	bfd_get_mach.....	129
2.13.2.10	bfd_arch_bits_per_byte.....	129
2.13.2.11	bfd_arch_bits_per_address.....	129
2.13.2.12	bfd_default_compatible.....	129
2.13.2.13	bfd_default_scan.....	129
2.13.2.14	bfd_get_arch_info.....	129
2.13.2.15	bfd_lookup_arch.....	130
2.13.2.16	bfd_printable_arch_mach.....	130
2.13.2.17	bfd_octets_per_byte.....	130
2.13.2.18	bfd_arch_mach_octets_per_byte.....	130
2.14	Opening and closing BFDs.....	130
2.14.1	Functions for opening and closing.....	130
2.14.1.1	bfd_fopen.....	130
2.14.1.2	bfd_openr.....	131
2.14.1.3	bfd_fdopenr.....	131
2.14.1.4	bfd_openstreamr.....	131
2.14.1.5	bfd_openr_iovec.....	131
2.14.1.6	bfd_openw.....	132
2.14.1.7	bfd_close.....	132
2.14.1.8	bfd_close_all_done.....	133
2.14.1.9	bfd_create.....	133
2.14.1.10	bfd_make_writable.....	133
2.14.1.11	bfd_make_readable.....	133
2.14.1.12	bfd_alloc.....	134
2.14.1.13	bfd_alloc2.....	134
2.14.1.14	bfd_zalloc.....	134
2.14.1.15	bfd_zalloc2.....	134
2.14.1.16	bfd_calc_gnu_debuglink_crc32.....	134
2.14.1.17	get_debug_link_info.....	135
2.14.1.18	separate_debug_file_exists.....	135
2.14.1.19	find_separate_debug_file.....	135
2.14.1.20	bfd_follow_gnu_debuglink.....	135

2.14.1.21	<code>bfd_create_gnu_debuglink_section</code>	135
2.14.1.22	<code>bfd_fill_in_gnu_debuglink_section</code>	136
2.15	Implementation details.....	136
2.15.1	Internal functions.....	136
2.15.1.1	<code>bfd_write_bigendian_4byte_int</code>	136
2.15.1.2	<code>bfd_put_size</code>	136
2.15.1.3	<code>bfd_get_size</code>	136
2.15.1.4	<code>bfd_h_put_size</code>	138
2.15.1.5	<code>bfd_log2</code>	139
2.16	File caching.....	139
2.16.1	Caching functions.....	139
2.16.1.1	<code>bfd_cache_init</code>	139
2.16.1.2	<code>bfd_cache_close</code>	140
2.16.1.3	<code>bfd_cache_close_all</code>	140
2.16.1.4	<code>bfd_open_file</code>	140
2.17	Linker Functions.....	140
2.17.1	Creating a linker hash table.....	141
2.17.2	Adding symbols to the hash table.....	141
2.17.2.1	Differing file formats.....	141
2.17.2.2	Adding symbols from an object file.....	142
2.17.2.3	Adding symbols from an archive.....	142
2.17.3	Performing the final link.....	143
2.17.3.1	Information provided by the linker.....	143
2.17.3.2	Relocating the section contents.....	144
2.17.3.3	Writing the symbol table.....	144
2.17.3.4	<code>bfd_link_split_section</code>	145
2.17.3.5	<code>bfd_section_already_linked</code>	145
2.17.3.6	<code>bfd_generic_define_common_symbol</code>	145
2.17.3.7	<code>bfd_find_version_for_sym</code>	146
2.18	Hash Tables.....	146
2.18.1	Creating and freeing a hash table.....	146
2.18.2	Looking up or entering a string.....	146
2.18.3	Traversing a hash table.....	147
2.18.4	Deriving a new hash table type.....	147
2.18.4.1	Define the derived structures.....	147
2.18.4.2	Write the derived creation routine.....	148
2.18.4.3	Write other derived routines.....	149
3	BFD back ends.....	150
3.1	What to Put Where.....	150
3.2	a.out backends.....	150
3.2.1	Relocations.....	151
3.2.2	Internal entry points.....	151
3.2.2.1	<code>aout_size_swap_exec_header_in</code>	151
3.2.2.2	<code>aout_size_swap_exec_header_out</code>	151
3.2.2.3	<code>aout_size_some_aout_object_p</code>	152
3.2.2.4	<code>aout_size_mkobject</code>	152
3.2.2.5	<code>aout_size_machine_type</code>	152

3.2.2.6	<code>aout_size_set_arch_mach</code>	152
3.2.2.7	<code>aout_size_new_section_hook</code>	153
3.3	coff backends	153
3.3.1	Porting to a new version of coff	153
3.3.2	How the coff backend works	153
3.3.2.1	File layout	153
3.3.2.2	Coff long section names	154
3.3.2.3	Bit twiddling	155
3.3.2.4	Symbol reading	155
3.3.2.5	Symbol writing	156
3.3.2.6	<code>coff_symbol_type</code>	156
3.3.2.7	<code>bfd_coff_backend_data</code>	158
3.3.2.8	Writing relocations	164
3.3.2.9	Reading linenumbers	164
3.3.2.10	Reading relocations	164
3.4	ELF backends	165
3.5	mmo backend	165
3.5.1	File layout	165
3.5.2	Symbol table format	167
3.5.3	mmo section mapping	169
BFD Index		179

1 Introduction

BFD is a package which allows applications to use the same routines to operate on object files whatever the object file format. A new object file format can be supported simply by creating a new BFD back end and adding it to the library.

BFD is split into two parts: the front end, and the back ends (one for each object file format).

- The front end of BFD provides the interface to the user. It manages memory and various canonical data structures. The front end also decides which back end to use and when to call back end routines.
- The back ends provide BFD its view of the real world. Each back end provides a set of calls which the BFD front end can use to maintain its canonical form. The back ends also may keep around information for their own use, for greater efficiency.

1.1 History

One spur behind BFD was the desire, on the part of the GNU 960 team at Intel Oregon, for interoperability of applications on their COFF and b.out file formats. Cygnus was providing GNU support for the team, and was contracted to provide the required functionality.

The name came from a conversation David Wallace was having with Richard Stallman about the library: RMS said that it would be quite hard—David said “BFD”. Stallman was right, but the name stuck.

At the same time, Ready Systems wanted much the same thing, but for different object file formats: IEEE-695, Oasys, Srecords, a.out and 68k coff.

BFD was first implemented by members of Cygnus Support; Steve Chamberlain (sac@cygnus.com), John Gilmore (gnu@cygnus.com), K. Richard Pixley (rich@cygnus.com) and David Henkel-Wallace (gumby@cygnus.com).

1.2 How To Use BFD

To use the library, include ‘bfd.h’ and link with ‘libbfd.a’.

BFD provides a common interface to the parts of an object file for a calling application.

When an application successfully opens a target file (object, archive, or whatever), a pointer to an internal structure is returned. This pointer points to a structure called `bfd`, described in ‘bfd.h’. Our convention is to call this pointer a BFD, and instances of it within code `abfd`. All operations on the target object file are applied as methods to the BFD. The mapping is defined within `bfd.h` in a set of macros, all beginning with ‘bfd_’ to reduce namespace pollution.

For example, this sequence does what you would probably expect: return the number of sections in an object file attached to a BFD `abfd`.

```
#include "bfd.h"

unsigned int number_of_sections (abfd)
bfd *abfd;
{
```

```
    return bfd_count_sections (abfd);  
}
```

The abstraction used within BFD is that an object file has:

- a header,
- a number of sections containing raw data (see [Section 2.6 \[Sections\]](#), page 22),
- a set of relocations (see [Section 2.10 \[Relocations\]](#), page 51), and
- some symbol information (see [Section 2.7 \[Symbols\]](#), page 40).

Also, BFDs opened for archives have the additional attribute of an index and contain subordinate BFDs. This approach is fine for a.out and coff, but loses efficiency when applied to formats such as S-records and IEEE-695.

1.3 What BFD Version 2 Can Do

When an object file is opened, BFD subroutines automatically determine the format of the input object file. They then build a descriptor in memory with pointers to routines that will be used to access elements of the object file's data structures.

As different information from the object files is required, BFD reads from different sections of the file and processes them. For example, a very common operation for the linker is processing symbol tables. Each BFD back end provides a routine for converting between the object file's representation of symbols and an internal canonical format. When the linker asks for the symbol table of an object file, it calls through a memory pointer to the routine from the relevant BFD back end which reads and converts the table into a canonical form. The linker then operates upon the canonical form. When the link is finished and the linker writes the output file's symbol table, another BFD back end routine is called to take the newly created symbol table and convert it into the chosen output format.

1.3.1 Information Loss

Information can be lost during output. The output formats supported by BFD do not provide identical facilities, and information which can be described in one form has nowhere to go in another format. One example of this is alignment information in b.out. There is nowhere in an a.out format file to store alignment information on the contained data, so when a file is linked from b.out and an a.out image is produced, alignment information will not propagate to the output file. (The linker will still use the alignment information internally, so the link is performed correctly).

Another example is COFF section names. COFF files may contain an unlimited number of sections, each one with a textual section name. If the target of the link is a format which does not have many sections (e.g., a.out) or has sections without names (e.g., the Oasys format), the link cannot be done simply. You can circumvent this problem by describing the desired input-to-output section mapping with the linker command language.

Information can be lost during canonicalization. The BFD internal canonical form of the external formats is not exhaustive; there are structures in input formats for which there is no direct representation internally. This means that the BFD back ends cannot maintain all possible data richness through the transformation between external to internal and back to external formats.

This limitation is only a problem when an application reads one format and writes another. Each BFD back end is responsible for maintaining as much data as possible, and the internal BFD canonical form has structures which are opaque to the BFD core, and exported only to the back ends. When a file is read in one format, the canonical form is generated for BFD and the application. At the same time, the back end saves away any information which may otherwise be lost. If the data is then written back in the same format, the back end routine will be able to use the canonical form provided by the BFD core as well as the information it prepared earlier. Since there is a great deal of commonality between back ends, there is no information lost when linking or copying big endian COFF to little endian COFF, or `a.out` to `b.out`. When a mixture of formats is linked, the information is only lost from the files whose format differs from the destination.

1.3.2 The BFD canonical object-file format

The greatest potential for loss of information occurs when there is the least overlap between the information provided by the source format, that stored by the canonical format, and that needed by the destination format. A brief description of the canonical form may help you understand which kinds of data you can count on preserving across conversions.

files Information stored on a per-file basis includes target machine architecture, particular implementation format type, a demand pageable bit, and a write protected bit. Information like Unix magic numbers is not stored here—only the magic numbers’ meaning, so a ZMAGIC file would have both the demand pageable bit and the write protected text bit set. The byte order of the target is stored on a per-file basis, so that big- and little-endian object files may be used with one another.

sections Each section in the input file contains the name of the section, the section’s original address in the object file, size and alignment information, various flags, and pointers into other BFD data structures.

symbols Each symbol contains a pointer to the information for the object file which originally defined it, its name, its value, and various flag bits. When a BFD back end reads in a symbol table, it relocates all symbols to make them relative to the base of the section where they were defined. Doing this ensures that each symbol points to its containing section. Each symbol also has a varying amount of hidden private data for the BFD back end. Since the symbol points to the original file, the private data format for that symbol is accessible. `ld` can operate on a collection of symbols of wildly different formats without problems. Normal global and simple local symbols are maintained on output, so an output file (no matter its format) will retain symbols pointing to functions and to global, static, and common variables. Some symbol information is not worth retaining; in `a.out`, type information is stored in the symbol table as long symbol names. This information would be useless to most COFF debuggers; the linker has command line switches to allow users to throw it away.

There is one word of type information within the symbol, so if the format supports symbol type information within symbols (for example, COFF, IEEE, Oasys) and the type is simple enough to fit within one word (nearly everything but aggregates), the information will be preserved.

relocation level

Each canonical BFD relocation record contains a pointer to the symbol to relocate to, the offset of the data to relocate, the section the data is in, and a pointer to a relocation type descriptor. Relocation is performed by passing messages through the relocation type descriptor and the symbol pointer. Therefore, relocations can be performed on output data using a relocation method that is only available in one of the input formats. For instance, Oasys provides a byte relocation format. A relocation record requesting this relocation type would point indirectly to a routine to perform this, so the relocation may be performed on a byte being written to a 68k COFF file, even though 68k COFF has no such relocation type.

line numbers

Object formats can contain, for debugging purposes, some form of mapping between symbols, source line numbers, and addresses in the output file. These addresses have to be relocated along with the symbol information. Each symbol with an associated list of line number records points to the first record of the list. The head of a line number list consists of a pointer to the symbol, which allows finding out the address of the function whose line number is being described. The rest of the list is made up of pairs: offsets into the section and line numbers. Any format which can simply derive this information can pass it successfully between formats (COFF, IEEE and Oasys).

2 BFD Front End

2.1 typedef bfd

A BFD has type `bfd`; objects of this type are the cornerstone of any application using BFD. Using BFD consists of making references though the BFD and to data in the BFD.

Here is the structure that defines the type `bfd`. It contains the major data about the file and pointers to the rest of the data.

```
enum bfd_direction
{
    no_direction = 0,
    read_direction = 1,
    write_direction = 2,
    both_direction = 3
};

struct bfd
{
    /* A unique identifier of the BFD */
    unsigned int id;

    /* The filename the application opened the BFD with. */
    const char *filename;

    /* A pointer to the target jump table. */
    const struct bfd_target *xvec;

    /* The IOSTREAM, and corresponding IO vector that provide access
       to the file backing the BFD. */
    void *iostream;
    const struct bfd_iovec *iovec;

    /* The caching routines use these to maintain a
       least-recently-used list of BFDs. */
    struct bfd *lru_prev, *lru_next;

    /* When a file is closed by the caching routines, BFD retains
       state information on the file here... */
    ufile_ptr where;

    /* File modified time, if mtime_set is TRUE. */
    long mtime;

    /* Reserved for an unimplemented file locking extension. */
    int ifd;
```

```
/* The format which belongs to the BFD. (object, core, etc.) */
bfd_format format;

/* The direction with which the BFD was opened. */
enum bfd_direction direction;

/* Format_specific flags. */
flagword flags;

/* Values that may appear in the flags field of a BFD. These also
   appear in the object_flags field of the bfd_target structure, where
   they indicate the set of flags used by that backend (not all flags
   are meaningful for all object file formats) (FIXME: at the moment,
   the object_flags values have mostly just been copied from backend
   to another, and are not necessarily correct). */

#define BFD_NO_FLAGS    0x00

/* BFD contains relocation entries. */
#define HAS_RELOC       0x01

/* BFD is directly executable. */
#define EXEC_P          0x02

/* BFD has line number information (basically used for F_LNNO in a
   COFF header). */
#define HAS_LINENO      0x04

/* BFD has debugging information. */
#define HAS_DEBUG       0x08

/* BFD has symbols. */
#define HAS_SYMS        0x10

/* BFD has local symbols (basically used for F_LSYMS in a COFF
   header). */
#define HAS_LOCALS      0x20

/* BFD is a dynamic object. */
#define DYNAMIC          0x40

/* Text section is write protected (if D_PAGED is not set, this is
   like an a.out NMAGIC file) (the linker sets this by default, but
   clears it for -r or -N). */
#define WP_TEXT         0x80
```

```
/* BFD is dynamically paged (this is like an a.out ZMAGIC file) (the
   linker sets this by default, but clears it for -r or -n or -N). */
#define D_PAGED          0x100

/* BFD is relaxable (this means that bfd_relax_section may be able to
   do something) (sometimes bfd_relax_section can do something even if
   this is not set). */
#define BFD_IS_RELAXABLE 0x200

/* This may be set before writing out a BFD to request using a
   traditional format. For example, this is used to request that when
   writing out an a.out object the symbols not be hashed to eliminate
   duplicates. */
#define BFD_TRADITIONAL_FORMAT 0x400

/* This flag indicates that the BFD contents are actually cached
   in memory. If this is set, iostream points to a bfd_in_memory
   struct. */
#define BFD_IN_MEMORY 0x800

/* The sections in this BFD specify a memory page. */
#define HAS_LOAD_PAGE 0x1000

/* This BFD has been created by the linker and doesn't correspond
   to any input file. */
#define BFD_LINKER_CREATED 0x2000

/* This may be set before writing out a BFD to request that it
   be written using values for UIDs, GIDs, timestamps, etc. that
   will be consistent from run to run. */
#define BFD_DETERMINISTIC_OUTPUT 0x4000

/* Compress sections in this BFD. */
#define BFD_COMPRESS 0x8000

/* Decompress sections in this BFD. */
#define BFD_DECOMPRESS 0x10000

/* Flags bits to be saved in bfd_preserve_save. */
#define BFD_FLAGS_SAVED \
  (BFD_IN_MEMORY | BFD_COMPRESS | BFD_DECOMPRESS)

/* Flags bits which are for BFD use only. */
#define BFD_FLAGS_FOR_BFD_USE_MASK \
  (BFD_IN_MEMORY | BFD_COMPRESS | BFD_DECOMPRESS | BFD_LINKER_CREATED \
   | BFD_TRADITIONAL_FORMAT | BFD_DETERMINISTIC_OUTPUT)
```

```
/* Currently my_archive is tested before adding origin to
   anything. I believe that this can become always an add of
   origin, with origin set to 0 for non archive files. */
ufile_ptr origin;

/* The origin in the archive of the proxy entry. This will
   normally be the same as origin, except for thin archives,
   when it will contain the current offset of the proxy in the
   thin archive rather than the offset of the bfd in its actual
   container. */
ufile_ptr proxy_origin;

/* A hash table for section names. */
struct bfd_hash_table section_htab;

/* Pointer to linked list of sections. */
struct bfd_section *sections;

/* The last section on the section list. */
struct bfd_section *section_last;

/* The number of sections. */
unsigned int section_count;

/* Stuff only useful for object files:
   The start address. */
bfd_vma start_address;

/* Used for input and output. */
unsigned int symcount;

/* Symbol table for output BFD (with symcount entries).
   Also used by the linker to cache input BFD symbols. */
struct bfd_symbol **outsymbols;

/* Used for slurped dynamic symbol tables. */
unsigned int dynsymcount;

/* Pointer to structure which contains architecture information. */
const struct bfd_arch_info *arch_info;

/* Stuff only useful for archives. */
void *arelt_data;
struct bfd *my_archive; /* The containing archive BFD. */
struct bfd *archive_next; /* The next BFD in the archive. */
struct bfd *archive_head; /* The first BFD in the archive. */
struct bfd *nested_archives; /* List of nested archive in a flattened
```



```

                                thin archive.  */

/* A chain of BFD structures involved in a link.  */
struct bfd *link_next;

/* A field used by _bfd_generic_link_add_archive_symbols.  This will
   be used only for archive elements.  */
int archive_pass;

/* Used by the back end to hold private data.  */
union
{
    struct aout_data_struct *aout_data;
    struct artdat *aout_ar_data;
    struct _oasys_data *oasys_obj_data;
    struct _oasys_ar_data *oasys_ar_data;
    struct coff_tdata *coff_obj_data;
    struct pe_tdata *pe_obj_data;
    struct xcoff_tdata *xcoff_obj_data;
    struct ecoff_tdata *ecoff_obj_data;
    struct ieee_data_struct *ieee_data;
    struct ieee_ar_data_struct *ieee_ar_data;
    struct srec_data_struct *srec_data;
    struct verilog_data_struct *verilog_data;
    struct ihex_data_struct *ihex_data;
    struct tekhex_data_struct *tekhex_data;
    struct elf_obj_tdata *elf_obj_data;
    struct nlm_obj_tdata *nlm_obj_data;
    struct bout_data_struct *bout_data;
    struct mmo_data_struct *mmo_data;
    struct sun_core_struct *sun_core_data;
    struct sco5_core_struct *sco5_core_data;
    struct trad_core_struct *trad_core_data;
    struct som_data_struct *som_data;
    struct hpux_core_struct *hpux_core_data;
    struct hppabsd_core_struct *hppabsd_core_data;
    struct sgi_core_struct *sgi_core_data;
    struct lynx_core_struct *lynx_core_data;
    struct osf_core_struct *osf_core_data;
    struct cisco_core_struct *cisco_core_data;
    struct versados_data_struct *versados_data;
    struct netbsd_core_struct *netbsd_core_data;
    struct mach_o_data_struct *mach_o_data;
    struct mach_o_fat_data_struct *mach_o_fat_data;
    struct plugin_data_struct *plugin_data;
    struct bfd_pef_data_struct *pef_data;
    struct bfd_pef_xlib_data_struct *pef_xlib_data;

```

```

    struct bfd_sym_data_struct *sym_data;
    void *any;
}
tdata;

/* Used by the application to hold private data. */
void *usrdata;

/* Where all the allocated stuff under this BFD goes. This is a
   struct objalloc *, but we use void * to avoid requiring the inclusion
   of objalloc.h. */
void *memory;

/* Is the file descriptor being cached? That is, can it be closed as
   needed, and re-opened when accessed later? */
unsigned int cacheable : 1;

/* Marks whether there was a default target specified when the
   BFD was opened. This is used to select which matching algorithm
   to use to choose the back end. */
unsigned int target_defaulted : 1;

/* ... and here: ('once' means at least once). */
unsigned int opened_once : 1;

/* Set if we have a locally maintained mtime value, rather than
   getting it from the file each time. */
unsigned int mtime_set : 1;

/* Flag set if symbols from this BFD should not be exported. */
unsigned int no_export : 1;

/* Remember when output has begun, to stop strange things
   from happening. */
unsigned int output_has_begun : 1;

/* Have archive map. */
unsigned int has_armap : 1;

/* Set if this is a thin archive. */
unsigned int is_thin_archive : 1;

/* Set if only required symbols should be added in the link hash table for
   this object. Used by VMS linkers. */
unsigned int selective_search : 1;
};

```

2.2 Error reporting

Most BFD functions return nonzero on success (check their individual documentation for precise semantics). On an error, they call `bfd_set_error` to set an error condition that callers can check by calling `bfd_get_error`. If that returns `bfd_error_system_call`, then check `errno`.

The easiest way to report a BFD error to the user is to use `bfd_perror`.

2.2.1 Type `bfd_error_type`

The values returned by `bfd_get_error` are defined by the enumerated type `bfd_error_type`.

```
typedef enum bfd_error
{
    bfd_error_no_error = 0,
    bfd_error_system_call,
    bfd_error_invalid_target,
    bfd_error_wrong_format,
    bfd_error_wrong_object_format,
    bfd_error_invalid_operation,
    bfd_error_no_memory,
    bfd_error_no_symbols,
    bfd_error_no_armap,
    bfd_error_no_more_archived_files,
    bfd_error_malformed_archive,
    bfd_error_file_not_recognized,
    bfd_error_file_ambiguously_recognized,
    bfd_error_no_contents,
    bfd_error_nonrepresentable_section,
    bfd_error_no_debug_section,
    bfd_error_bad_value,
    bfd_error_file_truncated,
    bfd_error_file_too_big,
    bfd_error_on_input,
    bfd_error_invalid_error_code
}
bfd_error_type;
```

2.2.1.1 `bfd_get_error`

Synopsis

```
bfd_error_type bfd_get_error (void);
```

Description

Return the current BFD error condition.

2.2.1.2 bfd_set_error

Synopsis

```
void bfd_set_error (bfd_error_type error_tag, ...);
```

Description

Set the BFD error condition to be *error_tag*. If *error_tag* is `bfd_error_on_input`, then this function takes two more parameters, the input bfd where the error occurred, and the `bfd_error_type` error.

2.2.1.3 bfd_errmsg

Synopsis

```
const char *bfd_errmsg (bfd_error_type error_tag);
```

Description

Return a string describing the error *error_tag*, or the system error if *error_tag* is `bfd_error_system_call`.

2.2.1.4 bfd_perror

Synopsis

```
void bfd_perror (const char *message);
```

Description

Print to the standard error stream a string describing the last BFD error that occurred, or the last system error if the last BFD error was a system call failure. If *message* is non-NULL and non-empty, the error string printed is preceded by *message*, a colon, and a space. It is followed by a newline.

2.2.2 BFD error handler

Some BFD functions want to print messages describing the problem. They call a BFD error handler function. This function may be overridden by the program.

The BFD error handler acts like `printf`.

```
typedef void (*bfd_error_handler_type) (const char *, ...);
```

2.2.2.1 bfd_set_error_handler

Synopsis

```
bfd_error_handler_type bfd_set_error_handler (bfd_error_handler_type);
```

Description

Set the BFD error handler function. Returns the previous function.

2.2.2.2 bfd_set_error_program_name

Synopsis

```
void bfd_set_error_program_name (const char *);
```

Description

Set the program name to use when printing a BFD error. This is printed before the error

message followed by a colon and space. The string must not be changed after it is passed to this function.

2.2.2.3 bfd_get_error_handler

Synopsis

```
bfd_error_handler_type bfd_get_error_handler (void);
```

Description

Return the BFD error handler function.

2.3 Miscellaneous

2.3.1 Miscellaneous functions

2.3.1.1 bfd_get_reloc_upper_bound

Synopsis

```
long bfd_get_reloc_upper_bound (bfd *abfd, asection *sect);
```

Description

Return the number of bytes required to store the relocation information associated with section *sect* attached to bfd *abfd*. If an error occurs, return -1.

2.3.1.2 bfd_canonicalize_reloc

Synopsis

```
long bfd_canonicalize_reloc
(bfd *abfd, asection *sec, arelent **loc, asymbol **syms);
```

Description

Call the back end associated with the open BFD *abfd* and translate the external form of the relocation information attached to *sec* into the internal canonical form. Place the table into memory at *loc*, which has been preallocated, usually by a call to `bfd_get_reloc_upper_bound`. Returns the number of relocs, or -1 on error.

The *syms* table is also needed for horrible internal magic reasons.

2.3.1.3 bfd_set_reloc

Synopsis

```
void bfd_set_reloc
(bfd *abfd, asection *sec, arelent **rel, unsigned int count);
```

Description

Set the relocation pointer and count within section *sec* to the values *rel* and *count*. The argument *abfd* is ignored.

2.3.1.4 bfd_set_file_flags

Synopsis

```
bfd_boolean bfd_set_file_flags (bfd *abfd, flagword flags);
```

Description

Set the flag word in the BFD *abfd* to the value *flags*.

Possible errors are:

- `bfd_error_wrong_format` - The target bfd was not of object format.
- `bfd_error_invalid_operation` - The target bfd was open for reading.
- `bfd_error_invalid_operation` - The flag word contained a bit which was not applicable to the type of file. E.g., an attempt was made to set the `D_PAGED` bit on a BFD format which does not support demand paging.

2.3.1.5 `bfd_get_arch_size`

Synopsis

```
int bfd_get_arch_size (bfd *abfd);
```

Description

Returns the architecture address size, in bits, as determined by the object file's format. For ELF, this information is included in the header.

Returns

Returns the arch size in bits if known, -1 otherwise.

2.3.1.6 `bfd_get_sign_extend_vma`

Synopsis

```
int bfd_get_sign_extend_vma (bfd *abfd);
```

Description

Indicates if the target architecture "naturally" sign extends an address. Some architectures implicitly sign extend address values when they are converted to types larger than the size of an address. For instance, `bfd_get_start_address()` will return an address sign extended to fill a `bfd_vma` when this is the case.

Returns

Returns 1 if the target architecture is known to sign extend addresses, 0 if the target architecture is known to not sign extend addresses, and -1 otherwise.

2.3.1.7 `bfd_set_start_address`

Synopsis

```
bfd_boolean bfd_set_start_address (bfd *abfd, bfd_vma vma);
```

Description

Make *vma* the entry point of output BFD *abfd*.

Returns

Returns TRUE on success, FALSE otherwise.

2.3.1.8 `bfd_get_gp_size`

Synopsis

```
unsigned int bfd_get_gp_size (bfd *abfd);
```

Description

Return the maximum size of objects to be optimized using the GP register under MIPS ECOFF. This is typically set by the `-G` argument to the compiler, assembler or linker.

2.3.1.9 bfd_set_gp_size

Synopsis

```
void bfd_set_gp_size (bfd *abfd, unsigned int i);
```

Description

Set the maximum size of objects to be optimized using the GP register under ECOFF or MIPS ELF. This is typically set by the `-G` argument to the compiler, assembler or linker.

2.3.1.10 bfd_scan_vma

Synopsis

```
bfd_vma bfd_scan_vma (const char *string, const char **end, int base);
```

Description

Convert, like `strtoul`, a numerical expression *string* into a `bfd_vma` integer, and return that integer. (Though without as many bells and whistles as `strtoul`.) The expression is assumed to be unsigned (i.e., positive). If given a *base*, it is used as the base for conversion. A base of 0 causes the function to interpret the string in hex if a leading "0x" or "0X" is found, otherwise in octal if a leading zero is found, otherwise in decimal.

If the value would overflow, the maximum `bfd_vma` value is returned.

2.3.1.11 bfd_copy_private_header_data

Synopsis

```
bfd_boolean bfd_copy_private_header_data (bfd *ibfd, bfd *obfd);
```

Description

Copy private BFD header information from the BFD *ibfd* to the the BFD *obfd*. This copies information that may require sections to exist, but does not require symbol tables. Return `true` on success, `false` on error. Possible error returns are:

- `bfd_error_no_memory` - Not enough memory exists to create private data for *obfd*.

```
#define bfd_copy_private_header_data(ibfd, obfd) \
    BFD_SEND (obfd, _bfd_copy_private_header_data, \
              (ibfd, obfd))
```

2.3.1.12 bfd_copy_private_bfd_data

Synopsis

```
bfd_boolean bfd_copy_private_bfd_data (bfd *ibfd, bfd *obfd);
```

Description

Copy private BFD information from the BFD *ibfd* to the the BFD *obfd*. Return `TRUE` on success, `FALSE` on error. Possible error returns are:

- `bfd_error_no_memory` - Not enough memory exists to create private data for *obfd*.

```
#define bfd_copy_private_bfd_data(ibfd, obfd) \
    BFD_SEND (obfd, _bfd_copy_private_bfd_data, \
              (ibfd, obfd))
```

2.3.1.13 bfd_merge_private_bfd_data

Synopsis

```
bfd_boolean bfd_merge_private_bfd_data (bfd *ibfd, bfd *obfd);
```

Description

Merge private BFD information from the BFD *ibfd* to the the output file BFD *obfd* when linking. Return TRUE on success, FALSE on error. Possible error returns are:

- `bfd_error_no_memory` - Not enough memory exists to create private data for *obfd*.

```
#define bfd_merge_private_bfd_data(ibfd, obfd) \
    BFD_SEND (obfd, _bfd_merge_private_bfd_data, \
              (ibfd, obfd))
```

2.3.1.14 bfd_set_private_flags

Synopsis

```
bfd_boolean bfd_set_private_flags (bfd *abfd, flagword flags);
```

Description

Set private BFD flag information in the BFD *abfd*. Return TRUE on success, FALSE on error. Possible error returns are:

- `bfd_error_no_memory` - Not enough memory exists to create private data for *obfd*.

```
#define bfd_set_private_flags(abfd, flags) \
    BFD_SEND (abfd, _bfd_set_private_flags, (abfd, flags))
```

2.3.1.15 Other functions

Description

The following functions exist but have not yet been documented.

```
#define bfd_sizeof_headers(abfd, info) \
    BFD_SEND (abfd, _bfd_sizeof_headers, (abfd, info))

#define bfd_find_nearest_line(abfd, sec, syms, off, file, func, line) \
    BFD_SEND (abfd, _bfd_find_nearest_line, \
              (abfd, sec, syms, off, file, func, line))

#define bfd_find_line(abfd, syms, sym, file, line) \
    BFD_SEND (abfd, _bfd_find_line, \
              (abfd, syms, sym, file, line))

#define bfd_find_inliner_info(abfd, file, func, line) \
    BFD_SEND (abfd, _bfd_find_inliner_info, \
              (abfd, file, func, line))

#define bfd_debug_info_start(abfd) \
    BFD_SEND (abfd, _bfd_debug_info_start, (abfd))

#define bfd_debug_info_end(abfd) \
    BFD_SEND (abfd, _bfd_debug_info_end, (abfd))
```



```
#define bfd_debug_info_accumulate(abfd, section) \
    BFD_SEND (abfd, _bfd_debug_info_accumulate, (abfd, section))

#define bfd_stat_arch_elt(abfd, stat) \
    BFD_SEND (abfd, _bfd_stat_arch_elt, (abfd, stat))

#define bfd_update_armap_timestamp(abfd) \
    BFD_SEND (abfd, _bfd_update_armap_timestamp, (abfd))

#define bfd_set_arch_mach(abfd, arch, mach) \
    BFD_SEND ( abfd, _bfd_set_arch_mach, (abfd, arch, mach))

#define bfd_relax_section(abfd, section, link_info, again) \
    BFD_SEND (abfd, _bfd_relax_section, (abfd, section, link_info, again))■

#define bfd_gc_sections(abfd, link_info) \
    BFD_SEND (abfd, _bfd_gc_sections, (abfd, link_info))

#define bfd_merge_sections(abfd, link_info) \
    BFD_SEND (abfd, _bfd_merge_sections, (abfd, link_info))

#define bfd_is_group_section(abfd, sec) \
    BFD_SEND (abfd, _bfd_is_group_section, (abfd, sec))

#define bfd_discard_group(abfd, sec) \
    BFD_SEND (abfd, _bfd_discard_group, (abfd, sec))

#define bfd_link_hash_table_create(abfd) \
    BFD_SEND (abfd, _bfd_link_hash_table_create, (abfd))

#define bfd_link_hash_table_free(abfd, hash) \
    BFD_SEND (abfd, _bfd_link_hash_table_free, (hash))

#define bfd_link_add_symbols(abfd, info) \
    BFD_SEND (abfd, _bfd_link_add_symbols, (abfd, info))

#define bfd_link_just_syms(abfd, sec, info) \
    BFD_SEND (abfd, _bfd_link_just_syms, (sec, info))

#define bfd_final_link(abfd, info) \
    BFD_SEND (abfd, _bfd_final_link, (abfd, info))

#define bfd_free_cached_info(abfd) \
    BFD_SEND (abfd, _bfd_free_cached_info, (abfd))

#define bfd_get_dynamic_symtab_upper_bound(abfd) \
```

```

        BFD_SEND (abfd, _bfd_get_dynamic_symtab_upper_bound, (abfd))

#define bfd_print_private_bfd_data(abfd, file)\
        BFD_SEND (abfd, _bfd_print_private_bfd_data, (abfd, file))

#define bfd_canonicalize_dynamic_symtab(abfd, asymbols) \
        BFD_SEND (abfd, _bfd_canonicalize_dynamic_symtab, (abfd, asymbols))■

#define bfd_get_synthetic_symtab(abfd, count, syms, dyncount, dynsyms, ret) \
        BFD_SEND (abfd, _bfd_get_synthetic_symtab, (abfd, count, syms, \
                                                    dyncount, dynsyms, ret))■

#define bfd_get_dynamic_reloc_upper_bound(abfd) \
        BFD_SEND (abfd, _bfd_get_dynamic_reloc_upper_bound, (abfd))

#define bfd_canonicalize_dynamic_reloc(abfd, arels, asyms) \
        BFD_SEND (abfd, _bfd_canonicalize_dynamic_reloc, (abfd, arels, asyms))■

extern bfd_byte *bfd_get_relocated_section_contents
(bfd *, struct bfd_link_info *, struct bfd_link_order *, bfd_byte *,
  bfd_boolean, asymbol **);

```

2.3.1.16 bfd_alt_mach_code

Synopsis

```

bfd_boolean bfd_alt_mach_code (bfd *abfd, int alternative);

```

Description

When more than one machine code number is available for the same machine type, this function can be used to switch between the preferred one (`alternative == 0`) and any others. Currently, only ELF supports this feature, with up to two alternate machine codes.

```

struct bfd_preserve
{
    void *marker;
    void *tdata;
    flagword flags;
    const struct bfd_arch_info *arch_info;
    struct bfd_section *sections;
    struct bfd_section *section_last;
    unsigned int section_count;
    struct bfd_hash_table section_htab;
};

```

2.3.1.17 bfd_preserve_save

Synopsis

```

bfd_boolean bfd_preserve_save (bfd *, struct bfd_preserve *);

```

Description

When testing an object for compatibility with a particular target back-end, the back-end `object_p` function needs to set up certain fields in the bfd on successfully recognizing the object. This typically happens in a piecemeal fashion, with failures possible at many points. On failure, the bfd is supposed to be restored to its initial state, which is virtually impossible. However, restoring a subset of the bfd state works in practice. This function stores the subset and reinitializes the bfd.

2.3.1.18 bfd_preserve_restore**Synopsis**

```
void bfd_preserve_restore (bfd *, struct bfd_preserve *);
```

Description

This function restores bfd state saved by `bfd_preserve_save`. If `MARKER` is non-NULL in struct `bfd_preserve` then that block and all subsequently `bfd_alloc`'d memory is freed.

2.3.1.19 bfd_preserve_finish**Synopsis**

```
void bfd_preserve_finish (bfd *, struct bfd_preserve *);
```

Description

This function should be called when the bfd state saved by `bfd_preserve_save` is no longer needed. ie. when the back-end `object_p` function returns with success.

2.3.1.20 bfd_emul_get_maxpagesize**Synopsis**

```
bfd_vma bfd_emul_get_maxpagesize (const char *);
```

Description

Returns the maximum page size, in bytes, as determined by emulation.

Returns

Returns the maximum page size in bytes for ELF, 0 otherwise.

2.3.1.21 bfd_emul_set_maxpagesize**Synopsis**

```
void bfd_emul_set_maxpagesize (const char *, bfd_vma);
```

Description

For ELF, set the maximum page size for the emulation. It is a no-op for other formats.

2.3.1.22 bfd_emul_get_commonpagesize**Synopsis**

```
bfd_vma bfd_emul_get_commonpagesize (const char *);
```

Description

Returns the common page size, in bytes, as determined by emulation.

Returns

Returns the common page size in bytes for ELF, 0 otherwise.

2.3.1.23 bfd_emul_set_commonpagesize

Synopsis

```
void bfd_emul_set_commonpagesize (const char *, bfd_vma);
```

Description

For ELF, set the common page size for the emulation. It is a no-op for other formats.

2.3.1.24 bfd_demangle

Synopsis

```
char *bfd_demangle (bfd *, const char *, int);
```

Description

Wrapper around cplus_demangle. Strips leading underscores and other such chars that would otherwise confuse the demangler. If passed a g++ v3 ABI mangled name, returns a buffer allocated with malloc holding the demangled name. Returns NULL otherwise and on memory alloc failure.

2.3.1.25 struct bfd_iovec

Description

The struct bfd_iovec contains the internal file I/O class. Each BFD has an instance of this class and all file I/O is routed through it (it is assumed that the instance implements all methods listed below).

```
struct bfd_iovec
{
    /* To avoid problems with macros, a "b" rather than "f"
       prefix is prepended to each method name. */
    /* Attempt to read/write NBYTES on ABFD's IOSTREAM storing/fetching
       bytes starting at PTR. Return the number of bytes actually
       transfered (a read past end-of-file returns less than NBYTES),
       or -1 (setting bfd_error) if an error occurs. */
    file_ptr (*bread) (struct bfd *abfd, void *ptr, file_ptr nbytes);
    file_ptr (*bwrite) (struct bfd *abfd, const void *ptr,
                        file_ptr nbytes);
    /* Return the current IOSTREAM file offset, or -1 (setting bfd_error
       if an error occurs. */
    file_ptr (*btell) (struct bfd *abfd);
    /* For the following, on successful completion a value of 0 is returned.
       Otherwise, a value of -1 is returned (and bfd_error is set). */
    int (*bseek) (struct bfd *abfd, file_ptr offset, int whence);
    int (*bclose) (struct bfd *abfd);
    int (*bflush) (struct bfd *abfd);
    int (*bstat) (struct bfd *abfd, struct stat *sb);
    /* Just like mmap: (void*)-1 on failure, mmaped address on success. */
    void *(*bmmap) (struct bfd *abfd, void *addr, bfd_size_type len,
                   int prot, int flags, file_ptr offset);
};

extern const struct bfd_iovec _bfd_memory_iovec;
```

2.3.1.26 `bfd_get_mtime`

Synopsis

```
long bfd_get_mtime (bfd *abfd);
```

Description

Return the file modification time (as read from the file system, or from the archive header for archive members).

2.3.1.27 `bfd_get_size`

Synopsis

```
file_ptr bfd_get_size (bfd *abfd);
```

Description

Return the file size (as read from file system) for the file associated with BFD *abfd*.

The initial motivation for, and use of, this routine is not so we can get the exact size of the object the BFD applies to, since that might not be generally possible (archive members for example). It would be ideal if someone could eventually modify it so that such results were guaranteed.

Instead, we want to ask questions like "is this NNN byte sized object I'm about to try read from file offset YYY reasonable?" As an example of where we might do this, some object formats use string tables for which the first `sizeof (long)` bytes of the table contain the size of the table itself, including the size bytes. If an application tries to read what it thinks is one of these string tables, without some way to validate the size, and for some reason the size is wrong (byte swapping error, wrong location for the string table, etc.), the only clue is likely to be a read error when it tries to read the table, or a "virtual memory exhausted" error when it tries to allocate 15 bazillion bytes of space for the 15 bazillion byte table it is about to read. This function at least allows us to answer the question, "is the size reasonable?".

2.3.1.28 `bfd_mmap`

Synopsis

```
void *bfd_mmap (bfd *abfd, void *addr, bfd_size_type len,  
               int prot, int flags, file_ptr offset);
```

Description

Return `mmap()`ed region of the file, if possible and implemented.

2.4 Memory Usage

BFD keeps all of its internal structures in obstacks. There is one obstack per open BFD file, into which the current state is stored. When a BFD is closed, the obstack is deleted, and so everything which has been allocated by BFD for the closing file is thrown away.

BFD does not free anything created by an application, but pointers into `bfd` structures become invalid on a `bfd_close`; for example, after a `bfd_close` the vector passed to `bfd_canonicalize_syntab` is still around, since it has been allocated by the application, but the data that it pointed to are lost.

The general rule is to not close a BFD until all operations dependent upon data from the BFD have been completed, or all the data from within the file has been copied. To help

with the management of memory, there is a function (`bfd_alloc_size`) which returns the number of bytes in obstacks associated with the supplied BFD. This could be used to select the greediest open BFD, close it to reclaim the memory, perform some operation and reopen the BFD again, to get a fresh copy of the data structures.

2.5 Initialization

2.5.1 Initialization functions

These are the functions that handle initializing a BFD.

2.5.1.1 `bfd_init`

Synopsis

```
void bfd_init (void);
```

Description

This routine must be called before any other BFD function to initialize magical internal data structures.

2.6 Sections

The raw data contained within a BFD is maintained through the section abstraction. A single BFD may have any number of sections. It keeps hold of them by pointing to the first; each one points to the next in the list.

Sections are supported in BFD in `section.c`.

2.6.1 Section input

When a BFD is opened for reading, the section structures are created and attached to the BFD.

Each section has a name which describes the section in the outside world—for example, `a.out` would contain at least three sections, called `.text`, `.data` and `.bss`.

Names need not be unique; for example a COFF file may have several sections named `.data`. Sometimes a BFD will contain more than the “natural” number of sections. A back end may attach other sections containing constructor data, or an application may add a section (using `bfd_make_section`) to the sections attached to an already open BFD. For example, the linker creates an extra section `COMMON` for each input file’s BFD to hold information about common storage.

The raw data is not necessarily read in when the section descriptor is created. Some targets may leave the data in place until a `bfd_get_section_contents` call is made. Other back ends may read in all the data at once. For example, an S-record file has to be read once to determine the size of the data. An IEEE-695 file doesn’t contain raw data in sections, but data and relocation expressions intermixed, so the data area has to be parsed to get out the data and relocations.

2.6.2 Section output

To write a new object style BFD, the various sections to be written have to be created. They are attached to the BFD in the same way as input sections; data is written to the sections using `bfd_set_section_contents`.

Any program that creates or combines sections (e.g., the assembler and linker) must use the `asection` fields `output_section` and `output_offset` to indicate the file sections to which each section must be written. (If the section is being created from scratch, `output_section` should probably point to the section itself and `output_offset` should probably be zero.)

The data to be written comes from input sections attached (via `output_section` pointers) to the output sections. The output section structure can be considered a filter for the input section: the output section determines the vma of the output data and the name, but the input section determines the offset into the output section of the data to be written.

E.g., to create a section "O", starting at 0x100, 0x123 long, containing two subsections, "A" at offset 0x0 (i.e., at vma 0x100) and "B" at offset 0x20 (i.e., at vma 0x120) the `asection` structures would look like:

section name	"A"		
output_offset	0x00		
size	0x20		
output_section	----->	section name	"O"
		vma	0x100
section name	"B"	size	0x123
output_offset	0x20		
size	0x103		
output_section	-----		

2.6.3 Link orders

The data within a section is stored in a *link_order*. These are much like the fixups in *gas*. The *link_order* abstraction allows a section to grow and shrink within itself.

A *link_order* knows how big it is, and which is the next *link_order* and where the raw data for it is; it also points to a list of relocations which apply to it.

The *link_order* is used by the linker to perform relaxing on final code. The compiler creates code which is as big as necessary to make it work without relaxing, and the user can select whether to relax. Sometimes relaxing takes a lot of time. The linker runs around the relocations to see if any are attached to data which can be shrunk, if so it does it on a *link_order* by *link_order* basis.

2.6.4 typedef asection

Here is the section structure:

```
typedef struct bfd_section
{
    /* The name of the section; the name isn't a copy, the pointer is
       the same as that passed to bfd_make_section.  */
    const char *name;

    /* A unique sequence number.  */
    int id;

    /* Which section in the bfd; 0..n-1 as sections are created in a bfd.  */
    int index;
};
```

```

int index;

/* The next section in the list belonging to the BFD, or NULL. */
struct bfd_section *next;

/* The previous section in the list belonging to the BFD, or NULL. */
struct bfd_section *prev;

/* The field flags contains attributes of the section. Some
   flags are read in from the object file, and some are
   synthesized from other information. */
flagword flags;

#define SEC_NO_FLAGS    0x000

/* Tells the OS to allocate space for this section when loading.
   This is clear for a section containing debug information only. */
#define SEC_ALLOC      0x001

/* Tells the OS to load the section from the file when loading.
   This is clear for a .bss section. */
#define SEC_LOAD       0x002

/* The section contains data still to be relocated, so there is
   some relocation information too. */
#define SEC_RELOC      0x004

/* A signal to the OS that the section contains read only data. */
#define SEC_READONLY   0x008

/* The section contains code only. */
#define SEC_CODE       0x010

/* The section contains data only. */
#define SEC_DATA       0x020

/* The section will reside in ROM. */
#define SEC_ROM        0x040

/* The section contains constructor information. This section
   type is used by the linker to create lists of constructors and
   destructors used by g++. When a back end sees a symbol
   which should be used in a constructor list, it creates a new
   section for the type of name (e.g., __CTOR_LIST__), attaches
   the symbol to it, and builds a relocation. To build the lists
   of constructors, all the linker has to do is catenate all the
   sections called __CTOR_LIST__ and relocate the data

```



```

        contained within - exactly the operations it would perform on
        standard data.  */
#define SEC_CONSTRUCTOR 0x080

    /* The section has contents - a data section could be
       SEC_ALLOC | SEC_HAS_CONTENTS; a debug section could be
       SEC_HAS_CONTENTS  */
#define SEC_HAS_CONTENTS 0x100

    /* An instruction to the linker to not output the section
       even if it has information which would normally be written.  */
#define SEC_NEVER_LOAD 0x200

    /* The section contains thread local data.  */
#define SEC_THREAD_LOCAL 0x400

    /* The section has GOT references.  This flag is only for the
       linker, and is currently only used by the elf32-hppa back end.
       It will be set if global offset table references were detected
       in this section, which indicate to the linker that the section
       contains PIC code, and must be handled specially when doing a
       static link.  */
#define SEC_HAS_GOT_REF 0x800

    /* The section contains common symbols (symbols may be defined
       multiple times, the value of a symbol is the amount of
       space it requires, and the largest symbol value is the one
       used).  Most targets have exactly one of these (which we
       translate to bfd_com_section_ptr), but ECOFF has two.  */
#define SEC_IS_COMMON 0x1000

    /* The section contains only debugging information.  For
       example, this is set for ELF .debug and .stab sections.
       strip tests this flag to see if a section can be
       discarded.  */
#define SEC_DEBUGGING 0x2000

    /* The contents of this section are held in memory pointed to
       by the contents field.  This is checked by bfd_get_section_contents,
       and the data is retrieved from memory if appropriate.  */
#define SEC_IN_MEMORY 0x4000

    /* The contents of this section are to be excluded by the
       linker for executable and shared objects unless those
       objects are to be further relocated.  */
#define SEC_EXCLUDE 0x8000

```

```

/* The contents of this section are to be sorted based on the sum of
   the symbol and addend values specified by the associated relocation
   entries. Entries without associated relocation entries will be
   appended to the end of the section in an unspecified order. */
#define SEC_SORT_ENTRIES 0x10000

/* When linking, duplicate sections of the same name should be
   discarded, rather than being combined into a single section as
   is usually done. This is similar to how common symbols are
   handled. See SEC_LINK_DUPLICATES below. */
#define SEC_LINK_ONCE 0x20000

/* If SEC_LINK_ONCE is set, this bitfield describes how the linker
   should handle duplicate sections. */
#define SEC_LINK_DUPLICATES 0xc0000

/* This value for SEC_LINK_DUPLICATES means that duplicate
   sections with the same name should simply be discarded. */
#define SEC_LINK_DUPLICATES_DISCARD 0x0

/* This value for SEC_LINK_DUPLICATES means that the linker
   should warn if there are any duplicate sections, although
   it should still only link one copy. */
#define SEC_LINK_DUPLICATES_ONE_ONLY 0x40000

/* This value for SEC_LINK_DUPLICATES means that the linker
   should warn if any duplicate sections are a different size. */
#define SEC_LINK_DUPLICATES_SAME_SIZE 0x80000

/* This value for SEC_LINK_DUPLICATES means that the linker
   should warn if any duplicate sections contain different
   contents. */
#define SEC_LINK_DUPLICATES_SAME_CONTENTS \
  (SEC_LINK_DUPLICATES_ONE_ONLY | SEC_LINK_DUPLICATES_SAME_SIZE)

/* This section was created by the linker as part of dynamic
   relocation or other arcane processing. It is skipped when
   going through the first-pass output, trusting that someone
   else up the line will take care of it later. */
#define SEC_LINKER_CREATED 0x100000

/* This section should not be subject to garbage collection.
   Also set to inform the linker that this section should not be
   listed in the link map as discarded. */
#define SEC_KEEP 0x200000

/* This section contains "short" data, and should be placed

```

```

    "near" the GP. */
#define SEC_SMALL_DATA 0x400000

    /* Attempt to merge identical entities in the section.
       Entity size is given in the entsize field. */
#define SEC_MERGE 0x800000

    /* If given with SEC_MERGE, entities to merge are zero terminated
       strings where entsize specifies character size instead of fixed
       size entries. */
#define SEC_STRINGS 0x1000000

    /* This section contains data about section groups. */
#define SEC_GROUP 0x2000000

    /* The section is a COFF shared library section. This flag is
       only for the linker. If this type of section appears in
       the input file, the linker must copy it to the output file
       without changing the vma or size. FIXME: Although this
       was originally intended to be general, it really is COFF
       specific (and the flag was renamed to indicate this). It
       might be cleaner to have some more general mechanism to
       allow the back end to control what the linker does with
       sections. */
#define SEC_COFF_SHARED_LIBRARY 0x4000000

    /* This section contains data which may be shared with other
       executables or shared objects. This is for COFF only. */
#define SEC_COFF_SHARED 0x8000000

    /* When a section with this flag is being linked, then if the size of
       the input section is less than a page, it should not cross a page
       boundary. If the size of the input section is one page or more,
       it should be aligned on a page boundary. This is for TI
       TMS320C54X only. */
#define SEC_TIC54X_BLOCK 0x10000000

    /* Conditionally link this section; do not link if there are no
       references found to any symbol in the section. This is for TI
       TMS320C54X only. */
#define SEC_TIC54X_CLINK 0x20000000

    /* Indicate that section has the no read flag set. This happens
       when memory read flag isn't set. */
#define SEC_COFF_NOREAD 0x40000000

    /* End of section flags. */

```

```

/* Some internal packed boolean fields. */

/* See the vma field. */
unsigned int user_set_vma : 1;

/* A mark flag used by some of the linker backends. */
unsigned int linker_mark : 1;

/* Another mark flag used by some of the linker backends. Set for
   output sections that have an input section. */
unsigned int linker_has_input : 1;

/* Mark flag used by some linker backends for garbage collection. */
unsigned int gc_mark : 1;

/* Section compression status. */
unsigned int compress_status : 2;
#define COMPRESS_SECTION_NONE 0
#define COMPRESS_SECTION_DONE 1
#define DECOMPRESS_SECTION_SIZED 2

/* The following flags are used by the ELF linker. */

/* Mark sections which have been allocated to segments. */
unsigned int segment_mark : 1;

/* Type of sec_info information. */
unsigned int sec_info_type:3;
#define ELF_INFO_TYPE_NONE 0
#define ELF_INFO_TYPE_STABS 1
#define ELF_INFO_TYPE_MERGE 2
#define ELF_INFO_TYPE_EH_FRAME 3
#define ELF_INFO_TYPE_JUST_SYMS 4

/* Nonzero if this section uses RELA relocations, rather than REL. */
unsigned int use_rela_p:1;

/* Bits used by various backends. The generic code doesn't touch
   these fields. */

unsigned int sec_flg0:1;
unsigned int sec_flg1:1;
unsigned int sec_flg2:1;
unsigned int sec_flg3:1;
unsigned int sec_flg4:1;
unsigned int sec_flg5:1;

```

```
/* End of internal packed boolean fields. */

/* The virtual memory address of the section - where it will be
   at run time. The symbols are relocated against this. The
   user_set_vma flag is maintained by bfd; if it's not set, the
   backend can assign addresses (for example, in a.out, where
   the default address for .data is dependent on the specific
   target and various flags). */
bfd_vma vma;

/* The load address of the section - where it would be in a
   rom image; really only used for writing section header
   information. */
bfd_vma lma;

/* The size of the section in octets, as it will be output.
   Contains a value even if the section has no contents (e.g., the
   size of .bss). */
bfd_size_type size;

/* For input sections, the original size on disk of the section, in
   octets. This field should be set for any section whose size is
   changed by linker relaxation. It is required for sections where
   the linker relaxation scheme doesn't cache altered section and
   reloc contents (stabs, eh_frame, SEC_MERGE, some coff relaxing
   targets), and thus the original size needs to be kept to read the
   section multiple times. For output sections, rawsize holds the
   section size calculated on a previous linker relaxation pass. */
bfd_size_type rawsize;

/* The compressed size of the section in octets. */
bfd_size_type compressed_size;

/* Relaxation table. */
struct relax_table *relax;

/* Count of used relaxation table entries. */
int relax_count;

/* If this section is going to be output, then this value is the
   offset in *bytes* into the output section of the first byte in the
   input section (byte ==> smallest addressable unit on the
   target). In most cases, if this was going to start at the
   100th octet (8-bit quantity) in the output section, this value
   would be 100. However, if the target byte size is 16 bits
```

```
(bfd_octets_per_byte is "2"), this value would be 50. */
bfd_vma output_offset;

/* The output section through which to map on output. */
struct bfd_section *output_section;

/* The alignment requirement of the section, as an exponent of 2 -
   e.g., 3 aligns to 2^3 (or 8). */
unsigned int alignment_power;

/* If an input section, a pointer to a vector of relocation
   records for the data in this section. */
struct reloc_cache_entry *relocation;

/* If an output section, a pointer to a vector of pointers to
   relocation records for the data in this section. */
struct reloc_cache_entry **orelocation;

/* The number of relocation records in one of the above. */
unsigned reloc_count;

/* Information below is back end specific - and not always used
   or updated. */

/* File position of section data. */
file_ptr filepos;

/* File position of relocation info. */
file_ptr rel_filepos;

/* File position of line data. */
file_ptr line_filepos;

/* Pointer to data for applications. */
void *userdata;

/* If the SEC_IN_MEMORY flag is set, this points to the actual
   contents. */
unsigned char *contents;

/* Attached line number information. */
alint *lineno;

/* Number of line number records. */
unsigned int lineno_count;

/* Entity size for merging purposes. */
```

```

    unsigned int entsize;

    /* Points to the kept section if this section is a link-once section,
       and is discarded. */
    struct bfd_section *kept_section;

    /* When a section is being output, this value changes as more
       linenumbers are written out. */
    file_ptr moving_line_filepos;

    /* What the section number is in the target world. */
    int target_index;

    void *used_by_bfd;

    /* If this is a constructor section then here is a list of the
       relocations created to relocate items within it. */
    struct relent_chain *constructor_chain;

    /* The BFD which owns the section. */
    bfd *owner;

    /* A symbol which points at this section only. */
    struct bfd_symbol *symbol;
    struct bfd_symbol **symbol_ptr_ptr;

    /* Early in the link process, map_head and map_tail are used to build
       a list of input sections attached to an output section. Later,
       output sections use these fields for a list of bfd_link_order
       structs. */
    union {
        struct bfd_link_order *link_order;
        struct bfd_section *s;
    } map_head, map_tail;
} asection;

/* Relax table contains information about instructions which can
   be removed by relaxation -- replacing a long address with a
   short address. */
struct relax_table {
    /* Address where bytes may be deleted. */
    bfd_vma addr;

    /* Number of bytes to be deleted. */
    int size;
};

```

```

/* These sections are global, and are managed by BFD.  The application
   and target back end are not permitted to change the values in
   these sections.  New code should use the section_ptr macros rather
   than referring directly to the const sections.  The const sections
   may eventually vanish.  */
#define BFD_ABS_SECTION_NAME "*ABS*"
#define BFD_UND_SECTION_NAME "*UND*"
#define BFD_COM_SECTION_NAME "*COM*"
#define BFD_IND_SECTION_NAME "*IND*"

/* The absolute section.  */
extern asection bfd_abs_section;
#define bfd_abs_section_ptr ((asection *) &bfd_abs_section)
#define bfd_is_abs_section(sec) ((sec) == bfd_abs_section_ptr)
/* Pointer to the undefined section.  */
extern asection bfd_und_section;
#define bfd_und_section_ptr ((asection *) &bfd_und_section)
#define bfd_is_und_section(sec) ((sec) == bfd_und_section_ptr)
/* Pointer to the common section.  */
extern asection bfd_com_section;
#define bfd_com_section_ptr ((asection *) &bfd_com_section)
/* Pointer to the indirect section.  */
extern asection bfd_ind_section;
#define bfd_ind_section_ptr ((asection *) &bfd_ind_section)
#define bfd_is_ind_section(sec) ((sec) == bfd_ind_section_ptr)

#define bfd_is_const_section(SEC) \
  ( ((SEC) == bfd_abs_section_ptr) \
    || ((SEC) == bfd_und_section_ptr) \
    || ((SEC) == bfd_com_section_ptr) \
    || ((SEC) == bfd_ind_section_ptr))

/* Macros to handle insertion and deletion of a bfd's sections.  These
   only handle the list pointers, ie. do not adjust section_count,
   target_index etc.  */
#define bfd_section_list_remove(ABFD, S) \
do \
{ \
    asection *_s = S; \
    asection *_next = _s->next; \
    asection *_prev = _s->prev; \
    if (_prev) \
        _prev->next = _next; \
    else \
        (ABFD)->sections = _next; \
    if (_next) \
        _next->prev = _prev; \
}

```



```

        else
            (ABFD)->section_last = _prev;
    }
    while (0)
#define bfd_section_list_append(ABFD, S) \
do
{
    asection *_s = S;
    bfd *_abfd = ABFD;
    _s->next = NULL;
    if (_abfd->section_last)
    {
        _s->prev = _abfd->section_last;
        _abfd->section_last->next = _s;
    }
    else
    {
        _s->prev = NULL;
        _abfd->sections = _s;
    }
    _abfd->section_last = _s;
}
while (0)
#define bfd_section_list_prepend(ABFD, S) \
do
{
    asection *_s = S;
    bfd *_abfd = ABFD;
    _s->prev = NULL;
    if (_abfd->sections)
    {
        _s->next = _abfd->sections;
        _abfd->sections->prev = _s;
    }
    else
    {
        _s->next = NULL;
        _abfd->section_last = _s;
    }
    _abfd->sections = _s;
}
while (0)
#define bfd_section_list_insert_after(ABFD, A, S) \
do
{
    asection *_a = A;
    asection *_s = S;

```

```

    asection *_next = _a->next;
    _s->next = _next;
    _s->prev = _a;
    _a->next = _s;
    if (_next)
        _next->prev = _s;
    else
        (ABFD)->section_last = _s;
}
while (0)
#define bfd_section_list_insert_before(ABFD, B, S) \
do
{
    asection *_b = B;
    asection *_s = S;
    asection *_prev = _b->prev;
    _s->prev = _prev;
    _s->next = _b;
    _b->prev = _s;
    if (_prev)
        _prev->next = _s;
    else
        (ABFD)->sections = _s;
}
while (0)
#define bfd_section_removed_from_list(ABFD, S) \
((S)->next == NULL ? (ABFD)->section_last != (S) : (S)->next->prev != (S))■

#define BFD_FAKE_SECTION(SEC, FLAGS, SYM, NAME, IDX) \
/* name, id, index, next, prev, flags, user_set_vma, */ \
{ NAME, IDX, 0, NULL, NULL, FLAGS, 0, \
\
/* linker_mark, linker_has_input, gc_mark, decompress_status, */ \
0, 0, 1, 0, \
\
/* segment_mark, sec_info_type, use_rela_p, */ \
0, 0, 0, \
\
/* sec_flg0, sec_flg1, sec_flg2, sec_flg3, sec_flg4, sec_flg5, */ \
0, 0, 0, 0, 0, 0, \
\
/* vma, lma, size, rawsize, compressed_size, relax, relax_count, */ \
0, 0, 0, 0, 0, 0, 0, \
\
/* output_offset, output_section, alignment_power, */ \
0, (struct bfd_section *) &SEC, 0, \
\

```

```

/* relocation, orelocation, reloc_count, filepos, rel_filepos, */ \
    NULL,          NULL,          0,          0,          0,      \
                                                                \
/* line_filepos, userdata, contents, lineno, lineno_count, */ \
    0,             NULL,          NULL,          NULL,    0,      \
                                                                \
/* entsize, kept_section, moving_line_filepos, */ \
    0,             NULL,          0,              \
                                                                \
/* target_index, used_by_bfd, constructor_chain, owner, */ \
    0,             NULL,          NULL,              NULL,      \
                                                                \
/* symbol,                                symbol_ptr_ptr, */ \
    (struct bfd_symbol *) SYM, &SEC.symbol,          \
                                                                \
/* map_head, map_tail */ \
    { NULL }, { NULL } \
}

```

2.6.5 Section prototypes

These are the functions exported by the section handling part of BFD.

2.6.5.1 bfd_section_list_clear

Synopsis

```
void bfd_section_list_clear (bfd *);
```

Description

Clears the section list, and also resets the section count and hash table entries.

2.6.5.2 bfd_get_section_by_name

Synopsis

```
asection *bfd_get_section_by_name (bfd *abfd, const char *name);
```

Description

Run through *abfd* and return the one of the *asections* whose name matches *name*, otherwise NULL. See [Section 2.6 \[Sections\], page 22](#), for more information.

This should only be used in special cases; the normal way to process all sections of a given name is to use `bfd_map_over_sections` and `strcmp` on the name (or better yet, base it on the section flags or something else) for each section.

2.6.5.3 bfd_get_section_by_name_if

Synopsis

```
asection *bfd_get_section_by_name_if
(bfd *abfd,
 const char *name,
 bfd_boolean (*func) (bfd *abfd, asection *sect, void *obj),
```

```
void *obj);
```

Description

Call the provided function *func* for each section attached to the BFD *abfd* whose name matches *name*, passing *obj* as an argument. The function will be called as if by

```
func (abfd, the_section, obj);
```

It returns the first section for which *func* returns true, otherwise NULL.

2.6.5.4 bfd_get_unique_section_name

Synopsis

```
char *bfd_get_unique_section_name
(bfd *abfd, const char *templat, int *count);
```

Description

Invent a section name that is unique in *abfd* by tacking a dot and a digit suffix onto the original *templat*. If *count* is non-NULL, then it specifies the first number tried as a suffix to generate a unique name. The value pointed to by *count* will be incremented in this case.

2.6.5.5 bfd_make_section_old_way

Synopsis

```
asection *bfd_make_section_old_way (bfd *abfd, const char *name);
```

Description

Create a new empty section called *name* and attach it to the end of the chain of sections for the BFD *abfd*. An attempt to create a section with a name which is already in use returns its pointer without changing the section chain.

It has the funny name since this is the way it used to be before it was rewritten....

Possible errors are:

- `bfd_error_invalid_operation` - If output has already started for this BFD.
- `bfd_error_no_memory` - If memory allocation fails.

2.6.5.6 bfd_make_section_anyway_with_flags

Synopsis

```
asection *bfd_make_section_anyway_with_flags
(bfd *abfd, const char *name, flagword flags);
```

Description

Create a new empty section called *name* and attach it to the end of the chain of sections for *abfd*. Create a new section even if there is already a section with that name. Also set the attributes of the new section to the value *flags*.

Return NULL and set `bfd_error` on error; possible errors are:

- `bfd_error_invalid_operation` - If output has already started for *abfd*.
- `bfd_error_no_memory` - If memory allocation fails.

2.6.5.7 bfd_make_section_anyway

Synopsis

```
asection *bfd_make_section_anyway (bfd *abfd, const char *name);
```

Description

Create a new empty section called *name* and attach it to the end of the chain of sections for *abfd*. Create a new section even if there is already a section with that name.

Return NULL and set `bfd_error` on error; possible errors are:

- `bfd_error_invalid_operation` - If output has already started for *abfd*.
- `bfd_error_no_memory` - If memory allocation fails.

2.6.5.8 bfd_make_section_with_flags

Synopsis

```
asection *bfd_make_section_with_flags
(bfd *, const char *name, flagword flags);
```

Description

Like `bfd_make_section_anyway`, but return NULL (without calling `bfd_set_error()`) without changing the section chain if there is already a section named *name*. Also set the attributes of the new section to the value *flags*. If there is an error, return NULL and set `bfd_error`.

2.6.5.9 bfd_make_section

Synopsis

```
asection *bfd_make_section (bfd *, const char *name);
```

Description

Like `bfd_make_section_anyway`, but return NULL (without calling `bfd_set_error()`) without changing the section chain if there is already a section named *name*. If there is an error, return NULL and set `bfd_error`.

2.6.5.10 bfd_set_section_flags

Synopsis

```
bfd_boolean bfd_set_section_flags
(bfd *abfd, asection *sec, flagword flags);
```

Description

Set the attributes of the section *sec* in the BFD *abfd* to the value *flags*. Return TRUE on success, FALSE on error. Possible error returns are:

- `bfd_error_invalid_operation` - The section cannot have one or more of the attributes requested. For example, a .bss section in *a.out* may not have the `SEC_HAS_CONTENTS` field set.

2.6.5.11 bfd_rename_section

Synopsis

```
void bfd_rename_section
(bfd *abfd, asection *sec, const char *newname);
```

Description

Rename section *sec* in *abfd* to *newname*.

2.6.5.12 bfd_map_over_sections

Synopsis

```
void bfd_map_over_sections
(bfd *abfd,
 void (*func) (bfd *abfd, asection *sect, void *obj),
 void *obj);
```

Description

Call the provided function *func* for each section attached to the BFD *abfd*, passing *obj* as an argument. The function will be called as if by

```
func (abfd, the_section, obj);
```

This is the preferred method for iterating over sections; an alternative would be to use a loop:

```
section *p;
for (p = abfd->sections; p != NULL; p = p->next)
    func (abfd, p, ...)
```

2.6.5.13 bfd_sections_find_if

Synopsis

```
asection *bfd_sections_find_if
(bfd *abfd,
 bfd_boolean (*operation) (bfd *abfd, asection *sect, void *obj),
 void *obj);
```

Description

Call the provided function *operation* for each section attached to the BFD *abfd*, passing *obj* as an argument. The function will be called as if by

```
operation (abfd, the_section, obj);
```

It returns the first section for which *operation* returns true.

2.6.5.14 bfd_set_section_size

Synopsis

```
bfd_boolean bfd_set_section_size
(bfd *abfd, asection *sec, bfd_size_type val);
```

Description

Set *sec* to the size *val*. If the operation is ok, then **TRUE** is returned, else **FALSE**.

Possible error returns:

- **bfd_error_invalid_operation** - Writing has started to the BFD, so setting the size is invalid.

2.6.5.15 bfd_set_section_contents

Synopsis

```
bfd_boolean bfd_set_section_contents
(bfd *abfd, asection *section, const void *data,
 file_ptr offset, bfd_size_type count);
```

Description

Sets the contents of the section *section* in BFD *abfd* to the data starting in memory at *data*. The data is written to the output section starting at offset *offset* for *count* octets.

Normally TRUE is returned, else FALSE. Possible error returns are:

- `bfd_error_no_contents` - The output section does not have the `SEC_HAS_CONTENTS` attribute, so nothing can be written to it.
- and some more too

This routine is front end to the back end function `_bfd_set_section_contents`.

2.6.5.16 bfd_get_section_contents**Synopsis**

```
bfd_boolean bfd_get_section_contents
(bfd *abfd, asection *section, void *location, file_ptr offset,
 bfd_size_type count);
```

Description

Read data from *section* in BFD *abfd* into memory starting at *location*. The data is read at an offset of *offset* from the start of the input section, and is read for *count* bytes.

If the contents of a constructor with the `SEC_CONSTRUCTOR` flag set are requested or if the section does not have the `SEC_HAS_CONTENTS` flag set, then the *location* is filled with zeroes. If no errors occur, TRUE is returned, else FALSE.

2.6.5.17 bfd_malloc_and_get_section**Synopsis**

```
bfd_boolean bfd_malloc_and_get_section
(bfd *abfd, asection *section, bfd_byte **buf);
```

Description

Read all data from *section* in BFD *abfd* into a buffer, **buf*, malloc'd by this function.

2.6.5.18 bfd_copy_private_section_data**Synopsis**

```
bfd_boolean bfd_copy_private_section_data
(bfd *ibfd, asection *isec, bfd *obfd, asection *osec);
```

Description

Copy private section information from *isec* in the BFD *ibfd* to the section *osec* in the BFD *obfd*. Return TRUE on success, FALSE on error. Possible error returns are:

- `bfd_error_no_memory` - Not enough memory exists to create private data for *osec*.
- ```
#define bfd_copy_private_section_data(ibfd, isection, obfd, osection) \
 BFD_SEND (obfd, _bfd_copy_private_section_data, \
 (ibfd, isection, obfd, osection))
```

**2.6.5.19 bfd\_generic\_is\_group\_section****Synopsis**

```
bfd_boolean bfd_generic_is_group_section (bfd *, const asection *sec);
```

#### Description

Returns TRUE if *sec* is a member of a group.

#### 2.6.5.20 bfd\_generic\_discard\_group

##### Synopsis

```
bfd_boolean bfd_generic_discard_group (bfd *abfd, asection *group);
```

#### Description

Remove all members of *group* from the output.

## 2.7 Symbols

BFD tries to maintain as much symbol information as it can when it moves information from file to file. BFD passes information to applications through the `asymbol` structure. When the application requests the symbol table, BFD reads the table in the native form and translates parts of it into the internal format. To maintain more than the information passed to applications, some targets keep some information “behind the scenes” in a structure only the particular back end knows about. For example, the coff back end keeps the original symbol table structure as well as the canonical structure when a BFD is read in. On output, the coff back end can reconstruct the output symbol table so that no information is lost, even information unique to coff which BFD doesn’t know or understand. If a coff symbol table were read, but were written through an a.out back end, all the coff specific information would be lost. The symbol table of a BFD is not necessarily read in until a canonicalize request is made. Then the BFD back end fills in a table provided by the application with pointers to the canonical information. To output symbols, the application provides BFD with a table of pointers to pointers to `asymbols`. This allows applications like the linker to output a symbol as it was read, since the “behind the scenes” information will be still available.

### 2.7.1 Reading symbols

There are two stages to reading a symbol table from a BFD: allocating storage, and the actual reading process. This is an excerpt from an application which reads the symbol table:

```
long storage_needed;
asymbol **symbol_table;
long number_of_symbols;
long i;

storage_needed = bfd_get_symtab_upper_bound (abfd);

if (storage_needed < 0)
 FAIL

if (storage_needed == 0)
 return;

symbol_table = xmalloc (storage_needed);
```



```

...
number_of_symbols =
 bfd_canonicalize_symtab (abfd, symbol_table);

if (number_of_symbols < 0)
 FAIL

for (i = 0; i < number_of_symbols; i++)
 process_symbol (symbol_table[i]);

```

All storage for the symbols themselves is in an objalloc connected to the BFD; it is freed when the BFD is closed.

### 2.7.2 Writing symbols

Writing of a symbol table is automatic when a BFD open for writing is closed. The application attaches a vector of pointers to pointers to symbols to the BFD being written, and fills in the symbol count. The close and cleanup code reads through the table provided and performs all the necessary operations. The BFD output code must always be provided with an “owned” symbol: one which has come from another BFD, or one which has been created using `bfd_make_empty_symbol`. Here is an example showing the creation of a symbol table with only one element:

```

#include "bfd.h"
int main (void)
{
 bfd *abfd;
 asymbol *ptrs[2];
 asymbol *new;

 abfd = bfd_openw ("foo", "a.out-sunos-big");
 bfd_set_format (abfd, bfd_object);
 new = bfd_make_empty_symbol (abfd);
 new->name = "dummy_symbol";
 new->section = bfd_make_section_old_way (abfd, ".text");
 new->flags = BSF_GLOBAL;
 new->value = 0x12345;

 ptrs[0] = new;
 ptrs[1] = 0;

 bfd_set_symtab (abfd, ptrs, 1);
 bfd_close (abfd);
 return 0;
}

./makesym
nm foo
00012345 A dummy_symbol

```

Many formats cannot represent arbitrary symbol information; for instance, the `a.out` object format does not allow an arbitrary number of sections. A symbol pointing to a section which is not one of `.text`, `.data` or `.bss` cannot be described.

### 2.7.3 Mini Symbols

Mini symbols provide read-only access to the symbol table. They use less memory space, but require more time to access. They can be useful for tools like `nm` or `objdump`, which may have to handle symbol tables of extremely large executables.

The `bfd_read_minisymbols` function will read the symbols into memory in an internal form. It will return a `void *` pointer to a block of memory, a symbol count, and the size of each symbol. The pointer is allocated using `malloc`, and should be freed by the caller when it is no longer needed.

The function `bfd_minisymbol_to_symbol` will take a pointer to a minisymbol, and a pointer to a structure returned by `bfd_make_empty_symbol`, and return a `asymbol` structure. The return value may or may not be the same as the value from `bfd_make_empty_symbol` which was passed in.

### 2.7.4 typedef asymbol

An `asymbol` has the form:

```
typedef struct bfd_symbol
{
 /* A pointer to the BFD which owns the symbol. This information
 is necessary so that a back end can work out what additional
 information (invisible to the application writer) is carried
 with the symbol.

 This field is *almost* redundant, since you can use section->owner
 instead, except that some symbols point to the global sections
 bfd_{abs,com,und}_section. This could be fixed by making
 these globals be per-bfd (or per-target-flavor). FIXME. */
 struct bfd *the_bfd; /* Use bfd_asymbol_bfd(sym) to access this field. */

 /* The text of the symbol. The name is left alone, and not copied; the
 application may not alter it. */
 const char *name;

 /* The value of the symbol. This really should be a union of a
 numeric value with a pointer, since some flags indicate that
 a pointer to another symbol is stored here. */
 symvalue value;

 /* Attributes of a symbol. */
#define BSF_NO_FLAGS 0x00

 /* The symbol has local scope; static in C. The value
```

```

 is the offset into the section of the data. */
#define BSF_LOCAL (1 << 0)

 /* The symbol has global scope; initialized data in C. The
 value is the offset into the section of the data. */
#define BSF_GLOBAL (1 << 1)

 /* The symbol has global scope and is exported. The value is
 the offset into the section of the data. */
#define BSF_EXPORT BSF_GLOBAL /* No real difference. */

 /* A normal C symbol would be one of:
 BSF_LOCAL, BSF_COMMON, BSF_UNDEFINED or
 BSF_GLOBAL. */

 /* The symbol is a debugging record. The value has an arbitrary
 meaning, unless BSF_DEBUGGING_RELOC is also set. */
#define BSF_DEBUGGING (1 << 2)

 /* The symbol denotes a function entry point. Used in ELF,
 perhaps others someday. */
#define BSF_FUNCTION (1 << 3)

 /* Used by the linker. */
#define BSF_KEEP (1 << 5)
#define BSF_KEEP_G (1 << 6)

 /* A weak global symbol, overridable without warnings by
 a regular global symbol of the same name. */
#define BSF_WEAK (1 << 7)

 /* This symbol was created to point to a section, e.g. ELF's
 STT_SECTION symbols. */
#define BSF_SECTION_SYM (1 << 8)

 /* The symbol used to be a common symbol, but now it is
 allocated. */
#define BSF_OLD_COMMON (1 << 9)

 /* In some files the type of a symbol sometimes alters its
 location in an output file - ie in coff a ISFCN symbol
 which is also C_EXT symbol appears where it was
 declared and not at the end of a section. This bit is set
 by the target BFD part to convey this information. */
#define BSF_NOT_AT_END (1 << 10)

 /* Signal that the symbol is the label of constructor section. */

```

```
#define BSF_CONSTRUCTOR (1 << 11)

/* Signal that the symbol is a warning symbol. The name is a
 warning. The name of the next symbol is the one to warn about;
 if a reference is made to a symbol with the same name as the next
 symbol, a warning is issued by the linker. */
#define BSF_WARNING (1 << 12)

/* Signal that the symbol is indirect. This symbol is an indirect
 pointer to the symbol with the same name as the next symbol. */
#define BSF_INDIRECT (1 << 13)

/* BSF_FILE marks symbols that contain a file name. This is used
 for ELF STT_FILE symbols. */
#define BSF_FILE (1 << 14)

/* Symbol is from dynamic linking information. */
#define BSF_DYNAMIC (1 << 15)

/* The symbol denotes a data object. Used in ELF, and perhaps
 others someday. */
#define BSF_OBJECT (1 << 16)

/* This symbol is a debugging symbol. The value is the offset
 into the section of the data. BSF_DEBUGGING should be set
 as well. */
#define BSF_DEBUGGING_RELOC (1 << 17)

/* This symbol is thread local. Used in ELF. */
#define BSF_THREAD_LOCAL (1 << 18)

/* This symbol represents a complex relocation expression,
 with the expression tree serialized in the symbol name. */
#define BSF_RELC (1 << 19)

/* This symbol represents a signed complex relocation expression,
 with the expression tree serialized in the symbol name. */
#define BSF_SRELC (1 << 20)

/* This symbol was created by bfd_get_synthetic_symtab. */
#define BSF_SYNTHETIC (1 << 21)

/* This symbol is an indirect code object. Unrelated to BSF_INDIRECT.
 The dynamic linker will compute the value of this symbol by
 calling the function that it points to. BSF_FUNCTION must
 also be also set. */
#define BSF_GNU_INDIRECT_FUNCTION (1 << 22)
```

```

 /* This symbol is a globally unique data object. The dynamic linker
 will make sure that in the entire process there is just one symbol
 with this name and type in use. BSF_OBJECT must also be set. */
#define BSF_GNU_UNIQUE (1 << 23)

 flagword flags;

 /* A pointer to the section to which this symbol is
 relative. This will always be non NULL, there are special
 sections for undefined and absolute symbols. */
 struct bfd_section *section;

 /* Back end special data. */
 union
 {
 void *p;
 bfd_vma i;
 }
 udata;
}
asymbol;

```

## 2.7.5 Symbol handling functions

### 2.7.5.1 bfd\_get\_symtab\_upper\_bound

#### Description

Return the number of bytes required to store a vector of pointers to `asymbols` for all the symbols in the BFD *abfd*, including a terminal NULL pointer. If there are no symbols in the BFD, then return 0. If an error occurs, return -1.

```

#define bfd_get_symtab_upper_bound(abfd) \
 BFD_SEND (abfd, _bfd_get_symtab_upper_bound, (abfd))

```

### 2.7.5.2 bfd\_is\_local\_label

#### Synopsis

```

bfd_boolean bfd_is_local_label (bfd *abfd, asymbol *sym);

```

#### Description

Return TRUE if the given symbol *sym* in the BFD *abfd* is a compiler generated local label, else return FALSE.

### 2.7.5.3 bfd\_is\_local\_label\_name

#### Synopsis

```

bfd_boolean bfd_is_local_label_name (bfd *abfd, const char *name);

```

#### Description

Return TRUE if a symbol with the name *name* in the BFD *abfd* is a compiler generated

local label, else return FALSE. This just checks whether the name has the form of a local label.

```
#define bfd_is_local_label_name(abfd, name) \
 BFD_SEND (abfd, _bfd_is_local_label_name, (abfd, name))
```

#### 2.7.5.4 bfd\_is\_target\_special\_symbol

##### Synopsis

```
bfd_boolean bfd_is_target_special_symbol (bfd *abfd, asymbol *sym);
```

##### Description

Return TRUE iff a symbol *sym* in the BFD *abfd* is something special to the particular target represented by the BFD. Such symbols should normally not be mentioned to the user.

```
#define bfd_is_target_special_symbol(abfd, sym) \
 BFD_SEND (abfd, _bfd_is_target_special_symbol, (abfd, sym))
```

#### 2.7.5.5 bfd\_canonicalize\_symtab

##### Description

Read the symbols from the BFD *abfd*, and fills in the vector *location* with pointers to the symbols and a trailing NULL. Return the actual number of symbol pointers, not including the NULL.

```
#define bfd_canonicalize_symtab(abfd, location) \
 BFD_SEND (abfd, _bfd_canonicalize_symtab, (abfd, location))
```

#### 2.7.5.6 bfd\_set\_symtab

##### Synopsis

```
bfd_boolean bfd_set_symtab
(bfd *abfd, asymbol **location, unsigned int count);
```

##### Description

Arrange that when the output BFD *abfd* is closed, the table *location* of *count* pointers to symbols will be written.

#### 2.7.5.7 bfd\_print\_symbol\_vandf

##### Synopsis

```
void bfd_print_symbol_vandf (bfd *abfd, void *file, asymbol *symbol);
```

##### Description

Print the value and flags of the *symbol* supplied to the stream *file*.

#### 2.7.5.8 bfd\_make\_empty\_symbol

##### Description

Create a new *asymbol* structure for the BFD *abfd* and return a pointer to it.

This routine is necessary because each back end has private information surrounding the `asymbol`. Building your own `asymbol` and pointing to it will not create the private information, and will cause problems later on.

```
#define bfd_make_empty_symbol(abfd) \
 BFD_SEND (abfd, _bfd_make_empty_symbol, (abfd))
```

### 2.7.5.9 `_bfd_generic_make_empty_symbol`

#### Synopsis

```
asymbol *_bfd_generic_make_empty_symbol (bfd *);
```

#### Description

Create a new `asymbol` structure for the BFD *abfd* and return a pointer to it. Used by core file routines, binary back-end and anywhere else where no private info is needed.

### 2.7.5.10 `bfd_make_debug_symbol`

#### Description

Create a new `asymbol` structure for the BFD *abfd*, to be used as a debugging symbol. Further details of its use have yet to be worked out.

```
#define bfd_make_debug_symbol(abfd,ptr,size) \
 BFD_SEND (abfd, _bfd_make_debug_symbol, (abfd, ptr, size))
```

### 2.7.5.11 `bfd_decode_symclass`

#### Description

Return a character corresponding to the symbol class of *symbol*, or '?' for an unknown class.

#### Synopsis

```
int bfd_decode_symclass (asymbol *symbol);
```

### 2.7.5.12 `bfd_is_undefined_symclass`

#### Description

Returns non-zero if the class symbol returned by `bfd_decode_symclass` represents an undefined symbol. Returns zero otherwise.

#### Synopsis

```
bfd_boolean bfd_is_undefined_symclass (int symclass);
```

### 2.7.5.13 `bfd_symbol_info`

#### Description

Fill in the basic info about symbol that nm needs. Additional info may be added by the back-ends after calling this function.

#### Synopsis

```
void bfd_symbol_info (asymbol *symbol, symbol_info *ret);
```

### 2.7.5.14 bfd\_copy\_private\_symbol\_data

#### Synopsis

```
bfd_boolean bfd_copy_private_symbol_data
 (bfd *ibfd, asymbol *isym, bfd *obfd, asymbol *osym);
```

#### Description

Copy private symbol information from *isym* in the BFD *ibfd* to the symbol *osym* in the BFD *obfd*. Return TRUE on success, FALSE on error. Possible error returns are:

- `bfd_error_no_memory` - Not enough memory exists to create private data for *osym*.

```
#define bfd_copy_private_symbol_data(ibfd, isymbol, obfd, osymbol) \
 BFD_SEND (obfd, _bfd_copy_private_symbol_data, \
 (ibfd, isymbol, obfd, osymbol))
```

## 2.8 Archives

### Description

An archive (or library) is just another BFD. It has a symbol table, although there's not much a user program will do with it.

The big difference between an archive BFD and an ordinary BFD is that the archive doesn't have sections. Instead it has a chain of BFDs that are considered its contents. These BFDs can be manipulated like any other. The BFDs contained in an archive opened for reading will all be opened for reading. You may put either input or output BFDs into an archive opened for output; they will be handled correctly when the archive is closed.

Use `bfd_openr_next_archived_file` to step through the contents of an archive opened for input. You don't have to read the entire archive if you don't want to! Read it until you find what you want.

Archive contents of output BFDs are chained through the `next` pointer in a BFD. The first one is findable through the `archive_head` slot of the archive. Set it with `bfd_set_archive_head` (q.v.). A given BFD may be in only one open output archive at a time.

As expected, the BFD archive code is more general than the archive code of any given environment. BFD archives may contain files of different formats (e.g., a.out and coff) and even different architectures. You may even place archives recursively into archives!

This can cause unexpected confusion, since some archive formats are more expressive than others. For instance, Intel COFF archives can preserve long filenames; SunOS a.out archives cannot. If you move a file from the first to the second format and back again, the filename may be truncated. Likewise, different a.out environments have different conventions as to how they truncate filenames, whether they preserve directory names in filenames, etc. When interoperating with native tools, be sure your files are homogeneous.

Beware: most of these formats do not react well to the presence of spaces in filenames. We do the best we can, but can't always handle this case due to restrictions in the format of archives. Many Unix utilities are braindead in regards to spaces and such in filenames anyway, so this shouldn't be much of a restriction.

Archives are supported in BFD in `archive.c`.

### 2.8.1 Archive functions



### 2.8.1.1 bfd\_get\_next\_mapent

#### Synopsis

```
symindex bfd_get_next_mapent
(bfd *abfd, symindex previous, carsym **sym);
```

#### Description

Step through archive *abfd*'s symbol table (if it has one). Successively update *sym* with the next symbol's information, returning that symbol's (internal) index into the symbol table.

Supply BFD\_NO\_MORE\_SYMBOLS as the *previous* entry to get the first one; returns BFD\_NO\_MORE\_SYMBOLS when you've already got the last one.

A *carsym* is a canonical archive symbol. The only user-visible element is its name, a null-terminated string.

### 2.8.1.2 bfd\_set\_archive\_head

#### Synopsis

```
bfd_boolean bfd_set_archive_head (bfd *output, bfd *new_head);
```

#### Description

Set the head of the chain of BFDs contained in the archive *output* to *new\_head*.

### 2.8.1.3 bfd\_openr\_next\_archived\_file

#### Synopsis

```
bfd *bfd_openr_next_archived_file (bfd *archive, bfd *previous);
```

#### Description

Provided a BFD, *archive*, containing an archive and NULL, open an input BFD on the first contained element and returns that. Subsequent calls should pass the archive and the previous return value to return a created BFD to the next contained element. NULL is returned when there are no more.

## 2.9 File formats

A format is a BFD concept of high level file contents type. The formats supported by BFD are:

- **bfd\_object**

The BFD may contain data, symbols, relocations and debug info.

- **bfd\_archive**

The BFD contains other BFDs and an optional index.

- **bfd\_core**

The BFD contains the result of an executable core dump.

### 2.9.1 File format functions

### 2.9.1.1 bfd\_check\_format

#### Synopsis

```
bfd_boolean bfd_check_format (bfd *abfd, bfd_format format);
```

#### Description

Verify if the file attached to the BFD *abfd* is compatible with the format *format* (i.e., one of `bfd_object`, `bfd_archive` or `bfd_core`).

If the BFD has been set to a specific target before the call, only the named target and format combination is checked. If the target has not been set, or has been set to `default`, then all the known target backends is interrogated to determine a match. If the default target matches, it is used. If not, exactly one target must recognize the file, or an error results.

The function returns `TRUE` on success, otherwise `FALSE` with one of the following error codes:

- `bfd_error_invalid_operation` - if *format* is not one of `bfd_object`, `bfd_archive` or `bfd_core`.
- `bfd_error_system_call` - if an error occurred during a read - even some file mismatches can cause `bfd_error_system_calls`.
- `file_not_recognised` - none of the backends recognised the file format.
- `bfd_error_file_ambiguously_recognized` - more than one backend recognised the file format.

### 2.9.1.2 bfd\_check\_format\_matches

#### Synopsis

```
bfd_boolean bfd_check_format_matches
(bfd *abfd, bfd_format format, char ***matching);
```

#### Description

Like `bfd_check_format`, except when it returns `FALSE` with `bfd_errno` set to `bfd_error_file_ambiguously_recognized`. In that case, if *matching* is not `NULL`, it will be filled in with a `NULL`-terminated list of the names of the formats that matched, allocated with `malloc`. Then the user may choose a format and try again.

When done with the list that *matching* points to, the caller should free it.

### 2.9.1.3 bfd\_set\_format

#### Synopsis

```
bfd_boolean bfd_set_format (bfd *abfd, bfd_format format);
```

#### Description

This function sets the file format of the BFD *abfd* to the format *format*. If the target set in the BFD does not support the format requested, the format is invalid, or the BFD is not open for writing, then an error occurs.

### 2.9.1.4 bfd\_format\_string

#### Synopsis

```
const char *bfd_format_string (bfd_format format);
```

**Description**

Return a pointer to a const string `invalid`, `object`, `archive`, `core`, or `unknown`, depending upon the value of *format*.

**2.10 Relocations**

BFD maintains relocations in much the same way it maintains symbols: they are left alone until required, then read in en-masse and translated into an internal form. A common routine `bfd_perform_relocation` acts upon the canonical form to do the fixup.

Relocations are maintained on a per section basis, while symbols are maintained on a per BFD basis.

All that a back end has to do to fit the BFD interface is to create a `struct reloc_cache_entry` for each relocation in a particular section, and fill in the right bits of the structures.

**2.10.1 typedef arelent**

This is the structure of a relocation entry:

```
typedef enum bfd_reloc_status
{
 /* No errors detected. */
 bfd_reloc_ok,

 /* The relocation was performed, but there was an overflow. */
 bfd_reloc_overflow,

 /* The address to relocate was not within the section supplied. */
 bfd_reloc_outofrange,

 /* Used by special functions. */
 bfd_reloc_continue,

 /* Unsupported relocation size requested. */
 bfd_reloc_notsupported,

 /* Unused. */
 bfd_reloc_other,

 /* The symbol to relocate against was undefined. */
 bfd_reloc_undefined,

 /* The relocation was performed, but may not be ok - presently
 generated only when linking i960 coff files with i960 b.out
 symbols. If this type is returned, the error_message argument
 to bfd_perform_relocation will be set. */
 bfd_reloc_dangerous
}
```

```

bfd_reloc_status_type;

typedef struct reloc_cache_entry
{
 /* A pointer into the canonical table of pointers. */
 struct bfd_symbol **sym_ptr_ptr;

 /* offset in section. */
 bfd_size_type address;

 /* addend for relocation value. */
 bfd_vma addend;

 /* Pointer to how to perform the required relocation. */
 reloc_howto_type *howto;
}
arelent;

```

### Description

Here is a description of each of the fields within an `arelent`:

- `sym_ptr_ptr`

The symbol table pointer points to a pointer to the symbol associated with the relocation request. It is the pointer into the table returned by the back end's `canonicalize_symtab` action. See [Section 2.7 \[Symbols\]](#), page 40. The symbol is referenced through a pointer to a pointer so that tools like the linker can fix up all the symbols of the same name by modifying only one pointer. The relocation routine looks in the symbol and uses the base of the section the symbol is attached to and the value of the symbol as the initial relocation offset. If the symbol pointer is zero, then the section provided is looked up.

- `address`

The `address` field gives the offset in bytes from the base of the section data which owns the relocation record to the first byte of relocatable information. The actual data relocated will be relative to this point; for example, a relocation type which modifies the bottom two bytes of a four byte word would not touch the first byte pointed to in a big endian world.

- `addend`

The `addend` is a value provided by the back end to be added (!) to the relocation offset. Its interpretation is dependent upon the howto. For example, on the 68k the code:

```

char foo[];
main()
{
 return foo[0x12345678];
}

```

Could be compiled into:

```

linkw fp,#-4
moveb @#12345678,d0
extbl d0
unlk fp
rts

```

This could create a reloc pointing to `foo`, but leave the offset in the data, something like:

```

RELOCATION RECORDS FOR [.text]:
offset type value
00000006 32 _foo

00000000 4e56 fffc ; linkw fp,#-4
00000004 1039 1234 5678 ; moveb @#12345678,d0
0000000a 49c0 ; extbl d0
0000000c 4e5e ; unlk fp
0000000e 4e75 ; rts

```

Using `coff` and an 88k, some instructions don't have enough space in them to represent the full address range, and pointers have to be loaded in two parts. So you'd get something like:

```

or.u r13,r0,hi16(_foo+0x12345678)
ld.b r2,r13,lo16(_foo+0x12345678)
jmp r1

```

This should create two relocs, both pointing to `_foo`, and with `0x12340000` in their addend field. The data would consist of:

```

RELOCATION RECORDS FOR [.text]:
offset type value
00000002 HVRT16 _foo+0x12340000
00000006 LVRT16 _foo+0x12340000

00000000 5da05678 ; or.u r13,r0,0x5678
00000004 1c4d5678 ; ld.b r2,r13,0x5678
00000008 f400c001 ; jmp r1

```

The relocation routine digs out the value from the data, adds it to the addend to get the original offset, and then adds the value of `_foo`. Note that all 32 bits have to be kept around somewhere, to cope with carry from bit 15 to bit 16.

One further example is the `sparc` and the `a.out` format. The `sparc` has a similar problem to the 88k, in that some instructions don't have room for an entire offset, but on the `sparc` the parts are created in odd sized lumps. The designers of the `a.out` format chose to not use the data within the section for storing part of the offset; all the offset is kept within the reloc. Anything in the data should be ignored.

```

save %sp,-112,%sp
sethi %hi(_foo+0x12345678),%g2
ldsb [%g2+%lo(_foo+0x12345678)],%i0
ret
restore

```

Both relocs contain a pointer to `foo`, and the offsets contain junk.

```

RELOCATION RECORDS FOR [.text]:
offset type value
00000004 HI22 _foo+0x12345678
00000008 L010 _foo+0x12345678

00000000 9de3bf90 ; save %sp,-112,%sp
00000004 05000000 ; sethi %hi(_foo+0),%g2
00000008 f048a000 ; ldsb [%g2+%lo(_foo+0)],%i0
0000000c 81c7e008 ; ret
00000010 81e80000 ; restore

```

- **howto**

The **howto** field can be imagined as a relocation instruction. It is a pointer to a structure which contains information on what to do with all of the other information in the reloc record and data section. A back end would normally have a relocation instruction set and turn relocations into pointers to the correct structure on input - but it would be possible to create each **howto** field on demand.

### 2.10.1.1 enum complain\_overflow

Indicates what sort of overflow checking should be done when performing a relocation.

```

enum complain_overflow
{
 /* Do not complain on overflow. */
 complain_overflow_dont,

 /* Complain if the value overflows when considered as a signed
 number one bit larger than the field. ie. A bitfield of N bits
 is allowed to represent -2**n to 2**n-1. */
 complain_overflow_bitfield,

 /* Complain if the value overflows when considered as a signed
 number. */
 complain_overflow_signed,

 /* Complain if the value overflows when considered as an
 unsigned number. */
 complain_overflow_unsigned
};

```

### 2.10.1.2 reloc\_howto\_type

The **reloc\_howto\_type** is a structure which contains all the information that libbfd needs to know to tie up a back end's data.

```

struct bfd_symbol; /* Forward declaration. */

struct reloc_howto_struct

```

```

{
 /* The type field has mainly a documentary use - the back end can
 do what it wants with it, though normally the back end's
 external idea of what a reloc number is stored
 in this field. For example, a PC relative word relocation
 in a coff environment has the type 023 - because that's
 what the outside world calls a R_PCRWORD reloc. */
 unsigned int type;

 /* The value the final relocation is shifted right by. This drops
 unwanted data from the relocation. */
 unsigned int rightshift;

 /* The size of the item to be relocated. This is *not* a
 power-of-two measure. To get the number of bytes operated
 on by a type of relocation, use bfd_get_reloc_size. */
 int size;

 /* The number of bits in the item to be relocated. This is used
 when doing overflow checking. */
 unsigned int bitsize;

 /* The relocation is relative to the field being relocated. */
 bfd_boolean pc_relative;

 /* The bit position of the reloc value in the destination.
 The relocated value is left shifted by this amount. */
 unsigned int bitpos;

 /* What type of overflow error should be checked for when
 relocating. */
 enum complain_overflow complain_on_overflow;

 /* If this field is non null, then the supplied function is
 called rather than the normal function. This allows really
 strange relocation methods to be accommodated (e.g., i960 callj
 instructions). */
 bfd_reloc_status_type (*special_function)
 (bfd *, arelent *, struct bfd_symbol *, void *, asection *,
 bfd *, char **);

 /* The textual name of the relocation type. */
 char *name;

 /* Some formats record a relocation addend in the section contents
 rather than with the relocation. For ELF formats this is the
 distinction between USE_REL and USE_RELA (though the code checks

```

```

 for USE_REL == 1/0). The value of this field is TRUE if the
 addend is recorded with the section contents; when performing a
 partial link (ld -r) the section contents (the data) will be
 modified. The value of this field is FALSE if addends are
 recorded with the relocation (in arelent.addend); when performing
 a partial link the relocation will be modified.
 All relocations for all ELF USE_RELA targets should set this field
 to FALSE (values of TRUE should be looked on with suspicion).
 However, the converse is not true: not all relocations of all ELF
 USE_REL targets set this field to TRUE. Why this is so is peculiar
 to each particular target. For relocs that aren't used in partial
 links (e.g. GOT stuff) it doesn't matter what this is set to. */
bfd_boolean partial_inplace;

/* src_mask selects the part of the instruction (or data) to be used
 in the relocation sum. If the target relocations don't have an
 addend in the reloc, eg. ELF USE_REL, src_mask will normally equal
 dst_mask to extract the addend from the section contents. If
 relocations do have an addend in the reloc, eg. ELF USE_RELA, this
 field should be zero. Non-zero values for ELF USE_RELA targets are
 bogus as in those cases the value in the dst_mask part of the
 section contents should be treated as garbage. */
bfd_vma src_mask;

/* dst_mask selects which parts of the instruction (or data) are
 replaced with a relocated value. */
bfd_vma dst_mask;

/* When some formats create PC relative instructions, they leave
 the value of the pc of the place being relocated in the offset
 slot of the instruction, so that a PC relative relocation can
 be made just by adding in an ordinary offset (e.g., sun3 a.out).
 Some formats leave the displacement part of an instruction
 empty (e.g., m88k bcs); this flag signals the fact. */
bfd_boolean pcrel_offset;
};

```

### 2.10.1.3 The HOWTO Macro

#### Description

The HOWTO define is horrible and will go away.

```

#define HOWTO(C, R, S, B, P, BI, O, SF, NAME, INPLACE, MASKSRC, MASKDST, PC) \
{ (unsigned) C, R, S, B, P, BI, O, SF, NAME, INPLACE, MASKSRC, MASKDST, PC }

```

#### Description

And will be replaced with the totally magic way. But for the moment, we are compatible, so do it this way.



```
#define NEWHOWTO(FUNCTION, NAME, SIZE, REL, IN) \
 HOWTO (0, 0, SIZE, 0, REL, 0, complain_overflow_dont, FUNCTION, \
 NAME, FALSE, 0, 0, IN)
```

**Description**

This is used to fill in an empty howto entry in an array.

```
#define EMPTY_HOWTO(C) \
 HOWTO ((C), 0, 0, 0, FALSE, 0, complain_overflow_dont, NULL, \
 NULL, FALSE, 0, 0, FALSE)
```

**Description**

Helper routine to turn a symbol into a relocation value.

```
#define HOWTO_PREPARE(relocation, symbol) \
{ \
 if (symbol != NULL) \
 { \
 if (bfd_is_com_section (symbol->section)) \
 { \
 relocation = 0; \
 } \
 else \
 { \
 relocation = symbol->value; \
 } \
 } \
}
```

**2.10.1.4 bfd\_get\_reloc\_size****Synopsis**

```
unsigned int bfd_get_reloc_size (reloc_howto_type *);
```

**Description**

For a `reloc_howto_type` that operates on a fixed number of bytes, this returns the number of bytes operated on.

**2.10.1.5 arelent\_chain****Description**

How relocs are tied together in an asection:

```
typedef struct relent_chain
{
 arelent relent;
 struct relent_chain *next;
}
arentent_chain;
```

### 2.10.1.6 bfd\_check\_overflow

#### Synopsis

```
bfd_reloc_status_type bfd_check_overflow
(enum complain_overflow how,
 unsigned int bitsize,
 unsigned int rightshift,
 unsigned int addrsz,
 bfd_vma relocation);
```

#### Description

Perform overflow checking on *relocation* which has *bitsize* significant bits and will be shifted right by *rightshift* bits, on a machine with addresses containing *addrsz* significant bits. The result is either of `bfd_reloc_ok` or `bfd_reloc_overflow`.

### 2.10.1.7 bfd\_perform\_relocation

#### Synopsis

```
bfd_reloc_status_type bfd_perform_relocation
(bfd *abfd,
 arelent *reloc_entry,
 void *data,
 asection *input_section,
 bfd *output_bfd,
 char **error_message);
```

#### Description

If *output\_bfd* is supplied to this function, the generated image will be relocatable; the relocations are copied to the output file after they have been changed to reflect the new state of the world. There are two ways of reflecting the results of partial linkage in an output file: by modifying the output data in place, and by modifying the relocation record. Some native formats (e.g., basic a.out and basic coff) have no way of specifying an addend in the relocation type, so the addend has to go in the output data. This is no big deal since in these formats the output data slot will always be big enough for the addend. Complex reloc types with addends were invented to solve just this problem. The *error\_message* argument is set to an error message if this return `bfd_reloc_dangerous`.

### 2.10.1.8 bfd\_install\_relocation

#### Synopsis

```
bfd_reloc_status_type bfd_install_relocation
(bfd *abfd,
 arelent *reloc_entry,
 void *data, bfd_vma data_start,
 asection *input_section,
 char **error_message);
```

#### Description

This looks remarkably like `bfd_perform_relocation`, except it does not expect that the section contents have been filled in. I.e., it's suitable for use when creating, rather than applying a relocation.

For now, this function should be considered reserved for the assembler.

## 2.10.2 The howto manager

When an application wants to create a relocation, but doesn't know what the target machine might call it, it can find out by using this bit of code.

### 2.10.2.1 bfd\_reloc\_code\_type

#### Description

The insides of a reloc code. The idea is that, eventually, there will be one enumerator for every type of relocation we ever do. Pass one of these values to `bfd_reloc_type_lookup`, and it'll return a howto pointer.

This does mean that the application must determine the correct enumerator value; you can't get a howto pointer from a random set of attributes.

Here are the possible values for `enum bfd_reloc_code_real`:

`BFD_RELOC_64`

`BFD_RELOC_32`

`BFD_RELOC_26`

`BFD_RELOC_24`

`BFD_RELOC_16`

`BFD_RELOC_14`

`BFD_RELOC_8`

Basic absolute relocations of N bits.

`BFD_RELOC_64_PCREL`

`BFD_RELOC_32_PCREL`

`BFD_RELOC_24_PCREL`

`BFD_RELOC_16_PCREL`

`BFD_RELOC_12_PCREL`

`BFD_RELOC_8_PCREL`

PC-relative relocations. Sometimes these are relative to the address of the relocation itself; sometimes they are relative to the start of the section containing the relocation. It depends on the specific target.

The 24-bit relocation is used in some Intel 960 configurations.

`BFD_RELOC_32_SECREL`

Section relative relocations. Some targets need this for DWARF2.

`BFD_RELOC_32_GOT_PCREL`

`BFD_RELOC_16_GOT_PCREL`

`BFD_RELOC_8_GOT_PCREL`

`BFD_RELOC_32_GOTOFF`

`BFD_RELOC_16_GOTOFF`

`BFD_RELOC_LO16_GOTOFF`

`BFD_RELOC_HI16_GOTOFF`

`BFD_RELOC_HI16_S_GOTOFF`

`BFD_RELOC_8_GOTOFF`

```

BFD_RELOC_64_PLT_PCREL
BFD_RELOC_32_PLT_PCREL
BFD_RELOC_24_PLT_PCREL
BFD_RELOC_16_PLT_PCREL
BFD_RELOC_8_PLT_PCREL
BFD_RELOC_64_PLTOFF
BFD_RELOC_32_PLTOFF
BFD_RELOC_16_PLTOFF
BFD_RELOC_L016_PLTOFF
BFD_RELOC_HI16_PLTOFF
BFD_RELOC_HI16_S_PLTOFF
BFD_RELOC_8_PLTOFF

```

For ELF.

```

BFD_RELOC_68K_GLOB_DAT
BFD_RELOC_68K_JMP_SLOT
BFD_RELOC_68K_RELATIVE
BFD_RELOC_68K_TLS_GD32
BFD_RELOC_68K_TLS_GD16
BFD_RELOC_68K_TLS_GD8
BFD_RELOC_68K_TLS_LDM32
BFD_RELOC_68K_TLS_LDM16
BFD_RELOC_68K_TLS_LDM8
BFD_RELOC_68K_TLS_LD032
BFD_RELOC_68K_TLS_LD016
BFD_RELOC_68K_TLS_LD08
BFD_RELOC_68K_TLS_IE32
BFD_RELOC_68K_TLS_IE16
BFD_RELOC_68K_TLS_IE8
BFD_RELOC_68K_TLS_LE32
BFD_RELOC_68K_TLS_LE16
BFD_RELOC_68K_TLS_LE8

```

Relocations used by 68K ELF.

```

BFD_RELOC_32_BASEREL
BFD_RELOC_16_BASEREL
BFD_RELOC_L016_BASEREL
BFD_RELOC_HI16_BASEREL
BFD_RELOC_HI16_S_BASEREL
BFD_RELOC_8_BASEREL
BFD_RELOC_RVA

```

Linkage-table relative.

```

BFD_RELOC_8_FFnn

```

Absolute 8-bit relocation, but used to form an address like 0xFFnn.

```

BFD_RELOC_32_PCREL_S2
BFD_RELOC_16_PCREL_S2

```

**BFD\_RELOC\_23\_PCREL\_S2**

These PC-relative relocations are stored as word displacements – i.e., byte displacements shifted right two bits. The 30-bit word displacement (<<32\_PCREL\_S2>> – 32 bits, shifted 2) is used on the SPARC. (SPARC tools generally refer to this as <<WDISP30>>.) The signed 16-bit displacement is used on the MIPS, and the 23-bit displacement is used on the Alpha.

**BFD\_RELOC\_HI22****BFD\_RELOC\_LO10**

High 22 bits and low 10 bits of 32-bit value, placed into lower bits of the target word. These are used on the SPARC.

**BFD\_RELOC\_GPREL16****BFD\_RELOC\_GPREL32**

For systems that allocate a Global Pointer register, these are displacements off that register. These relocation types are handled specially, because the value the register will have is decided relatively late.

**BFD\_RELOC\_I960\_CALLJ**

Reloc types used for i960/b.out.

**BFD\_RELOC\_NONE****BFD\_RELOC\_SPARC\_WDISP22****BFD\_RELOC\_SPARC22****BFD\_RELOC\_SPARC13****BFD\_RELOC\_SPARC\_GOT10****BFD\_RELOC\_SPARC\_GOT13****BFD\_RELOC\_SPARC\_GOT22****BFD\_RELOC\_SPARC\_PC10****BFD\_RELOC\_SPARC\_PC22****BFD\_RELOC\_SPARC\_WPLT30****BFD\_RELOC\_SPARC\_COPY****BFD\_RELOC\_SPARC\_GLOB\_DAT****BFD\_RELOC\_SPARC\_JMP\_SLOT****BFD\_RELOC\_SPARC\_RELATIVE****BFD\_RELOC\_SPARC\_UA16****BFD\_RELOC\_SPARC\_UA32****BFD\_RELOC\_SPARC\_UA64****BFD\_RELOC\_SPARC\_GOTDATA\_HIX22****BFD\_RELOC\_SPARC\_GOTDATA\_LOX10****BFD\_RELOC\_SPARC\_GOTDATA\_OP\_HIX22****BFD\_RELOC\_SPARC\_GOTDATA\_OP\_LOX10****BFD\_RELOC\_SPARC\_GOTDATA\_OP****BFD\_RELOC\_SPARC\_JMP\_IREL****BFD\_RELOC\_SPARC\_IRELATIVE**

SPARC ELF relocations. There is probably some overlap with other relocation types already defined.

BFD\_RELOC\_SPARC\_BASE13

BFD\_RELOC\_SPARC\_BASE22

I think these are specific to SPARC a.out (e.g., Sun 4).

BFD\_RELOC\_SPARC\_64

BFD\_RELOC\_SPARC\_10

BFD\_RELOC\_SPARC\_11

BFD\_RELOC\_SPARC\_OL010

BFD\_RELOC\_SPARC\_HH22

BFD\_RELOC\_SPARC\_HM10

BFD\_RELOC\_SPARC\_LM22

BFD\_RELOC\_SPARC\_PC\_HH22

BFD\_RELOC\_SPARC\_PC\_HM10

BFD\_RELOC\_SPARC\_PC\_LM22

BFD\_RELOC\_SPARC\_WDISP16

BFD\_RELOC\_SPARC\_WDISP19

BFD\_RELOC\_SPARC\_7

BFD\_RELOC\_SPARC\_6

BFD\_RELOC\_SPARC\_5

BFD\_RELOC\_SPARC\_DISP64

BFD\_RELOC\_SPARC\_PLT32

BFD\_RELOC\_SPARC\_PLT64

BFD\_RELOC\_SPARC\_HIX22

BFD\_RELOC\_SPARC\_LOX10

BFD\_RELOC\_SPARC\_H44

BFD\_RELOC\_SPARC\_M44

BFD\_RELOC\_SPARC\_L44

BFD\_RELOC\_SPARC\_REGISTER

SPARC64 relocations

BFD\_RELOC\_SPARC\_REV32

SPARC little endian relocation

BFD\_RELOC\_SPARC\_TLS\_GD\_HI22

BFD\_RELOC\_SPARC\_TLS\_GD\_L010

BFD\_RELOC\_SPARC\_TLS\_GD\_ADD

BFD\_RELOC\_SPARC\_TLS\_GD\_CALL

BFD\_RELOC\_SPARC\_TLS\_LDM\_HI22

BFD\_RELOC\_SPARC\_TLS\_LDM\_L010

BFD\_RELOC\_SPARC\_TLS\_LDM\_ADD

BFD\_RELOC\_SPARC\_TLS\_LDM\_CALL

BFD\_RELOC\_SPARC\_TLS\_LDO\_HIX22

BFD\_RELOC\_SPARC\_TLS\_LDO\_LOX10

BFD\_RELOC\_SPARC\_TLS\_LDO\_ADD

BFD\_RELOC\_SPARC\_TLS\_IE\_HI22

BFD\_RELOC\_SPARC\_TLS\_IE\_L010

BFD\_RELOC\_SPARC\_TLS\_IE\_LD

BFD\_RELOC\_SPARC\_TLS\_IE\_LDX

BFD\_RELOC\_SPARC\_TLS\_IE\_ADD  
 BFD\_RELOC\_SPARC\_TLS\_LE\_HIX22  
 BFD\_RELOC\_SPARC\_TLS\_LE\_LOX10  
 BFD\_RELOC\_SPARC\_TLS\_DTPMOD32  
 BFD\_RELOC\_SPARC\_TLS\_DTPMOD64  
 BFD\_RELOC\_SPARC\_TLS\_DTPOFF32  
 BFD\_RELOC\_SPARC\_TLS\_DTPOFF64  
 BFD\_RELOC\_SPARC\_TLS\_TPOFF32  
 BFD\_RELOC\_SPARC\_TLS\_TPOFF64  
 SPARC TLS relocations

BFD\_RELOC\_SPU\_IMM7  
 BFD\_RELOC\_SPU\_IMM8  
 BFD\_RELOC\_SPU\_IMM10  
 BFD\_RELOC\_SPU\_IMM10W  
 BFD\_RELOC\_SPU\_IMM16  
 BFD\_RELOC\_SPU\_IMM16W  
 BFD\_RELOC\_SPU\_IMM18  
 BFD\_RELOC\_SPU\_PCREL9a  
 BFD\_RELOC\_SPU\_PCREL9b  
 BFD\_RELOC\_SPU\_PCREL16  
 BFD\_RELOC\_SPU\_LO16  
 BFD\_RELOC\_SPU\_HI16  
 BFD\_RELOC\_SPU\_PPU32  
 BFD\_RELOC\_SPU\_PPU64  
 BFD\_RELOC\_SPU\_ADD\_PIC  
 SPU Relocations.

BFD\_RELOC\_ALPHA\_GPDISP\_HI16

Alpha ECOFF and ELF relocations. Some of these treat the symbol or "addend" in some special way. For GPDISP\_HI16 ("gpdisp") relocations, the symbol is ignored when writing; when reading, it will be the absolute section symbol. The addend is the displacement in bytes of the "lda" instruction from the "ldah" instruction (which is at the address of this reloc).

BFD\_RELOC\_ALPHA\_GPDISP\_LO16

For GPDISP\_LO16 ("ignore") relocations, the symbol is handled as with GPDISP\_HI16 relocations. The addend is ignored when writing the relocations out, and is filled in with the file's GP value on reading, for convenience.

BFD\_RELOC\_ALPHA\_GPDISP

The ELF GPDISP relocation is exactly the same as the GPDISP\_HI16 relocation except that there is no accompanying GPDISP\_LO16 relocation.

BFD\_RELOC\_ALPHA\_LITERAL

BFD\_RELOC\_ALPHA\_ELF\_LITERAL

**BFD\_RELOC\_ALPHA\_LITUSE**

The Alpha LITERAL/LITUSE relocs are produced by a symbol reference; the assembler turns it into a LDQ instruction to load the address of the symbol, and then fills in a register in the real instruction.

The LITERAL reloc, at the LDQ instruction, refers to the .lita section symbol. The addend is ignored when writing, but is filled in with the file's GP value on reading, for convenience, as with the GPDISP\_LO16 reloc.

The ELF\_LITERAL reloc is somewhere between 16\_GOTOFF and GPDISP\_LO16. It should refer to the symbol to be referenced, as with 16\_GOTOFF, but it generates output not based on the position within the .got section, but relative to the GP value chosen for the file during the final link stage.

The LITUSE reloc, on the instruction using the loaded address, gives information to the linker that it might be able to use to optimize away some literal section references. The symbol is ignored (read as the absolute section symbol), and the "addend" indicates the type of instruction using the register: 1 - "memory" fmt insn 2 - byte-manipulation (byte offset reg) 3 - jsr (target of branch)

**BFD\_RELOC\_ALPHA\_HINT**

The HINT relocation indicates a value that should be filled into the "hint" field of a jmp/jsr/ret instruction, for possible branch- prediction logic which may be provided on some processors.

**BFD\_RELOC\_ALPHA\_LINKAGE**

The LINKAGE relocation outputs a linkage pair in the object file, which is filled by the linker.

**BFD\_RELOC\_ALPHA\_CODEADDR**

The CODEADDR relocation outputs a STO\_CA in the object file, which is filled by the linker.

**BFD\_RELOC\_ALPHA\_GPREL\_HI16****BFD\_RELOC\_ALPHA\_GPREL\_LO16**

The GPREL\_HI/LO relocations together form a 32-bit offset from the GP register.

**BFD\_RELOC\_ALPHA\_BRSGP**

Like BFD\_RELOC\_23\_PCREL\_S2, except that the source and target must share a common GP, and the target address is adjusted for STO\_ALPHA\_STD\_GPLOAD.

**BFD\_RELOC\_ALPHA\_NOP**

The NOP relocation outputs a NOP if the longword displacement between two procedure entry points is  $< 2^{21}$ .

**BFD\_RELOC\_ALPHA\_BSR**

The BSR relocation outputs a BSR if the longword displacement between two procedure entry points is  $< 2^{21}$ .

**BFD\_RELOC\_ALPHA\_LDA**

The LDA relocation outputs a LDA if the longword displacement between two procedure entry points is  $< 2^{16}$ .



**BFD\_RELOC\_ALPHA\_BOH**

The BOH relocation outputs a BSR if the longword displacement between two procedure entry points is  $< 2^{21}$ , or else a hint.

**BFD\_RELOC\_ALPHA\_TLSGD****BFD\_RELOC\_ALPHA\_TLSDM****BFD\_RELOC\_ALPHA\_DTPMOD64****BFD\_RELOC\_ALPHA\_GOTDTPREL16****BFD\_RELOC\_ALPHA\_DTPREL64****BFD\_RELOC\_ALPHA\_DTPREL\_HI16****BFD\_RELOC\_ALPHA\_DTPREL\_LO16****BFD\_RELOC\_ALPHA\_DTPREL16****BFD\_RELOC\_ALPHA\_GOTTPREL16****BFD\_RELOC\_ALPHA\_TPREL64****BFD\_RELOC\_ALPHA\_TPREL\_HI16****BFD\_RELOC\_ALPHA\_TPREL\_LO16****BFD\_RELOC\_ALPHA\_TPREL16**

Alpha thread-local storage relocations.

**BFD\_RELOC\_MIPS\_JMP**

Bits 27..2 of the relocation address shifted right 2 bits; simple reloc otherwise.

**BFD\_RELOC\_MIPS16\_JMP**

The MIPS16 jump instruction.

**BFD\_RELOC\_MIPS16\_GPREL**

MIPS16 GP relative reloc.

**BFD\_RELOC\_HI16**

High 16 bits of 32-bit value; simple reloc.

**BFD\_RELOC\_HI16\_S**

High 16 bits of 32-bit value but the low 16 bits will be sign extended and added to form the final result. If the low 16 bits form a negative number, we need to add one to the high value to compensate for the borrow when the low bits are added.

**BFD\_RELOC\_LO16**

Low 16 bits.

**BFD\_RELOC\_HI16\_PCREL**

High 16 bits of 32-bit pc-relative value

**BFD\_RELOC\_HI16\_S\_PCREL**

High 16 bits of 32-bit pc-relative value, adjusted

**BFD\_RELOC\_LO16\_PCREL**

Low 16 bits of pc-relative value

**BFD\_RELOC\_MIPS16\_GOT16****BFD\_RELOC\_MIPS16\_CALL16**

Equivalent of BFD\_RELOC\_MIPS\_\*, but with the MIPS16 layout of 16-bit immediate fields

BFD\_RELOC\_MIPS16\_HI16

MIPS16 high 16 bits of 32-bit value.

BFD\_RELOC\_MIPS16\_HI16\_S

MIPS16 high 16 bits of 32-bit value but the low 16 bits will be sign extended and added to form the final result. If the low 16 bits form a negative number, we need to add one to the high value to compensate for the borrow when the low bits are added.

BFD\_RELOC\_MIPS16\_LO16

MIPS16 low 16 bits.

BFD\_RELOC\_MIPS\_LITERAL

Relocation against a MIPS literal section.

BFD\_RELOC\_MIPS\_GOT16

BFD\_RELOC\_MIPS\_CALL16

BFD\_RELOC\_MIPS\_GOT\_HI16

BFD\_RELOC\_MIPS\_GOT\_LO16

BFD\_RELOC\_MIPS\_CALL\_HI16

BFD\_RELOC\_MIPS\_CALL\_LO16

BFD\_RELOC\_MIPS\_SUB

BFD\_RELOC\_MIPS\_GOT\_PAGE

BFD\_RELOC\_MIPS\_GOT\_OFST

BFD\_RELOC\_MIPS\_GOT\_DISP

BFD\_RELOC\_MIPS\_SHIFT5

BFD\_RELOC\_MIPS\_SHIFT6

BFD\_RELOC\_MIPS\_INSERT\_A

BFD\_RELOC\_MIPS\_INSERT\_B

BFD\_RELOC\_MIPS\_DELETE

BFD\_RELOC\_MIPS\_HIGHEST

BFD\_RELOC\_MIPS\_HIGHER

BFD\_RELOC\_MIPS\_SCN\_DISP

BFD\_RELOC\_MIPS\_REL16

BFD\_RELOC\_MIPS\_RELGOT

BFD\_RELOC\_MIPS\_JALR

BFD\_RELOC\_MIPS\_TLS\_DTPMOD32

BFD\_RELOC\_MIPS\_TLS\_DTPREL32

BFD\_RELOC\_MIPS\_TLS\_DTPMOD64

BFD\_RELOC\_MIPS\_TLS\_DTPREL64

BFD\_RELOC\_MIPS\_TLS\_GD

BFD\_RELOC\_MIPS\_TLS\_LDM

BFD\_RELOC\_MIPS\_TLS\_DTPREL\_HI16

BFD\_RELOC\_MIPS\_TLS\_DTPREL\_LO16

BFD\_RELOC\_MIPS\_TLS\_GOTTREL

BFD\_RELOC\_MIPS\_TLS\_TPREL32

BFD\_RELOC\_MIPS\_TLS\_TPREL64

BFD\_RELOC\_MIPS\_TLS\_TPREL\_HI16

BFD\_RELOC\_MIPS\_TLS\_TPREL\_LO16

MIPS ELF relocations.

BFD\_RELOC\_MIPS\_COPY

BFD\_RELOC\_MIPS\_JUMP\_SLOT

MIPS ELF relocations (VxWorks and PLT extensions).

BFD\_RELOC\_MOXIE\_10\_PCREL

Moxie ELF relocations.

BFD\_RELOC\_FRV\_LABEL16

BFD\_RELOC\_FRV\_LABEL24

BFD\_RELOC\_FRV\_LO16

BFD\_RELOC\_FRV\_HI16

BFD\_RELOC\_FRV\_GPREL12

BFD\_RELOC\_FRV\_GPRELU12

BFD\_RELOC\_FRV\_GPREL32

BFD\_RELOC\_FRV\_GPRELHI

BFD\_RELOC\_FRV\_GPRELLO

BFD\_RELOC\_FRV\_GOT12

BFD\_RELOC\_FRV\_GOTHI

BFD\_RELOC\_FRV\_GOTLO

BFD\_RELOC\_FRV\_FUNCDESC

BFD\_RELOC\_FRV\_FUNCDESC\_GOT12

BFD\_RELOC\_FRV\_FUNCDESC\_GOTHI

BFD\_RELOC\_FRV\_FUNCDESC\_GOTLO

BFD\_RELOC\_FRV\_FUNCDESC\_VALUE

BFD\_RELOC\_FRV\_FUNCDESC\_GOTOFF12

BFD\_RELOC\_FRV\_FUNCDESC\_GOTOFFHI

BFD\_RELOC\_FRV\_FUNCDESC\_GOTOFFLO

BFD\_RELOC\_FRV\_GOTOFF12

BFD\_RELOC\_FRV\_GOTOFFHI

BFD\_RELOC\_FRV\_GOTOFFLO

BFD\_RELOC\_FRV\_GETTLSOFF

BFD\_RELOC\_FRV\_TLSDESC\_VALUE

BFD\_RELOC\_FRV\_GOTTLSDESC12

BFD\_RELOC\_FRV\_GOTTLSDESCHI

BFD\_RELOC\_FRV\_GOTTLSDESCLO

BFD\_RELOC\_FRV\_TLSMOFF12

BFD\_RELOC\_FRV\_TLSMOFFHI

BFD\_RELOC\_FRV\_TLSMOFFLO

BFD\_RELOC\_FRV\_GOTTLSOFF12

BFD\_RELOC\_FRV\_GOTTLSOFFHI

BFD\_RELOC\_FRV\_GOTTLSOFFLO

BFD\_RELOC\_FRV\_TLSOFF

BFD\_RELOC\_FRV\_TLSDESC\_RELAX

BFD\_RELOC\_FRV\_GETTLSOFF\_RELAX

BFD\_RELOC\_FRV\_TLSOFF\_RELAX

BFD\_RELOC\_FRV\_TLSMOFF

Fujitsu Frv Relocations.

**BFD\_RELOC\_MN10300\_GOTOFF24**

This is a 24bit GOT-relative reloc for the mn10300.

**BFD\_RELOC\_MN10300\_GOT32**

This is a 32bit GOT-relative reloc for the mn10300, offset by two bytes in the instruction.

**BFD\_RELOC\_MN10300\_GOT24**

This is a 24bit GOT-relative reloc for the mn10300, offset by two bytes in the instruction.

**BFD\_RELOC\_MN10300\_GOT16**

This is a 16bit GOT-relative reloc for the mn10300, offset by two bytes in the instruction.

**BFD\_RELOC\_MN10300\_COPY**

Copy symbol at runtime.

**BFD\_RELOC\_MN10300\_GLOB\_DAT**

Create GOT entry.

**BFD\_RELOC\_MN10300\_JMP\_SLOT**

Create PLT entry.

**BFD\_RELOC\_MN10300\_RELATIVE**

Adjust by program base.

**BFD\_RELOC\_MN10300\_SYM\_DIFF**

Together with another reloc targeted at the same location, allows for a value that is the difference of two symbols in the same section.

**BFD\_RELOC\_MN10300\_ALIGN**

The addend of this reloc is an alignment power that must be honoured at the offset's location, regardless of linker relaxation.

**BFD\_RELOC\_386\_GOT32****BFD\_RELOC\_386\_PLT32****BFD\_RELOC\_386\_COPY****BFD\_RELOC\_386\_GLOB\_DAT****BFD\_RELOC\_386\_JUMP\_SLOT****BFD\_RELOC\_386\_RELATIVE****BFD\_RELOC\_386\_GOTOFF****BFD\_RELOC\_386\_GOTPC****BFD\_RELOC\_386\_TLS\_TPOFF****BFD\_RELOC\_386\_TLS\_IE****BFD\_RELOC\_386\_TLS\_GOTIE****BFD\_RELOC\_386\_TLS\_LE****BFD\_RELOC\_386\_TLS\_GD****BFD\_RELOC\_386\_TLS\_LDM****BFD\_RELOC\_386\_TLS\_LDO\_32**

BFD\_RELOC\_386\_TLS\_IE\_32  
BFD\_RELOC\_386\_TLS\_LE\_32  
BFD\_RELOC\_386\_TLS\_DTPMOD32  
BFD\_RELOC\_386\_TLS\_DTPOFF32  
BFD\_RELOC\_386\_TLS\_TPOFF32  
BFD\_RELOC\_386\_TLS\_GOTDESC  
BFD\_RELOC\_386\_TLS\_DESC\_CALL  
BFD\_RELOC\_386\_TLS\_DESC  
BFD\_RELOC\_386\_IRELATIVE

i386/elf relocations

BFD\_RELOC\_X86\_64\_GOT32  
BFD\_RELOC\_X86\_64\_PLT32  
BFD\_RELOC\_X86\_64\_COPY  
BFD\_RELOC\_X86\_64\_GLOB\_DAT  
BFD\_RELOC\_X86\_64\_JUMP\_SLOT  
BFD\_RELOC\_X86\_64\_RELATIVE  
BFD\_RELOC\_X86\_64\_GOTPCREL  
BFD\_RELOC\_X86\_64\_32S  
BFD\_RELOC\_X86\_64\_DTPMOD64  
BFD\_RELOC\_X86\_64\_DTPOFF64  
BFD\_RELOC\_X86\_64\_TPOFF64  
BFD\_RELOC\_X86\_64\_TLSGD  
BFD\_RELOC\_X86\_64\_TLSLD  
BFD\_RELOC\_X86\_64\_DTPOFF32  
BFD\_RELOC\_X86\_64\_GOTTPOFF  
BFD\_RELOC\_X86\_64\_TPOFF32  
BFD\_RELOC\_X86\_64\_GOTOFF64  
BFD\_RELOC\_X86\_64\_GOTPC32  
BFD\_RELOC\_X86\_64\_GOT64  
BFD\_RELOC\_X86\_64\_GOTPCREL64  
BFD\_RELOC\_X86\_64\_GOTPC64  
BFD\_RELOC\_X86\_64\_GOTPLT64  
BFD\_RELOC\_X86\_64\_PLTOFF64  
BFD\_RELOC\_X86\_64\_GOTPC32\_TLSDESC  
BFD\_RELOC\_X86\_64\_TLSDESC\_CALL  
BFD\_RELOC\_X86\_64\_TLSDESC  
BFD\_RELOC\_X86\_64\_IRELATIVE

x86-64/elf relocations

BFD\_RELOC\_NS32K\_IMM\_8  
BFD\_RELOC\_NS32K\_IMM\_16  
BFD\_RELOC\_NS32K\_IMM\_32  
BFD\_RELOC\_NS32K\_IMM\_8\_PCREL  
BFD\_RELOC\_NS32K\_IMM\_16\_PCREL  
BFD\_RELOC\_NS32K\_IMM\_32\_PCREL  
BFD\_RELOC\_NS32K\_DISP\_8  
BFD\_RELOC\_NS32K\_DISP\_16

BFD\_RELOC\_NS32K\_DISP\_32  
BFD\_RELOC\_NS32K\_DISP\_8\_PCREL  
BFD\_RELOC\_NS32K\_DISP\_16\_PCREL  
BFD\_RELOC\_NS32K\_DISP\_32\_PCREL  
    ns32k relocations

BFD\_RELOC\_PDP11\_DISP\_8\_PCREL  
BFD\_RELOC\_PDP11\_DISP\_6\_PCREL  
    PDP11 relocations

BFD\_RELOC\_PJ\_CODE\_HI16  
BFD\_RELOC\_PJ\_CODE\_LO16  
BFD\_RELOC\_PJ\_CODE\_DIR16  
BFD\_RELOC\_PJ\_CODE\_DIR32  
BFD\_RELOC\_PJ\_CODE\_REL16  
BFD\_RELOC\_PJ\_CODE\_REL32

    Picojava relocs. Not all of these appear in object files.

BFD\_RELOC\_PPC\_B26  
BFD\_RELOC\_PPC\_BA26  
BFD\_RELOC\_PPC\_TOC16  
BFD\_RELOC\_PPC\_B16  
BFD\_RELOC\_PPC\_B16\_BRTAKEN  
BFD\_RELOC\_PPC\_B16\_BRNTAKEN  
BFD\_RELOC\_PPC\_BA16  
BFD\_RELOC\_PPC\_BA16\_BRTAKEN  
BFD\_RELOC\_PPC\_BA16\_BRNTAKEN  
BFD\_RELOC\_PPC\_COPY  
BFD\_RELOC\_PPC\_GLOB\_DAT  
BFD\_RELOC\_PPC\_JMP\_SLOT  
BFD\_RELOC\_PPC\_RELATIVE  
BFD\_RELOC\_PPC\_LOCAL24PC  
BFD\_RELOC\_PPC\_EMB\_NADDR32  
BFD\_RELOC\_PPC\_EMB\_NADDR16  
BFD\_RELOC\_PPC\_EMB\_NADDR16\_LO  
BFD\_RELOC\_PPC\_EMB\_NADDR16\_HI  
BFD\_RELOC\_PPC\_EMB\_NADDR16\_HA  
BFD\_RELOC\_PPC\_EMB\_SDAI16  
BFD\_RELOC\_PPC\_EMB\_SDA2I16  
BFD\_RELOC\_PPC\_EMB\_SDA2REL  
BFD\_RELOC\_PPC\_EMB\_SDA21  
BFD\_RELOC\_PPC\_EMB\_MRKREF  
BFD\_RELOC\_PPC\_EMB\_RELSEC16  
BFD\_RELOC\_PPC\_EMB\_RELST\_LO  
BFD\_RELOC\_PPC\_EMB\_RELST\_HI  
BFD\_RELOC\_PPC\_EMB\_RELST\_HA  
BFD\_RELOC\_PPC\_EMB\_BIT\_FLD  
BFD\_RELOC\_PPC\_EMB\_RELSDA

```

BFD_RELOC_PPC64_HIGHER
BFD_RELOC_PPC64_HIGHER_S
BFD_RELOC_PPC64_HIGHEST
BFD_RELOC_PPC64_HIGHEST_S
BFD_RELOC_PPC64_TOC16_LO
BFD_RELOC_PPC64_TOC16_HI
BFD_RELOC_PPC64_TOC16_HA
BFD_RELOC_PPC64_TOC
BFD_RELOC_PPC64_PLTGOT16
BFD_RELOC_PPC64_PLTGOT16_LO
BFD_RELOC_PPC64_PLTGOT16_HI
BFD_RELOC_PPC64_PLTGOT16_HA
BFD_RELOC_PPC64_ADDR16_DS
BFD_RELOC_PPC64_ADDR16_LO_DS
BFD_RELOC_PPC64_GOT16_DS
BFD_RELOC_PPC64_GOT16_LO_DS
BFD_RELOC_PPC64_PLT16_LO_DS
BFD_RELOC_PPC64_SECTOFF_DS
BFD_RELOC_PPC64_SECTOFF_LO_DS
BFD_RELOC_PPC64_TOC16_DS
BFD_RELOC_PPC64_TOC16_LO_DS
BFD_RELOC_PPC64_PLTGOT16_DS
BFD_RELOC_PPC64_PLTGOT16_LO_DS

```

Power(rs6000) and PowerPC relocations.

```

BFD_RELOC_PPC_TLS
BFD_RELOC_PPC_TLSGD
BFD_RELOC_PPC_TLSLD
BFD_RELOC_PPC_DTPMOD
BFD_RELOC_PPC_TPREL16
BFD_RELOC_PPC_TPREL16_LO
BFD_RELOC_PPC_TPREL16_HI
BFD_RELOC_PPC_TPREL16_HA
BFD_RELOC_PPC_TPREL
BFD_RELOC_PPC_DTPREL16
BFD_RELOC_PPC_DTPREL16_LO
BFD_RELOC_PPC_DTPREL16_HI
BFD_RELOC_PPC_DTPREL16_HA
BFD_RELOC_PPC_DTPREL
BFD_RELOC_PPC_GOT_TLSGD16
BFD_RELOC_PPC_GOT_TLSGD16_LO
BFD_RELOC_PPC_GOT_TLSGD16_HI
BFD_RELOC_PPC_GOT_TLSGD16_HA
BFD_RELOC_PPC_GOT_TLSLD16
BFD_RELOC_PPC_GOT_TLSLD16_LO
BFD_RELOC_PPC_GOT_TLSLD16_HI
BFD_RELOC_PPC_GOT_TLSLD16_HA

```

BFD\_RELOC\_PPC\_GOT\_TPREL16  
 BFD\_RELOC\_PPC\_GOT\_TPREL16\_LO  
 BFD\_RELOC\_PPC\_GOT\_TPREL16\_HI  
 BFD\_RELOC\_PPC\_GOT\_TPREL16\_HA  
 BFD\_RELOC\_PPC\_GOT\_DTPREL16  
 BFD\_RELOC\_PPC\_GOT\_DTPREL16\_LO  
 BFD\_RELOC\_PPC\_GOT\_DTPREL16\_HI  
 BFD\_RELOC\_PPC\_GOT\_DTPREL16\_HA  
 BFD\_RELOC\_PPC64\_TPREL16\_DS  
 BFD\_RELOC\_PPC64\_TPREL16\_LO\_DS  
 BFD\_RELOC\_PPC64\_TPREL16\_HIGHER  
 BFD\_RELOC\_PPC64\_TPREL16\_HIGHERA  
 BFD\_RELOC\_PPC64\_TPREL16\_HIGHEST  
 BFD\_RELOC\_PPC64\_TPREL16\_HIGHESTA  
 BFD\_RELOC\_PPC64\_DTPREL16\_DS  
 BFD\_RELOC\_PPC64\_DTPREL16\_LO\_DS  
 BFD\_RELOC\_PPC64\_DTPREL16\_HIGHER  
 BFD\_RELOC\_PPC64\_DTPREL16\_HIGHERA  
 BFD\_RELOC\_PPC64\_DTPREL16\_HIGHEST  
 BFD\_RELOC\_PPC64\_DTPREL16\_HIGHESTA

PowerPC and PowerPC64 thread-local storage relocations.

BFD\_RELOC\_I370\_D12  
 IBM 370/390 relocations

BFD\_RELOC\_CTOR

The type of reloc used to build a constructor table - at the moment probably a 32 bit wide absolute relocation, but the target can choose. It generally does map to one of the other relocation types.

BFD\_RELOC\_ARM\_PCREL\_BRANCH

ARM 26 bit pc-relative branch. The lowest two bits must be zero and are not stored in the instruction.

BFD\_RELOC\_ARM\_PCREL\_BLX

ARM 26 bit pc-relative branch. The lowest bit must be zero and is not stored in the instruction. The 2nd lowest bit comes from a 1 bit field in the instruction.

BFD\_RELOC\_THUMB\_PCREL\_BLX

Thumb 22 bit pc-relative branch. The lowest bit must be zero and is not stored in the instruction. The 2nd lowest bit comes from a 1 bit field in the instruction.

BFD\_RELOC\_ARM\_PCREL\_CALL

ARM 26-bit pc-relative branch for an unconditional BL or BLX instruction.

BFD\_RELOC\_ARM\_PCREL\_JUMP

ARM 26-bit pc-relative branch for B or conditional BL instruction.

BFD\_RELOC\_THUMB\_PCREL\_BRANCH7  
 BFD\_RELOC\_THUMB\_PCREL\_BRANCH9



BFD\_RELOC\_THUMB\_PCREL\_BRANCH12

BFD\_RELOC\_THUMB\_PCREL\_BRANCH20

BFD\_RELOC\_THUMB\_PCREL\_BRANCH23

BFD\_RELOC\_THUMB\_PCREL\_BRANCH25

Thumb 7-, 9-, 12-, 20-, 23-, and 25-bit pc-relative branches. The lowest bit must be zero and is not stored in the instruction. Note that the corresponding ELF R\_ARM\_THM\_JUMPnn constant has an "nn" one smaller in all cases. Note further that BRANCH23 corresponds to R\_ARM\_THM\_CALL.

BFD\_RELOC\_ARM\_OFFSET\_IMM

12-bit immediate offset, used in ARM-format ldr and str instructions.

BFD\_RELOC\_ARM\_THUMB\_OFFSET

5-bit immediate offset, used in Thumb-format ldr and str instructions.

BFD\_RELOC\_ARM\_TARGET1

Pc-relative or absolute relocation depending on target. Used for entries in .init\_array sections.

BFD\_RELOC\_ARM\_ROSEGREL32

Read-only segment base relative address.

BFD\_RELOC\_ARM\_SBREL32

Data segment base relative address.

BFD\_RELOC\_ARM\_TARGET2

This reloc is used for references to RTTI data from exception handling tables. The actual definition depends on the target. It may be a pc-relative or some form of GOT-indirect relocation.

BFD\_RELOC\_ARM\_PREL31

31-bit PC relative address.

BFD\_RELOC\_ARM\_MOVW

BFD\_RELOC\_ARM\_MOVT

BFD\_RELOC\_ARM\_MOVW\_PCREL

BFD\_RELOC\_ARM\_MOVT\_PCREL

BFD\_RELOC\_ARM\_THUMB\_MOVW

BFD\_RELOC\_ARM\_THUMB\_MOVT

BFD\_RELOC\_ARM\_THUMB\_MOVW\_PCREL

BFD\_RELOC\_ARM\_THUMB\_MOVT\_PCREL

Low and High halfword relocations for MOVW and MOVT instructions.

BFD\_RELOC\_ARM\_JUMP\_SLOT

BFD\_RELOC\_ARM\_GLOB\_DAT

BFD\_RELOC\_ARM\_GOT32

BFD\_RELOC\_ARM\_PLT32

BFD\_RELOC\_ARM\_RELATIVE

BFD\_RELOC\_ARM\_GOTOFF

BFD\_RELOC\_ARM\_GOTPC

**BFD\_RELOC\_ARM\_GOT\_PREL**

Relocations for setting up GOTs and PLTs for shared libraries.

BFD\_RELOC\_ARM\_TLS\_GD32  
BFD\_RELOC\_ARM\_TLS\_LD032  
BFD\_RELOC\_ARM\_TLS\_LDM32  
BFD\_RELOC\_ARM\_TLS\_DTPOFF32  
BFD\_RELOC\_ARM\_TLS\_DTPMOD32  
BFD\_RELOC\_ARM\_TLS\_TPOFF32  
BFD\_RELOC\_ARM\_TLS\_IE32  
BFD\_RELOC\_ARM\_TLS\_LE32  
BFD\_RELOC\_ARM\_TLS\_GOTDESC  
BFD\_RELOC\_ARM\_TLS\_CALL  
BFD\_RELOC\_ARM\_THM\_TLS\_CALL  
BFD\_RELOC\_ARM\_TLS\_DESCSEQ  
BFD\_RELOC\_ARM\_THM\_TLS\_DESCSEQ  
BFD\_RELOC\_ARM\_TLS\_DESC

ARM thread-local storage relocations.

BFD\_RELOC\_ARM\_ALU\_PC\_G0\_NC  
BFD\_RELOC\_ARM\_ALU\_PC\_G0  
BFD\_RELOC\_ARM\_ALU\_PC\_G1\_NC  
BFD\_RELOC\_ARM\_ALU\_PC\_G1  
BFD\_RELOC\_ARM\_ALU\_PC\_G2  
BFD\_RELOC\_ARM\_LDR\_PC\_G0  
BFD\_RELOC\_ARM\_LDR\_PC\_G1  
BFD\_RELOC\_ARM\_LDR\_PC\_G2  
BFD\_RELOC\_ARM\_LDRS\_PC\_G0  
BFD\_RELOC\_ARM\_LDRS\_PC\_G1  
BFD\_RELOC\_ARM\_LDRS\_PC\_G2  
BFD\_RELOC\_ARM\_LDC\_PC\_G0  
BFD\_RELOC\_ARM\_LDC\_PC\_G1  
BFD\_RELOC\_ARM\_LDC\_PC\_G2  
BFD\_RELOC\_ARM\_ALU\_SB\_G0\_NC  
BFD\_RELOC\_ARM\_ALU\_SB\_G0  
BFD\_RELOC\_ARM\_ALU\_SB\_G1\_NC  
BFD\_RELOC\_ARM\_ALU\_SB\_G1  
BFD\_RELOC\_ARM\_ALU\_SB\_G2  
BFD\_RELOC\_ARM\_LDR\_SB\_G0  
BFD\_RELOC\_ARM\_LDR\_SB\_G1  
BFD\_RELOC\_ARM\_LDR\_SB\_G2  
BFD\_RELOC\_ARM\_LDRS\_SB\_G0  
BFD\_RELOC\_ARM\_LDRS\_SB\_G1  
BFD\_RELOC\_ARM\_LDRS\_SB\_G2  
BFD\_RELOC\_ARM\_LDC\_SB\_G0  
BFD\_RELOC\_ARM\_LDC\_SB\_G1  
BFD\_RELOC\_ARM\_LDC\_SB\_G2

ARM group relocations.

BFD\_RELOC\_ARM\_V4BX

Annotation of BX instructions.

BFD\_RELOC\_ARM\_IRELATIVE

ARM support for STT\_GNU\_IFUNC.

BFD\_RELOC\_ARM\_IMMEDIATE

BFD\_RELOC\_ARM\_ADRL\_IMMEDIATE

BFD\_RELOC\_ARM\_T32\_IMMEDIATE

BFD\_RELOC\_ARM\_T32\_ADD\_IMM

BFD\_RELOC\_ARM\_T32\_IMM12

BFD\_RELOC\_ARM\_T32\_ADD\_PC12

BFD\_RELOC\_ARM\_SHIFT\_IMM

BFD\_RELOC\_ARM\_SMC

BFD\_RELOC\_ARM\_HVC

BFD\_RELOC\_ARM\_SWI

BFD\_RELOC\_ARM\_MULTT

BFD\_RELOC\_ARM\_CP\_OFF\_IMM

BFD\_RELOC\_ARM\_CP\_OFF\_IMM\_S2

BFD\_RELOC\_ARM\_T32\_CP\_OFF\_IMM

BFD\_RELOC\_ARM\_T32\_CP\_OFF\_IMM\_S2

BFD\_RELOC\_ARM\_ADR\_IMM

BFD\_RELOC\_ARM\_LDR\_IMM

BFD\_RELOC\_ARM\_LITERAL

BFD\_RELOC\_ARM\_IN\_POOL

BFD\_RELOC\_ARM\_OFFSET\_IMM8

BFD\_RELOC\_ARM\_T32\_OFFSET\_U8

BFD\_RELOC\_ARM\_T32\_OFFSET\_IMM

BFD\_RELOC\_ARM\_HWLITERAL

BFD\_RELOC\_ARM\_THUMB\_ADD

BFD\_RELOC\_ARM\_THUMB\_IMM

BFD\_RELOC\_ARM\_THUMB\_SHIFT

These relocs are only used within the ARM assembler. They are not (at present) written to any object files.

BFD\_RELOC\_SH\_PCDISP8BY2

BFD\_RELOC\_SH\_PCDISP12BY2

BFD\_RELOC\_SH\_IMM3

BFD\_RELOC\_SH\_IMM3U

BFD\_RELOC\_SH\_DISP12

BFD\_RELOC\_SH\_DISP12BY2

BFD\_RELOC\_SH\_DISP12BY4

BFD\_RELOC\_SH\_DISP12BY8

BFD\_RELOC\_SH\_DISP20

BFD\_RELOC\_SH\_DISP20BY8

BFD\_RELOC\_SH\_IMM4

BFD\_RELOC\_SH\_IMM4BY2

BFD\_RELOC\_SH\_IMM4BY4

BFD\_RELOC\_SH\_IMM8  
BFD\_RELOC\_SH\_IMM8BY2  
BFD\_RELOC\_SH\_IMM8BY4  
BFD\_RELOC\_SH\_PCRELIMM8BY2  
BFD\_RELOC\_SH\_PCRELIMM8BY4  
BFD\_RELOC\_SH\_SWITCH16  
BFD\_RELOC\_SH\_SWITCH32  
BFD\_RELOC\_SH\_USES  
BFD\_RELOC\_SH\_COUNT  
BFD\_RELOC\_SH\_ALIGN  
BFD\_RELOC\_SH\_CODE  
BFD\_RELOC\_SH\_DATA  
BFD\_RELOC\_SH\_LABEL  
BFD\_RELOC\_SH\_LOOP\_START  
BFD\_RELOC\_SH\_LOOP\_END  
BFD\_RELOC\_SH\_COPY  
BFD\_RELOC\_SH\_GLOB\_DAT  
BFD\_RELOC\_SH\_JMP\_SLOT  
BFD\_RELOC\_SH\_RELATIVE  
BFD\_RELOC\_SH\_GOTPC  
BFD\_RELOC\_SH\_GOT\_LOW16  
BFD\_RELOC\_SH\_GOT\_MEDLOW16  
BFD\_RELOC\_SH\_GOT\_MEDHI16  
BFD\_RELOC\_SH\_GOT\_HI16  
BFD\_RELOC\_SH\_GOTPLT\_LOW16  
BFD\_RELOC\_SH\_GOTPLT\_MEDLOW16  
BFD\_RELOC\_SH\_GOTPLT\_MEDHI16  
BFD\_RELOC\_SH\_GOTPLT\_HI16  
BFD\_RELOC\_SH\_PLT\_LOW16  
BFD\_RELOC\_SH\_PLT\_MEDLOW16  
BFD\_RELOC\_SH\_PLT\_MEDHI16  
BFD\_RELOC\_SH\_PLT\_HI16  
BFD\_RELOC\_SH\_GOTOFF\_LOW16  
BFD\_RELOC\_SH\_GOTOFF\_MEDLOW16  
BFD\_RELOC\_SH\_GOTOFF\_MEDHI16  
BFD\_RELOC\_SH\_GOTOFF\_HI16  
BFD\_RELOC\_SH\_GOTPC\_LOW16  
BFD\_RELOC\_SH\_GOTPC\_MEDLOW16  
BFD\_RELOC\_SH\_GOTPC\_MEDHI16  
BFD\_RELOC\_SH\_GOTPC\_HI16  
BFD\_RELOC\_SH\_COPY64  
BFD\_RELOC\_SH\_GLOB\_DAT64  
BFD\_RELOC\_SH\_JMP\_SLOT64  
BFD\_RELOC\_SH\_RELATIVE64  
BFD\_RELOC\_SH\_GOT10BY4  
BFD\_RELOC\_SH\_GOT10BY8  
BFD\_RELOC\_SH\_GOTPLT10BY4

```

BFD_RELOC_SH_GOTPLT10BY8
BFD_RELOC_SH_GOTPLT32
BFD_RELOC_SH_SHMEDIA_CODE
BFD_RELOC_SH_IMMU5
BFD_RELOC_SH_IMMS6
BFD_RELOC_SH_IMMS6BY32
BFD_RELOC_SH_IMMU6
BFD_RELOC_SH_IMMS10
BFD_RELOC_SH_IMMS10BY2
BFD_RELOC_SH_IMMS10BY4
BFD_RELOC_SH_IMMS10BY8
BFD_RELOC_SH_IMMS16
BFD_RELOC_SH_IMMU16
BFD_RELOC_SH_IMM_LOW16
BFD_RELOC_SH_IMM_LOW16_PCREL
BFD_RELOC_SH_IMM_MEDLOW16
BFD_RELOC_SH_IMM_MEDLOW16_PCREL
BFD_RELOC_SH_IMM_MEDHI16
BFD_RELOC_SH_IMM_MEDHI16_PCREL
BFD_RELOC_SH_IMM_HI16
BFD_RELOC_SH_IMM_HI16_PCREL
BFD_RELOC_SH_PT_16
BFD_RELOC_SH_TLS_GD_32
BFD_RELOC_SH_TLS_LD_32
BFD_RELOC_SH_TLS_LDO_32
BFD_RELOC_SH_TLS_IE_32
BFD_RELOC_SH_TLS_LE_32
BFD_RELOC_SH_TLS_DTPMOD32
BFD_RELOC_SH_TLS_DTPOFF32
BFD_RELOC_SH_TLS_TPOFF32
BFD_RELOC_SH_GOT20
BFD_RELOC_SH_GOTOFF20
BFD_RELOC_SH_GOTFUNCDESC
BFD_RELOC_SH_GOTFUNCDESC20
BFD_RELOC_SH_GOTOFFFUNCDESC
BFD_RELOC_SH_GOTOFFFUNCDESC20
BFD_RELOC_SH_FUNCDESC

```

Renesas / SuperH SH relocs. Not all of these appear in object files.

```
BFD_RELOC_ARC_B22_PCREL
```

ARC Cores relocs. ARC 22 bit pc-relative branch. The lowest two bits must be zero and are not stored in the instruction. The high 20 bits are installed in bits 26 through 7 of the instruction.

```
BFD_RELOC_ARC_B26
```

ARC 26 bit absolute branch. The lowest two bits must be zero and are not stored in the instruction. The high 24 bits are installed in bits 23 through 0.

BFD\_RELOC\_BFIN\_16\_IMM

ADI Blackfin 16 bit immediate absolute reloc.

BFD\_RELOC\_BFIN\_16\_HIGH

ADI Blackfin 16 bit immediate absolute reloc higher 16 bits.

BFD\_RELOC\_BFIN\_4\_PCREL

ADI Blackfin 'a' part of LSETUP.

BFD\_RELOC\_BFIN\_5\_PCREL

ADI Blackfin.

BFD\_RELOC\_BFIN\_16\_LOW

ADI Blackfin 16 bit immediate absolute reloc lower 16 bits.

BFD\_RELOC\_BFIN\_10\_PCREL

ADI Blackfin.

BFD\_RELOC\_BFIN\_11\_PCREL

ADI Blackfin 'b' part of LSETUP.

BFD\_RELOC\_BFIN\_12\_PCREL\_JUMP

ADI Blackfin.

BFD\_RELOC\_BFIN\_12\_PCREL\_JUMP\_S

ADI Blackfin Short jump, pcrel.

BFD\_RELOC\_BFIN\_24\_PCREL\_CALL\_X

ADI Blackfin Call.x not implemented.

BFD\_RELOC\_BFIN\_24\_PCREL\_JUMP\_L

ADI Blackfin Long Jump pcrel.

BFD\_RELOC\_BFIN\_GOT17M4

BFD\_RELOC\_BFIN\_GOTHI

BFD\_RELOC\_BFIN\_GOTLO

BFD\_RELOC\_BFIN\_FUNCDESC

BFD\_RELOC\_BFIN\_FUNCDESC\_GOT17M4

BFD\_RELOC\_BFIN\_FUNCDESC\_GOTHI

BFD\_RELOC\_BFIN\_FUNCDESC\_GOTLO

BFD\_RELOC\_BFIN\_FUNCDESC\_VALUE

BFD\_RELOC\_BFIN\_FUNCDESC\_GTOFF17M4

BFD\_RELOC\_BFIN\_FUNCDESC\_GTOFFHI

BFD\_RELOC\_BFIN\_FUNCDESC\_GTOFFLO

BFD\_RELOC\_BFIN\_GTOFF17M4

BFD\_RELOC\_BFIN\_GTOFFHI

BFD\_RELOC\_BFIN\_GTOFFLO

ADI Blackfin FD-PIC relocations.

BFD\_RELOC\_BFIN\_GOT

ADI Blackfin GOT relocation.

BFD\_RELOC\_BFIN\_PLTPC  
ADI Blackfin PLTPC relocation.

BFD\_ARELOC\_BFIN\_PUSH  
ADI Blackfin arithmetic relocation.

BFD\_ARELOC\_BFIN\_CONST  
ADI Blackfin arithmetic relocation.

BFD\_ARELOC\_BFIN\_ADD  
ADI Blackfin arithmetic relocation.

BFD\_ARELOC\_BFIN\_SUB  
ADI Blackfin arithmetic relocation.

BFD\_ARELOC\_BFIN\_MULT  
ADI Blackfin arithmetic relocation.

BFD\_ARELOC\_BFIN\_DIV  
ADI Blackfin arithmetic relocation.

BFD\_ARELOC\_BFIN\_MOD  
ADI Blackfin arithmetic relocation.

BFD\_ARELOC\_BFIN\_LSHIFT  
ADI Blackfin arithmetic relocation.

BFD\_ARELOC\_BFIN\_RSHIFT  
ADI Blackfin arithmetic relocation.

BFD\_ARELOC\_BFIN\_AND  
ADI Blackfin arithmetic relocation.

BFD\_ARELOC\_BFIN\_OR  
ADI Blackfin arithmetic relocation.

BFD\_ARELOC\_BFIN\_XOR  
ADI Blackfin arithmetic relocation.

BFD\_ARELOC\_BFIN\_LAND  
ADI Blackfin arithmetic relocation.

BFD\_ARELOC\_BFIN\_LOR  
ADI Blackfin arithmetic relocation.

BFD\_ARELOC\_BFIN\_LEN  
ADI Blackfin arithmetic relocation.

BFD\_ARELOC\_BFIN\_NEG  
ADI Blackfin arithmetic relocation.

BFD\_ARELOC\_BFIN\_COMP  
ADI Blackfin arithmetic relocation.

**BFD\_ARELOC\_BFIN\_PAGE**

ADI Blackfin arithmetic relocation.

**BFD\_ARELOC\_BFIN\_HWPAGE**

ADI Blackfin arithmetic relocation.

**BFD\_ARELOC\_BFIN\_ADDR**

ADI Blackfin arithmetic relocation.

**BFD\_RELOC\_D10V\_10\_PCREL\_R**

Mitsubishi D10V relocs. This is a 10-bit reloc with the right 2 bits assumed to be 0.

**BFD\_RELOC\_D10V\_10\_PCREL\_L**

Mitsubishi D10V relocs. This is a 10-bit reloc with the right 2 bits assumed to be 0. This is the same as the previous reloc except it is in the left container, i.e., shifted left 15 bits.

**BFD\_RELOC\_D10V\_18**

This is an 18-bit reloc with the right 2 bits assumed to be 0.

**BFD\_RELOC\_D10V\_18\_PCREL**

This is an 18-bit reloc with the right 2 bits assumed to be 0.

**BFD\_RELOC\_D30V\_6**

Mitsubishi D30V relocs. This is a 6-bit absolute reloc.

**BFD\_RELOC\_D30V\_9\_PCREL**

This is a 6-bit pc-relative reloc with the right 3 bits assumed to be 0.

**BFD\_RELOC\_D30V\_9\_PCREL\_R**

This is a 6-bit pc-relative reloc with the right 3 bits assumed to be 0. Same as the previous reloc but on the right side of the container.

**BFD\_RELOC\_D30V\_15**

This is a 12-bit absolute reloc with the right 3 bits assumed to be 0.

**BFD\_RELOC\_D30V\_15\_PCREL**

This is a 12-bit pc-relative reloc with the right 3 bits assumed to be 0.

**BFD\_RELOC\_D30V\_15\_PCREL\_R**

This is a 12-bit pc-relative reloc with the right 3 bits assumed to be 0. Same as the previous reloc but on the right side of the container.

**BFD\_RELOC\_D30V\_21**

This is an 18-bit absolute reloc with the right 3 bits assumed to be 0.

**BFD\_RELOC\_D30V\_21\_PCREL**

This is an 18-bit pc-relative reloc with the right 3 bits assumed to be 0.

**BFD\_RELOC\_D30V\_21\_PCREL\_R**

This is an 18-bit pc-relative reloc with the right 3 bits assumed to be 0. Same as the previous reloc but on the right side of the container.



BFD\_RELOC\_D30V\_32

This is a 32-bit absolute reloc.

BFD\_RELOC\_D30V\_32\_PCREL

This is a 32-bit pc-relative reloc.

BFD\_RELOC\_DLX\_HI16\_S

DLX relocs

BFD\_RELOC\_DLX\_LO16

DLX relocs

BFD\_RELOC\_DLX\_JMP26

DLX relocs

BFD\_RELOC\_M32C\_HI8

BFD\_RELOC\_M32C\_RL\_JUMP

BFD\_RELOC\_M32C\_RL\_1ADDR

BFD\_RELOC\_M32C\_RL\_2ADDR

Renesas M16C/M32C Relocations.

BFD\_RELOC\_M32R\_24

Renesas M32R (formerly Mitsubishi M32R) relocs. This is a 24 bit absolute address.

BFD\_RELOC\_M32R\_10\_PCREL

This is a 10-bit pc-relative reloc with the right 2 bits assumed to be 0.

BFD\_RELOC\_M32R\_18\_PCREL

This is an 18-bit reloc with the right 2 bits assumed to be 0.

BFD\_RELOC\_M32R\_26\_PCREL

This is a 26-bit reloc with the right 2 bits assumed to be 0.

BFD\_RELOC\_M32R\_HI16\_ULO

This is a 16-bit reloc containing the high 16 bits of an address used when the lower 16 bits are treated as unsigned.

BFD\_RELOC\_M32R\_HI16\_SLO

This is a 16-bit reloc containing the high 16 bits of an address used when the lower 16 bits are treated as signed.

BFD\_RELOC\_M32R\_LO16

This is a 16-bit reloc containing the lower 16 bits of an address.

BFD\_RELOC\_M32R\_SDA16

This is a 16-bit reloc containing the small data area offset for use in add3, load, and store instructions.

BFD\_RELOC\_M32R\_GOT24

BFD\_RELOC\_M32R\_26\_PLTREL

BFD\_RELOC\_M32R\_COPY

BFD\_RELOC\_M32R\_GLOB\_DAT

BFD\_RELOC\_M32R\_JMP\_SLOT  
 BFD\_RELOC\_M32R\_RELATIVE  
 BFD\_RELOC\_M32R\_GOTOFF  
 BFD\_RELOC\_M32R\_GOTOFF\_HI\_UL0  
 BFD\_RELOC\_M32R\_GOTOFF\_HI\_SLO  
 BFD\_RELOC\_M32R\_GOTOFF\_LO  
 BFD\_RELOC\_M32R\_GOTPC24  
 BFD\_RELOC\_M32R\_GOT16\_HI\_UL0  
 BFD\_RELOC\_M32R\_GOT16\_HI\_SLO  
 BFD\_RELOC\_M32R\_GOT16\_LO  
 BFD\_RELOC\_M32R\_GOTPC\_HI\_UL0  
 BFD\_RELOC\_M32R\_GOTPC\_HI\_SLO  
 BFD\_RELOC\_M32R\_GOTPC\_LO

For PIC.

BFD\_RELOC\_V850\_9\_PCREL

This is a 9-bit reloc

BFD\_RELOC\_V850\_22\_PCREL

This is a 22-bit reloc

BFD\_RELOC\_V850\_SDA\_16\_16\_OFFSET

This is a 16 bit offset from the short data area pointer.

BFD\_RELOC\_V850\_SDA\_15\_16\_OFFSET

This is a 16 bit offset (of which only 15 bits are used) from the short data area pointer.

BFD\_RELOC\_V850\_ZDA\_16\_16\_OFFSET

This is a 16 bit offset from the zero data area pointer.

BFD\_RELOC\_V850\_ZDA\_15\_16\_OFFSET

This is a 16 bit offset (of which only 15 bits are used) from the zero data area pointer.

BFD\_RELOC\_V850\_TDA\_6\_8\_OFFSET

This is an 8 bit offset (of which only 6 bits are used) from the tiny data area pointer.

BFD\_RELOC\_V850\_TDA\_7\_8\_OFFSET

This is an 8bit offset (of which only 7 bits are used) from the tiny data area pointer.

BFD\_RELOC\_V850\_TDA\_7\_7\_OFFSET

This is a 7 bit offset from the tiny data area pointer.

BFD\_RELOC\_V850\_TDA\_16\_16\_OFFSET

This is a 16 bit offset from the tiny data area pointer.

BFD\_RELOC\_V850\_TDA\_4\_5\_OFFSET

This is a 5 bit offset (of which only 4 bits are used) from the tiny data area pointer.

BFD\_RELOC\_V850\_TDA\_4\_4\_OFFSET

This is a 4 bit offset from the tiny data area pointer.

**BFD\_RELOC\_V850\_SDA\_16\_16\_SPLIT\_OFFSET**

This is a 16 bit offset from the short data area pointer, with the bits placed non-contiguously in the instruction.

**BFD\_RELOC\_V850\_ZDA\_16\_16\_SPLIT\_OFFSET**

This is a 16 bit offset from the zero data area pointer, with the bits placed non-contiguously in the instruction.

**BFD\_RELOC\_V850\_CALLT\_6\_7\_OFFSET**

This is a 6 bit offset from the call table base pointer.

**BFD\_RELOC\_V850\_CALLT\_16\_16\_OFFSET**

This is a 16 bit offset from the call table base pointer.

**BFD\_RELOC\_V850\_LONGCALL**

Used for relaxing indirect function calls.

**BFD\_RELOC\_V850\_LONGJUMP**

Used for relaxing indirect jumps.

**BFD\_RELOC\_V850\_ALIGN**

Used to maintain alignment whilst relaxing.

**BFD\_RELOC\_V850\_LO16\_SPLIT\_OFFSET**

This is a variation of BFD\_RELOC\_LO16 that can be used in v850e ld.bu instructions.

**BFD\_RELOC\_V850\_16\_PCREL**

This is a 16-bit reloc.

**BFD\_RELOC\_V850\_17\_PCREL**

This is a 17-bit reloc.

**BFD\_RELOC\_V850\_23**

This is a 23-bit reloc.

**BFD\_RELOC\_V850\_32\_PCREL**

This is a 32-bit reloc.

**BFD\_RELOC\_V850\_32\_ABS**

This is a 32-bit reloc.

**BFD\_RELOC\_V850\_16\_SPLIT\_OFFSET**

This is a 16-bit reloc.

**BFD\_RELOC\_V850\_16\_S1**

This is a 16-bit reloc.

**BFD\_RELOC\_V850\_LO16\_S1**

Low 16 bits. 16 bit shifted by 1.

**BFD\_RELOC\_V850\_CALLT\_15\_16\_OFFSET**

This is a 16 bit offset from the call table base pointer.

BFD\_RELOC\_V850\_32\_GOTPCREL  
DSO relocations.

BFD\_RELOC\_V850\_16\_GOT  
DSO relocations.

BFD\_RELOC\_V850\_32\_GOT  
DSO relocations.

BFD\_RELOC\_V850\_22\_PLT\_PCREL  
DSO relocations.

BFD\_RELOC\_V850\_32\_PLT\_PCREL  
DSO relocations.

BFD\_RELOC\_V850\_COPY  
DSO relocations.

BFD\_RELOC\_V850\_GLOB\_DAT  
DSO relocations.

BFD\_RELOC\_V850\_JMP\_SLOT  
DSO relocations.

BFD\_RELOC\_V850\_RELATIVE  
DSO relocations.

BFD\_RELOC\_V850\_16\_GOTOFF  
DSO relocations.

BFD\_RELOC\_V850\_32\_GOTOFF  
DSO relocations.

BFD\_RELOC\_V850\_CODE  
start code.

BFD\_RELOC\_V850\_DATA  
start data in text.

BFD\_RELOC\_MN10300\_32\_PCREL  
This is a 32bit pcrel reloc for the mn10300, offset by two bytes in the instruction.

BFD\_RELOC\_MN10300\_16\_PCREL  
This is a 16bit pcrel reloc for the mn10300, offset by two bytes in the instruction.

BFD\_RELOC\_TIC30\_LDP  
This is a 8bit DP reloc for the tms320c30, where the most significant 8 bits of a 24 bit word are placed into the least significant 8 bits of the opcode.

BFD\_RELOC\_TIC54X\_PARTLS7  
This is a 7bit reloc for the tms320c54x, where the least significant 7 bits of a 16 bit word are placed into the least significant 7 bits of the opcode.

**BFD\_RELOC\_TIC54X\_PARTMS9**

This is a 9bit DP reloc for the tms320c54x, where the most significant 9 bits of a 16 bit word are placed into the least significant 9 bits of the opcode.

**BFD\_RELOC\_TIC54X\_23**

This is an extended address 23-bit reloc for the tms320c54x.

**BFD\_RELOC\_TIC54X\_16\_OF\_23**

This is a 16-bit reloc for the tms320c54x, where the least significant 16 bits of a 23-bit extended address are placed into the opcode.

**BFD\_RELOC\_TIC54X\_MS7\_OF\_23**

This is a reloc for the tms320c54x, where the most significant 7 bits of a 23-bit extended address are placed into the opcode.

**BFD\_RELOC\_C6000\_PCR\_S21****BFD\_RELOC\_C6000\_PCR\_S12****BFD\_RELOC\_C6000\_PCR\_S10****BFD\_RELOC\_C6000\_PCR\_S7****BFD\_RELOC\_C6000\_ABS\_S16****BFD\_RELOC\_C6000\_ABS\_L16****BFD\_RELOC\_C6000\_ABS\_H16****BFD\_RELOC\_C6000\_SBR\_U15\_B****BFD\_RELOC\_C6000\_SBR\_U15\_H****BFD\_RELOC\_C6000\_SBR\_U15\_W****BFD\_RELOC\_C6000\_SBR\_S16****BFD\_RELOC\_C6000\_SBR\_L16\_B****BFD\_RELOC\_C6000\_SBR\_L16\_H****BFD\_RELOC\_C6000\_SBR\_L16\_W****BFD\_RELOC\_C6000\_SBR\_H16\_B****BFD\_RELOC\_C6000\_SBR\_H16\_H****BFD\_RELOC\_C6000\_SBR\_H16\_W****BFD\_RELOC\_C6000\_SBR\_GOT\_U15\_W****BFD\_RELOC\_C6000\_SBR\_GOT\_L16\_W****BFD\_RELOC\_C6000\_SBR\_GOT\_H16\_W****BFD\_RELOC\_C6000\_DSBT\_INDEX****BFD\_RELOC\_C6000\_PREL31****BFD\_RELOC\_C6000\_COPY****BFD\_RELOC\_C6000\_JUMP\_SLOT****BFD\_RELOC\_C6000\_EHTYPE****BFD\_RELOC\_C6000\_PCR\_H16****BFD\_RELOC\_C6000\_PCR\_L16****BFD\_RELOC\_C6000\_ALIGN****BFD\_RELOC\_C6000\_FPHEAD****BFD\_RELOC\_C6000\_NOCMP**

TMS320C6000 relocations.

**BFD\_RELOC\_FR30\_48**

This is a 48 bit reloc for the FR30 that stores 32 bits.

BFD\_RELOC\_FR30\_20

This is a 32 bit reloc for the FR30 that stores 20 bits split up into two sections.

BFD\_RELOC\_FR30\_6\_IN\_4

This is a 16 bit reloc for the FR30 that stores a 6 bit word offset in 4 bits.

BFD\_RELOC\_FR30\_8\_IN\_8

This is a 16 bit reloc for the FR30 that stores an 8 bit byte offset into 8 bits.

BFD\_RELOC\_FR30\_9\_IN\_8

This is a 16 bit reloc for the FR30 that stores a 9 bit short offset into 8 bits.

BFD\_RELOC\_FR30\_10\_IN\_8

This is a 16 bit reloc for the FR30 that stores a 10 bit word offset into 8 bits.

BFD\_RELOC\_FR30\_9\_PCREL

This is a 16 bit reloc for the FR30 that stores a 9 bit pc relative short offset into 8 bits.

BFD\_RELOC\_FR30\_12\_PCREL

This is a 16 bit reloc for the FR30 that stores a 12 bit pc relative short offset into 11 bits.

BFD\_RELOC\_MCORE\_PCREL\_IMM8BY4

BFD\_RELOC\_MCORE\_PCREL\_IMM11BY2

BFD\_RELOC\_MCORE\_PCREL\_IMM4BY2

BFD\_RELOC\_MCORE\_PCREL\_32

BFD\_RELOC\_MCORE\_PCREL\_JSR\_IMM11BY2

BFD\_RELOC\_MCORE\_RVA

Motorola Mcore relocations.

BFD\_RELOC\_MEP\_8

BFD\_RELOC\_MEP\_16

BFD\_RELOC\_MEP\_32

BFD\_RELOC\_MEP\_PCREL8A2

BFD\_RELOC\_MEP\_PCREL12A2

BFD\_RELOC\_MEP\_PCREL17A2

BFD\_RELOC\_MEP\_PCREL24A2

BFD\_RELOC\_MEP\_PCABS24A2

BFD\_RELOC\_MEP\_LOW16

BFD\_RELOC\_MEP\_HI16U

BFD\_RELOC\_MEP\_HI16S

BFD\_RELOC\_MEP\_GPREL

BFD\_RELOC\_MEP\_TPREL

BFD\_RELOC\_MEP\_TPREL7

BFD\_RELOC\_MEP\_TPREL7A2

BFD\_RELOC\_MEP\_TPREL7A4

BFD\_RELOC\_MEP\_UIMM24

BFD\_RELOC\_MEP\_ADDR24A4

BFD\_RELOC\_MEP\_GNU\_VTINHERIT

**BFD\_RELOC\_MEP\_GNU\_VTENTRY**

Toshiba Media Processor Relocations.

**BFD\_RELOC\_MMIX\_GETA****BFD\_RELOC\_MMIX\_GETA\_1****BFD\_RELOC\_MMIX\_GETA\_2****BFD\_RELOC\_MMIX\_GETA\_3**

These are relocations for the GETA instruction.

**BFD\_RELOC\_MMIX\_CBRANCH****BFD\_RELOC\_MMIX\_CBRANCH\_J****BFD\_RELOC\_MMIX\_CBRANCH\_1****BFD\_RELOC\_MMIX\_CBRANCH\_2****BFD\_RELOC\_MMIX\_CBRANCH\_3**

These are relocations for a conditional branch instruction.

**BFD\_RELOC\_MMIX\_PUSHJ****BFD\_RELOC\_MMIX\_PUSHJ\_1****BFD\_RELOC\_MMIX\_PUSHJ\_2****BFD\_RELOC\_MMIX\_PUSHJ\_3****BFD\_RELOC\_MMIX\_PUSHJ\_STUBBABLE**

These are relocations for the PUSHJ instruction.

**BFD\_RELOC\_MMIX\_JMP****BFD\_RELOC\_MMIX\_JMP\_1****BFD\_RELOC\_MMIX\_JMP\_2****BFD\_RELOC\_MMIX\_JMP\_3**

These are relocations for the JMP instruction.

**BFD\_RELOC\_MMIX\_ADDR19**

This is a relocation for a relative address as in a GETA instruction or a branch.

**BFD\_RELOC\_MMIX\_ADDR27**

This is a relocation for a relative address as in a JMP instruction.

**BFD\_RELOC\_MMIX\_REG\_OR\_BYTE**

This is a relocation for an instruction field that may be a general register or a value 0..255.

**BFD\_RELOC\_MMIX\_REG**

This is a relocation for an instruction field that may be a general register.

**BFD\_RELOC\_MMIX\_BASE\_PLUS\_OFFSET**

This is a relocation for two instruction fields holding a register and an offset, the equivalent of the relocation.

**BFD\_RELOC\_MMIX\_LOCAL**

This relocation is an assertion that the expression is not allocated as a global register. It does not modify contents.

**BFD\_RELOC\_AVR\_7\_PCREL**

This is a 16 bit reloc for the AVR that stores 8 bit pc relative short offset into 7 bits.

**BFD\_RELOC\_AVR\_13\_PCREL**

This is a 16 bit reloc for the AVR that stores 13 bit pc relative short offset into 12 bits.

**BFD\_RELOC\_AVR\_16\_PM**

This is a 16 bit reloc for the AVR that stores 17 bit value (usually program memory address) into 16 bits.

**BFD\_RELOC\_AVR\_L08\_LDI**

This is a 16 bit reloc for the AVR that stores 8 bit value (usually data memory address) into 8 bit immediate value of LDI insn.

**BFD\_RELOC\_AVR\_HI8\_LDI**

This is a 16 bit reloc for the AVR that stores 8 bit value (high 8 bit of data memory address) into 8 bit immediate value of LDI insn.

**BFD\_RELOC\_AVR\_HH8\_LDI**

This is a 16 bit reloc for the AVR that stores 8 bit value (most high 8 bit of program memory address) into 8 bit immediate value of LDI insn.

**BFD\_RELOC\_AVR\_MS8\_LDI**

This is a 16 bit reloc for the AVR that stores 8 bit value (most high 8 bit of 32 bit value) into 8 bit immediate value of LDI insn.

**BFD\_RELOC\_AVR\_L08\_LDI\_NEG**

This is a 16 bit reloc for the AVR that stores negated 8 bit value (usually data memory address) into 8 bit immediate value of SUBI insn.

**BFD\_RELOC\_AVR\_HI8\_LDI\_NEG**

This is a 16 bit reloc for the AVR that stores negated 8 bit value (high 8 bit of data memory address) into 8 bit immediate value of SUBI insn.

**BFD\_RELOC\_AVR\_HH8\_LDI\_NEG**

This is a 16 bit reloc for the AVR that stores negated 8 bit value (most high 8 bit of program memory address) into 8 bit immediate value of LDI or SUBI insn.

**BFD\_RELOC\_AVR\_MS8\_LDI\_NEG**

This is a 16 bit reloc for the AVR that stores negated 8 bit value (msb of 32 bit value) into 8 bit immediate value of LDI insn.

**BFD\_RELOC\_AVR\_L08\_LDI\_PM**

This is a 16 bit reloc for the AVR that stores 8 bit value (usually command address) into 8 bit immediate value of LDI insn.

**BFD\_RELOC\_AVR\_L08\_LDI\_GS**

This is a 16 bit reloc for the AVR that stores 8 bit value (command address) into 8 bit immediate value of LDI insn. If the address is beyond the 128k boundary, the linker inserts a jump stub for this reloc in the lower 128k.



**BFD\_RELOC\_AVR\_HI8\_LDI\_PM**

This is a 16 bit reloc for the AVR that stores 8 bit value (high 8 bit of command address) into 8 bit immediate value of LDI insn.

**BFD\_RELOC\_AVR\_HI8\_LDI\_GS**

This is a 16 bit reloc for the AVR that stores 8 bit value (high 8 bit of command address) into 8 bit immediate value of LDI insn. If the address is beyond the 128k boundary, the linker inserts a jump stub for this reloc below 128k.

**BFD\_RELOC\_AVR\_HH8\_LDI\_PM**

This is a 16 bit reloc for the AVR that stores 8 bit value (most high 8 bit of command address) into 8 bit immediate value of LDI insn.

**BFD\_RELOC\_AVR\_LO8\_LDI\_PM\_NEG**

This is a 16 bit reloc for the AVR that stores negated 8 bit value (usually command address) into 8 bit immediate value of SUBI insn.

**BFD\_RELOC\_AVR\_HI8\_LDI\_PM\_NEG**

This is a 16 bit reloc for the AVR that stores negated 8 bit value (high 8 bit of 16 bit command address) into 8 bit immediate value of SUBI insn.

**BFD\_RELOC\_AVR\_HH8\_LDI\_PM\_NEG**

This is a 16 bit reloc for the AVR that stores negated 8 bit value (high 6 bit of 22 bit command address) into 8 bit immediate value of SUBI insn.

**BFD\_RELOC\_AVR\_CALL**

This is a 32 bit reloc for the AVR that stores 23 bit value into 22 bits.

**BFD\_RELOC\_AVR\_LDI**

This is a 16 bit reloc for the AVR that stores all needed bits for absolute addressing with ldi with overflow check to linktime

**BFD\_RELOC\_AVR\_6**

This is a 6 bit reloc for the AVR that stores offset for ldd/std instructions

**BFD\_RELOC\_AVR\_6\_ADIW**

This is a 6 bit reloc for the AVR that stores offset for adiw/sbiw instructions

**BFD\_RELOC\_RX\_NEG8****BFD\_RELOC\_RX\_NEG16****BFD\_RELOC\_RX\_NEG24****BFD\_RELOC\_RX\_NEG32****BFD\_RELOC\_RX\_16\_OP****BFD\_RELOC\_RX\_24\_OP****BFD\_RELOC\_RX\_32\_OP****BFD\_RELOC\_RX\_8U****BFD\_RELOC\_RX\_16U****BFD\_RELOC\_RX\_24U****BFD\_RELOC\_RX\_DIR3U\_PCREL****BFD\_RELOC\_RX\_DIFF**

BFD\_RELOC\_RX\_GPRELB  
BFD\_RELOC\_RX\_GPRELW  
BFD\_RELOC\_RX\_GPRELL  
BFD\_RELOC\_RX\_SYM  
BFD\_RELOC\_RX\_OP\_SUBTRACT  
BFD\_RELOC\_RX\_OP\_NEG  
BFD\_RELOC\_RX\_ABS8  
BFD\_RELOC\_RX\_ABS16  
BFD\_RELOC\_RX\_ABS16\_REV  
BFD\_RELOC\_RX\_ABS32  
BFD\_RELOC\_RX\_ABS32\_REV  
BFD\_RELOC\_RX\_ABS16U  
BFD\_RELOC\_RX\_ABS16UW  
BFD\_RELOC\_RX\_ABS16UL  
BFD\_RELOC\_RX\_RELAX  
    Renesas RX Relocations.

BFD\_RELOC\_390\_12  
    Direct 12 bit.

BFD\_RELOC\_390\_GOT12  
    12 bit GOT offset.

BFD\_RELOC\_390\_PLT32  
    32 bit PC relative PLT address.

BFD\_RELOC\_390\_COPY  
    Copy symbol at runtime.

BFD\_RELOC\_390\_GLOB\_DAT  
    Create GOT entry.

BFD\_RELOC\_390\_JMP\_SLOT  
    Create PLT entry.

BFD\_RELOC\_390\_RELATIVE  
    Adjust by program base.

BFD\_RELOC\_390\_GOTPC  
    32 bit PC relative offset to GOT.

BFD\_RELOC\_390\_GOT16  
    16 bit GOT offset.

BFD\_RELOC\_390\_PC16DBL  
    PC relative 16 bit shifted by 1.

BFD\_RELOC\_390\_PLT16DBL  
    16 bit PC rel. PLT shifted by 1.

BFD\_RELOC\_390\_PC32DBL  
    PC relative 32 bit shifted by 1.

BFD\_RELOC\_390\_PLT32DBL  
32 bit PC rel. PLT shifted by 1.

BFD\_RELOC\_390\_GOTPCDBL  
32 bit PC rel. GOT shifted by 1.

BFD\_RELOC\_390\_GOT64  
64 bit GOT offset.

BFD\_RELOC\_390\_PLT64  
64 bit PC relative PLT address.

BFD\_RELOC\_390\_GOTENT  
32 bit rel. offset to GOT entry.

BFD\_RELOC\_390\_GOTOFF64  
64 bit offset to GOT.

BFD\_RELOC\_390\_GOTPLT12  
12-bit offset to symbol-entry within GOT, with PLT handling.

BFD\_RELOC\_390\_GOTPLT16  
16-bit offset to symbol-entry within GOT, with PLT handling.

BFD\_RELOC\_390\_GOTPLT32  
32-bit offset to symbol-entry within GOT, with PLT handling.

BFD\_RELOC\_390\_GOTPLT64  
64-bit offset to symbol-entry within GOT, with PLT handling.

BFD\_RELOC\_390\_GOTPLTENT  
32-bit rel. offset to symbol-entry within GOT, with PLT handling.

BFD\_RELOC\_390\_PLTOFF16  
16-bit rel. offset from the GOT to a PLT entry.

BFD\_RELOC\_390\_PLTOFF32  
32-bit rel. offset from the GOT to a PLT entry.

BFD\_RELOC\_390\_PLTOFF64  
64-bit rel. offset from the GOT to a PLT entry.

BFD\_RELOC\_390\_TLS\_LOAD

BFD\_RELOC\_390\_TLS\_GDCALL

BFD\_RELOC\_390\_TLS\_LDCALL

BFD\_RELOC\_390\_TLS\_GD32

BFD\_RELOC\_390\_TLS\_GD64

BFD\_RELOC\_390\_TLS\_GOTIE12

BFD\_RELOC\_390\_TLS\_GOTIE32

BFD\_RELOC\_390\_TLS\_GOTIE64

BFD\_RELOC\_390\_TLS\_LDM32

BFD\_RELOC\_390\_TLS\_LDM64

BFD\_RELOC\_390\_TLS\_IE32  
 BFD\_RELOC\_390\_TLS\_IE64  
 BFD\_RELOC\_390\_TLS\_IEENT  
 BFD\_RELOC\_390\_TLS\_LE32  
 BFD\_RELOC\_390\_TLS\_LE64  
 BFD\_RELOC\_390\_TLS\_LD032  
 BFD\_RELOC\_390\_TLS\_LD064  
 BFD\_RELOC\_390\_TLS\_DTPMOD  
 BFD\_RELOC\_390\_TLS\_DTPOFF  
 BFD\_RELOC\_390\_TLS\_TPOFF

s390 tls relocations.

BFD\_RELOC\_390\_20  
 BFD\_RELOC\_390\_GOT20  
 BFD\_RELOC\_390\_GOTPLT20  
 BFD\_RELOC\_390\_TLS\_GOTIE20

Long displacement extension.

BFD\_RELOC\_SCORE\_GPREL15

Score relocations Low 16 bit for load/store

BFD\_RELOC\_SCORE\_DUMMY2

BFD\_RELOC\_SCORE\_JMP

This is a 24-bit reloc with the right 1 bit assumed to be 0

BFD\_RELOC\_SCORE\_BRANCH

This is a 19-bit reloc with the right 1 bit assumed to be 0

BFD\_RELOC\_SCORE\_IMM30

This is a 32-bit reloc for 48-bit instructions.

BFD\_RELOC\_SCORE\_IMM32

This is a 32-bit reloc for 48-bit instructions.

BFD\_RELOC\_SCORE16\_JMP

This is a 11-bit reloc with the right 1 bit assumed to be 0

BFD\_RELOC\_SCORE16\_BRANCH

This is a 8-bit reloc with the right 1 bit assumed to be 0

BFD\_RELOC\_SCORE\_BCMP

This is a 9-bit reloc with the right 1 bit assumed to be 0

BFD\_RELOC\_SCORE\_GOT15

BFD\_RELOC\_SCORE\_GOT\_L016

BFD\_RELOC\_SCORE\_CALL15

BFD\_RELOC\_SCORE\_DUMMY\_HI16

Undocumented Score relocs

BFD\_RELOC\_IP2K\_FR9

Scenix IP2K - 9-bit register number / data address

**BFD\_RELOC\_IP2K\_BANK**

Scenix IP2K - 4-bit register/data bank number

**BFD\_RELOC\_IP2K\_ADDR16CJP**

Scenix IP2K - low 13 bits of instruction word address

**BFD\_RELOC\_IP2K\_PAGE3**

Scenix IP2K - high 3 bits of instruction word address

**BFD\_RELOC\_IP2K\_LO8DATA**

**BFD\_RELOC\_IP2K\_HI8DATA**

**BFD\_RELOC\_IP2K\_EX8DATA**

Scenix IP2K - ext/low/high 8 bits of data address

**BFD\_RELOC\_IP2K\_LO8INSN**

**BFD\_RELOC\_IP2K\_HI8INSN**

Scenix IP2K - low/high 8 bits of instruction word address

**BFD\_RELOC\_IP2K\_PC\_SKIP**

Scenix IP2K - even/odd PC modifier to modify snb pcl.0

**BFD\_RELOC\_IP2K\_TEXT**

Scenix IP2K - 16 bit word address in text section.

**BFD\_RELOC\_IP2K\_FR\_OFFSET**

Scenix IP2K - 7-bit sp or dp offset

**BFD\_RELOC\_VPE4KMATH\_DATA**

**BFD\_RELOC\_VPE4KMATH\_INSN**

Scenix VPE4K coprocessor - data/insn-space addressing

**BFD\_RELOC\_VTABLE\_INHERIT**

**BFD\_RELOC\_VTABLE\_ENTRY**

These two relocations are used by the linker to determine which of the entries in a C++ virtual function table are actually used. When the `-gc-sections` option is given, the linker will zero out the entries that are not used, so that the code for those functions need not be included in the output.

**VTABLE\_INHERIT** is a zero-space relocation used to describe to the linker the inheritance tree of a C++ virtual function table. The relocation's symbol should be the parent class' vtable, and the relocation should be located at the child vtable.

**VTABLE\_ENTRY** is a zero-space relocation that describes the use of a virtual function table entry. The reloc's symbol should refer to the table of the class mentioned in the code. Off of that base, an offset describes the entry that is being used. For Rela hosts, this offset is stored in the reloc's addend. For Rel hosts, we are forced to put this offset in the reloc's section offset.

**BFD\_RELOC\_IA64\_IMM14**

**BFD\_RELOC\_IA64\_IMM22**

**BFD\_RELOC\_IA64\_IMM64**

**BFD\_RELOC\_IA64\_DIR32MSB**

BFD\_RELOC\_IA64\_DIR32LSB  
BFD\_RELOC\_IA64\_DIR64MSB  
BFD\_RELOC\_IA64\_DIR64LSB  
BFD\_RELOC\_IA64\_GPREL22  
BFD\_RELOC\_IA64\_GPREL64I  
BFD\_RELOC\_IA64\_GPREL32MSB  
BFD\_RELOC\_IA64\_GPREL32LSB  
BFD\_RELOC\_IA64\_GPREL64MSB  
BFD\_RELOC\_IA64\_GPREL64LSB  
BFD\_RELOC\_IA64\_LTOFF22  
BFD\_RELOC\_IA64\_LTOFF64I  
BFD\_RELOC\_IA64\_PLTOFF22  
BFD\_RELOC\_IA64\_PLTOFF64I  
BFD\_RELOC\_IA64\_PLTOFF64MSB  
BFD\_RELOC\_IA64\_PLTOFF64LSB  
BFD\_RELOC\_IA64\_FPTR64I  
BFD\_RELOC\_IA64\_FPTR32MSB  
BFD\_RELOC\_IA64\_FPTR32LSB  
BFD\_RELOC\_IA64\_FPTR64MSB  
BFD\_RELOC\_IA64\_FPTR64LSB  
BFD\_RELOC\_IA64\_PCREL21B  
BFD\_RELOC\_IA64\_PCREL21BI  
BFD\_RELOC\_IA64\_PCREL21M  
BFD\_RELOC\_IA64\_PCREL21F  
BFD\_RELOC\_IA64\_PCREL22  
BFD\_RELOC\_IA64\_PCREL60B  
BFD\_RELOC\_IA64\_PCREL64I  
BFD\_RELOC\_IA64\_PCREL32MSB  
BFD\_RELOC\_IA64\_PCREL32LSB  
BFD\_RELOC\_IA64\_PCREL64MSB  
BFD\_RELOC\_IA64\_PCREL64LSB  
BFD\_RELOC\_IA64\_LTOFF\_FPTR22  
BFD\_RELOC\_IA64\_LTOFF\_FPTR64I  
BFD\_RELOC\_IA64\_LTOFF\_FPTR32MSB  
BFD\_RELOC\_IA64\_LTOFF\_FPTR32LSB  
BFD\_RELOC\_IA64\_LTOFF\_FPTR64MSB  
BFD\_RELOC\_IA64\_LTOFF\_FPTR64LSB  
BFD\_RELOC\_IA64\_SEGREL32MSB  
BFD\_RELOC\_IA64\_SEGREL32LSB  
BFD\_RELOC\_IA64\_SEGREL64MSB  
BFD\_RELOC\_IA64\_SEGREL64LSB  
BFD\_RELOC\_IA64\_SECREL32MSB  
BFD\_RELOC\_IA64\_SECREL32LSB  
BFD\_RELOC\_IA64\_SECREL64MSB  
BFD\_RELOC\_IA64\_SECREL64LSB  
BFD\_RELOC\_IA64\_REL32MSB  
BFD\_RELOC\_IA64\_REL32LSB

BFD\_RELOC\_IA64\_REL64MSB  
 BFD\_RELOC\_IA64\_REL64LSB  
 BFD\_RELOC\_IA64\_LTV32MSB  
 BFD\_RELOC\_IA64\_LTV32LSB  
 BFD\_RELOC\_IA64\_LTV64MSB  
 BFD\_RELOC\_IA64\_LTV64LSB  
 BFD\_RELOC\_IA64\_IPLTMSB  
 BFD\_RELOC\_IA64\_IPLTLSB  
 BFD\_RELOC\_IA64\_COPY  
 BFD\_RELOC\_IA64\_LTOFF22X  
 BFD\_RELOC\_IA64\_LDXMOV  
 BFD\_RELOC\_IA64\_TPREL14  
 BFD\_RELOC\_IA64\_TPREL22  
 BFD\_RELOC\_IA64\_TPREL64I  
 BFD\_RELOC\_IA64\_TPREL64MSB  
 BFD\_RELOC\_IA64\_TPREL64LSB  
 BFD\_RELOC\_IA64\_LTOFF\_TPREL22  
 BFD\_RELOC\_IA64\_DTPMOD64MSB  
 BFD\_RELOC\_IA64\_DTPMOD64LSB  
 BFD\_RELOC\_IA64\_LTOFF\_DTPMOD22  
 BFD\_RELOC\_IA64\_DTPREL14  
 BFD\_RELOC\_IA64\_DTPREL22  
 BFD\_RELOC\_IA64\_DTPREL64I  
 BFD\_RELOC\_IA64\_DTPREL32MSB  
 BFD\_RELOC\_IA64\_DTPREL32LSB  
 BFD\_RELOC\_IA64\_DTPREL64MSB  
 BFD\_RELOC\_IA64\_DTPREL64LSB  
 BFD\_RELOC\_IA64\_LTOFF\_DTPREL22

Intel IA64 Relocations.

BFD\_RELOC\_M68HC11\_HI8

Motorola 68HC11 reloc. This is the 8 bit high part of an absolute address.

BFD\_RELOC\_M68HC11\_LO8

Motorola 68HC11 reloc. This is the 8 bit low part of an absolute address.

BFD\_RELOC\_M68HC11\_3B

Motorola 68HC11 reloc. This is the 3 bit of a value.

BFD\_RELOC\_M68HC11\_RL\_JUMP

Motorola 68HC11 reloc. This reloc marks the beginning of a jump/call instruction. It is used for linker relaxation to correctly identify beginning of instruction and change some branches to use PC-relative addressing mode.

BFD\_RELOC\_M68HC11\_RL\_GROUP

Motorola 68HC11 reloc. This reloc marks a group of several instructions that gcc generates and for which the linker relaxation pass can modify and/or remove some of them.

**BFD\_RELOC\_M68HC11\_L016**

Motorola 68HC11 reloc. This is the 16-bit lower part of an address. It is used for 'call' instruction to specify the symbol address without any special transformation (due to memory bank window).

**BFD\_RELOC\_M68HC11\_PAGE**

Motorola 68HC11 reloc. This is a 8-bit reloc that specifies the page number of an address. It is used by 'call' instruction to specify the page number of the symbol.

**BFD\_RELOC\_M68HC11\_24**

Motorola 68HC11 reloc. This is a 24-bit reloc that represents the address with a 16-bit value and a 8-bit page number. The symbol address is transformed to follow the 16K memory bank of 68HC12 (seen as mapped in the window).

**BFD\_RELOC\_M68HC12\_5B**

Motorola 68HC12 reloc. This is the 5 bits of a value.

**BFD\_RELOC\_16C\_NUM08****BFD\_RELOC\_16C\_NUM08\_C****BFD\_RELOC\_16C\_NUM16****BFD\_RELOC\_16C\_NUM16\_C****BFD\_RELOC\_16C\_NUM32****BFD\_RELOC\_16C\_NUM32\_C****BFD\_RELOC\_16C\_DISP04****BFD\_RELOC\_16C\_DISP04\_C****BFD\_RELOC\_16C\_DISP08****BFD\_RELOC\_16C\_DISP08\_C****BFD\_RELOC\_16C\_DISP16****BFD\_RELOC\_16C\_DISP16\_C****BFD\_RELOC\_16C\_DISP24****BFD\_RELOC\_16C\_DISP24\_C****BFD\_RELOC\_16C\_DISP24a****BFD\_RELOC\_16C\_DISP24a\_C****BFD\_RELOC\_16C\_REG04****BFD\_RELOC\_16C\_REG04\_C****BFD\_RELOC\_16C\_REG04a****BFD\_RELOC\_16C\_REG04a\_C****BFD\_RELOC\_16C\_REG14****BFD\_RELOC\_16C\_REG14\_C****BFD\_RELOC\_16C\_REG16****BFD\_RELOC\_16C\_REG16\_C****BFD\_RELOC\_16C\_REG20****BFD\_RELOC\_16C\_REG20\_C****BFD\_RELOC\_16C\_ABS20****BFD\_RELOC\_16C\_ABS20\_C****BFD\_RELOC\_16C\_ABS24****BFD\_RELOC\_16C\_ABS24\_C****BFD\_RELOC\_16C\_IMM04**



BFD\_RELOC\_16C\_IMM04\_C  
BFD\_RELOC\_16C\_IMM16  
BFD\_RELOC\_16C\_IMM16\_C  
BFD\_RELOC\_16C\_IMM20  
BFD\_RELOC\_16C\_IMM20\_C  
BFD\_RELOC\_16C\_IMM24  
BFD\_RELOC\_16C\_IMM24\_C  
BFD\_RELOC\_16C\_IMM32  
BFD\_RELOC\_16C\_IMM32\_C  
NS CR16C Relocations.

BFD\_RELOC\_CR16\_NUM8  
BFD\_RELOC\_CR16\_NUM16  
BFD\_RELOC\_CR16\_NUM32  
BFD\_RELOC\_CR16\_NUM32a  
BFD\_RELOC\_CR16\_REGREL0  
BFD\_RELOC\_CR16\_REGREL4  
BFD\_RELOC\_CR16\_REGREL4a  
BFD\_RELOC\_CR16\_REGREL14  
BFD\_RELOC\_CR16\_REGREL14a  
BFD\_RELOC\_CR16\_REGREL16  
BFD\_RELOC\_CR16\_REGREL20  
BFD\_RELOC\_CR16\_REGREL20a  
BFD\_RELOC\_CR16\_ABS20  
BFD\_RELOC\_CR16\_ABS24  
BFD\_RELOC\_CR16\_IMM4  
BFD\_RELOC\_CR16\_IMM8  
BFD\_RELOC\_CR16\_IMM16  
BFD\_RELOC\_CR16\_IMM20  
BFD\_RELOC\_CR16\_IMM24  
BFD\_RELOC\_CR16\_IMM32  
BFD\_RELOC\_CR16\_IMM32a  
BFD\_RELOC\_CR16\_DISP4  
BFD\_RELOC\_CR16\_DISP8  
BFD\_RELOC\_CR16\_DISP16  
BFD\_RELOC\_CR16\_DISP20  
BFD\_RELOC\_CR16\_DISP24  
BFD\_RELOC\_CR16\_DISP24a  
BFD\_RELOC\_CR16\_SWITCH8  
BFD\_RELOC\_CR16\_SWITCH16  
BFD\_RELOC\_CR16\_SWITCH32  
BFD\_RELOC\_CR16\_GOT\_REGREL20  
BFD\_RELOC\_CR16\_GOTC\_REGREL20  
BFD\_RELOC\_CR16\_GLOB\_DAT  
NS CR16 Relocations.

BFD\_RELOC\_CRX\_REL4  
BFD\_RELOC\_CRX\_REL8

BFD\_RELOC\_CRX\_REL8\_CMP  
BFD\_RELOC\_CRX\_REL16  
BFD\_RELOC\_CRX\_REL24  
BFD\_RELOC\_CRX\_REL32  
BFD\_RELOC\_CRX\_REGREL12  
BFD\_RELOC\_CRX\_REGREL22  
BFD\_RELOC\_CRX\_REGREL28  
BFD\_RELOC\_CRX\_REGREL32  
BFD\_RELOC\_CRX\_ABS16  
BFD\_RELOC\_CRX\_ABS32  
BFD\_RELOC\_CRX\_NUM8  
BFD\_RELOC\_CRX\_NUM16  
BFD\_RELOC\_CRX\_NUM32  
BFD\_RELOC\_CRX\_IMM16  
BFD\_RELOC\_CRX\_IMM32  
BFD\_RELOC\_CRX\_SWITCH8  
BFD\_RELOC\_CRX\_SWITCH16  
BFD\_RELOC\_CRX\_SWITCH32

NS CRX Relocations.

BFD\_RELOC\_CRIS\_BDISP8  
BFD\_RELOC\_CRIS\_UNSIGNED\_5  
BFD\_RELOC\_CRIS\_SIGNED\_6  
BFD\_RELOC\_CRIS\_UNSIGNED\_6  
BFD\_RELOC\_CRIS\_SIGNED\_8  
BFD\_RELOC\_CRIS\_UNSIGNED\_8  
BFD\_RELOC\_CRIS\_SIGNED\_16  
BFD\_RELOC\_CRIS\_UNSIGNED\_16  
BFD\_RELOC\_CRIS\_LAPCQ\_OFFSET  
BFD\_RELOC\_CRIS\_UNSIGNED\_4

These relocs are only used within the CRIS assembler. They are not (at present) written to any object files.

BFD\_RELOC\_CRIS\_COPY  
BFD\_RELOC\_CRIS\_GLOB\_DAT  
BFD\_RELOC\_CRIS\_JUMP\_SLOT  
BFD\_RELOC\_CRIS\_RELATIVE

Relocs used in ELF shared libraries for CRIS.

BFD\_RELOC\_CRIS\_32\_GOT  
32-bit offset to symbol-entry within GOT.

BFD\_RELOC\_CRIS\_16\_GOT  
16-bit offset to symbol-entry within GOT.

BFD\_RELOC\_CRIS\_32\_GOTPLT  
32-bit offset to symbol-entry within GOT, with PLT handling.

BFD\_RELOC\_CRIS\_16\_GOTPLT  
16-bit offset to symbol-entry within GOT, with PLT handling.

BFD\_RELOC\_CRIS\_32\_GOTREL

32-bit offset to symbol, relative to GOT.

BFD\_RELOC\_CRIS\_32\_PLT\_GOTREL

32-bit offset to symbol with PLT entry, relative to GOT.

BFD\_RELOC\_CRIS\_32\_PLT\_PCREL

32-bit offset to symbol with PLT entry, relative to this relocation.

BFD\_RELOC\_CRIS\_32\_GOT\_GD

BFD\_RELOC\_CRIS\_16\_GOT\_GD

BFD\_RELOC\_CRIS\_32\_GD

BFD\_RELOC\_CRIS\_DTP

BFD\_RELOC\_CRIS\_32\_DTPREL

BFD\_RELOC\_CRIS\_16\_DTPREL

BFD\_RELOC\_CRIS\_32\_GOT\_TPREL

BFD\_RELOC\_CRIS\_16\_GOT\_TPREL

BFD\_RELOC\_CRIS\_32\_TPREL

BFD\_RELOC\_CRIS\_16\_TPREL

BFD\_RELOC\_CRIS\_DTPMOD

BFD\_RELOC\_CRIS\_32\_IE

Relocs used in TLS code for CRIS.

BFD\_RELOC\_860\_COPY

BFD\_RELOC\_860\_GLOB\_DAT

BFD\_RELOC\_860\_JUMP\_SLOT

BFD\_RELOC\_860\_RELATIVE

BFD\_RELOC\_860\_PC26

BFD\_RELOC\_860\_PLT26

BFD\_RELOC\_860\_PC16

BFD\_RELOC\_860\_LOW0

BFD\_RELOC\_860\_SPLIT0

BFD\_RELOC\_860\_LOW1

BFD\_RELOC\_860\_SPLIT1

BFD\_RELOC\_860\_LOW2

BFD\_RELOC\_860\_SPLIT2

BFD\_RELOC\_860\_LOW3

BFD\_RELOC\_860\_LOGOTO

BFD\_RELOC\_860\_SPGOTO

BFD\_RELOC\_860\_LOGOT1

BFD\_RELOC\_860\_SPGOT1

BFD\_RELOC\_860\_LOGOTOFF0

BFD\_RELOC\_860\_SPGOTOFF0

BFD\_RELOC\_860\_LOGOTOFF1

BFD\_RELOC\_860\_SPGOTOFF1

BFD\_RELOC\_860\_LOGOTOFF2

BFD\_RELOC\_860\_LOGOTOFF3

BFD\_RELOC\_860\_LOPC

BFD\_RELOC\_860\_HIGHADJ  
BFD\_RELOC\_860\_HAGOT  
BFD\_RELOC\_860\_HAGTOFF  
BFD\_RELOC\_860\_HAPC  
BFD\_RELOC\_860\_HIGH  
BFD\_RELOC\_860\_HIGOT  
BFD\_RELOC\_860\_HIGTOFF

Intel i860 Relocations.

BFD\_RELOC\_OPENRISC\_ABS\_26  
BFD\_RELOC\_OPENRISC\_REL\_26

OpenRISC Relocations.

BFD\_RELOC\_H8\_DIR16A8  
BFD\_RELOC\_H8\_DIR16R8  
BFD\_RELOC\_H8\_DIR24A8  
BFD\_RELOC\_H8\_DIR24R8  
BFD\_RELOC\_H8\_DIR32A16

H8 elf Relocations.

BFD\_RELOC\_XSTORMY16\_REL\_12  
BFD\_RELOC\_XSTORMY16\_12  
BFD\_RELOC\_XSTORMY16\_24  
BFD\_RELOC\_XSTORMY16\_FPTR16

Sony Xstormy16 Relocations.

BFD\_RELOC\_RELC

Self-describing complex relocations.

BFD\_RELOC\_XC16X\_PAG  
BFD\_RELOC\_XC16X\_POF  
BFD\_RELOC\_XC16X\_SEG  
BFD\_RELOC\_XC16X\_SOF

Infineon Relocations.

BFD\_RELOC\_VAX\_GLOB\_DAT  
BFD\_RELOC\_VAX\_JMP\_SLOT  
BFD\_RELOC\_VAX\_RELATIVE

Relocations used by VAX ELF.

BFD\_RELOC\_MT\_PC16

Morpho MT - 16 bit immediate relocation.

BFD\_RELOC\_MT\_HI16

Morpho MT - Hi 16 bits of an address.

BFD\_RELOC\_MT\_LO16

Morpho MT - Low 16 bits of an address.

BFD\_RELOC\_MT\_GNU\_VTINHERIT

Morpho MT - Used to tell the linker which vtable entries are used.

**BFD\_RELOC\_MT\_GNU\_VTENTRY**

Morpho MT - Used to tell the linker which vtable entries are used.

**BFD\_RELOC\_MT\_PCINSN8**

Morpho MT - 8 bit immediate relocation.

**BFD\_RELOC\_MSP430\_10\_PCREL**

**BFD\_RELOC\_MSP430\_16\_PCREL**

**BFD\_RELOC\_MSP430\_16**

**BFD\_RELOC\_MSP430\_16\_PCREL\_BYTE**

**BFD\_RELOC\_MSP430\_16\_BYTE**

**BFD\_RELOC\_MSP430\_2X\_PCREL**

**BFD\_RELOC\_MSP430\_RL\_PCREL**

msp430 specific relocation codes

**BFD\_RELOC\_IQ2000\_OFFSET\_16**

**BFD\_RELOC\_IQ2000\_OFFSET\_21**

**BFD\_RELOC\_IQ2000\_UHI16**

IQ2000 Relocations.

**BFD\_RELOC\_XTENSA\_RTLD**

Special Xtensa relocation used only by PLT entries in ELF shared objects to indicate that the runtime linker should set the value to one of its own internal functions or data structures.

**BFD\_RELOC\_XTENSA\_GLOB\_DAT**

**BFD\_RELOC\_XTENSA\_JMP\_SLOT**

**BFD\_RELOC\_XTENSA\_RELATIVE**

Xtensa relocations for ELF shared objects.

**BFD\_RELOC\_XTENSA\_PLT**

Xtensa relocation used in ELF object files for symbols that may require PLT entries. Otherwise, this is just a generic 32-bit relocation.

**BFD\_RELOC\_XTENSA\_DIFF8**

**BFD\_RELOC\_XTENSA\_DIFF16**

**BFD\_RELOC\_XTENSA\_DIFF32**

Xtensa relocations to mark the difference of two local symbols. These are only needed to support linker relaxation and can be ignored when not relaxing. The field is set to the value of the difference assuming no relaxation. The relocation encodes the position of the first symbol so the linker can determine whether to adjust the field value.

**BFD\_RELOC\_XTENSA\_SLOT0\_OP**

**BFD\_RELOC\_XTENSA\_SLOT1\_OP**

**BFD\_RELOC\_XTENSA\_SLOT2\_OP**

**BFD\_RELOC\_XTENSA\_SLOT3\_OP**

**BFD\_RELOC\_XTENSA\_SLOT4\_OP**

**BFD\_RELOC\_XTENSA\_SLOT5\_OP**

**BFD\_RELOC\_XTENSA\_SLOT6\_OP**

BFD\_RELOC\_XTENSA\_SLOT7\_OP  
 BFD\_RELOC\_XTENSA\_SLOT8\_OP  
 BFD\_RELOC\_XTENSA\_SLOT9\_OP  
 BFD\_RELOC\_XTENSA\_SLOT10\_OP  
 BFD\_RELOC\_XTENSA\_SLOT11\_OP  
 BFD\_RELOC\_XTENSA\_SLOT12\_OP  
 BFD\_RELOC\_XTENSA\_SLOT13\_OP  
 BFD\_RELOC\_XTENSA\_SLOT14\_OP

Generic Xtensa relocations for instruction operands. Only the slot number is encoded in the relocation. The relocation applies to the last PC-relative immediate operand, or if there are no PC-relative immediates, to the last immediate operand.

BFD\_RELOC\_XTENSA\_SLOT0\_ALT  
 BFD\_RELOC\_XTENSA\_SLOT1\_ALT  
 BFD\_RELOC\_XTENSA\_SLOT2\_ALT  
 BFD\_RELOC\_XTENSA\_SLOT3\_ALT  
 BFD\_RELOC\_XTENSA\_SLOT4\_ALT  
 BFD\_RELOC\_XTENSA\_SLOT5\_ALT  
 BFD\_RELOC\_XTENSA\_SLOT6\_ALT  
 BFD\_RELOC\_XTENSA\_SLOT7\_ALT  
 BFD\_RELOC\_XTENSA\_SLOT8\_ALT  
 BFD\_RELOC\_XTENSA\_SLOT9\_ALT  
 BFD\_RELOC\_XTENSA\_SLOT10\_ALT  
 BFD\_RELOC\_XTENSA\_SLOT11\_ALT  
 BFD\_RELOC\_XTENSA\_SLOT12\_ALT  
 BFD\_RELOC\_XTENSA\_SLOT13\_ALT  
 BFD\_RELOC\_XTENSA\_SLOT14\_ALT

Alternate Xtensa relocations. Only the slot is encoded in the relocation. The meaning of these relocations is opcode-specific.

BFD\_RELOC\_XTENSA\_OP0  
 BFD\_RELOC\_XTENSA\_OP1  
 BFD\_RELOC\_XTENSA\_OP2

Xtensa relocations for backward compatibility. These have all been replaced by BFD\_RELOC\_XTENSA\_SLOT0\_OP.

BFD\_RELOC\_XTENSA\_ASM\_EXPAND

Xtensa relocation to mark that the assembler expanded the instructions from an original target. The expansion size is encoded in the reloc size.

BFD\_RELOC\_XTENSA\_ASM\_SIMPLIFY

Xtensa relocation to mark that the linker should simplify assembler-expanded instructions. This is commonly used internally by the linker after analysis of a BFD\_RELOC\_XTENSA\_ASM\_EXPAND.

BFD\_RELOC\_XTENSA\_TLSDESC\_FN  
 BFD\_RELOC\_XTENSA\_TLSDESC\_ARG  
 BFD\_RELOC\_XTENSA\_TLS\_DTPOFF  
 BFD\_RELOC\_XTENSA\_TLS\_TPOFF

BFD\_RELOC\_XTENSA\_TLS\_FUNC  
 BFD\_RELOC\_XTENSA\_TLS\_ARG  
 BFD\_RELOC\_XTENSA\_TLS\_CALL

Xtensa TLS relocations.

BFD\_RELOC\_Z80\_DISP8  
 8 bit signed offset in (ix+d) or (iy+d).

BFD\_RELOC\_Z8K\_DISP7  
 DJNZ offset.

BFD\_RELOC\_Z8K\_CALLR  
 CALR offset.

BFD\_RELOC\_Z8K\_IMM4L  
 4 bit value.

BFD\_RELOC\_LM32\_CALL  
 BFD\_RELOC\_LM32\_BRANCH  
 BFD\_RELOC\_LM32\_16\_GOT  
 BFD\_RELOC\_LM32\_GOTOFF\_HI16  
 BFD\_RELOC\_LM32\_GOTOFF\_LO16  
 BFD\_RELOC\_LM32\_COPY  
 BFD\_RELOC\_LM32\_GLOB\_DAT  
 BFD\_RELOC\_LM32\_JMP\_SLOT  
 BFD\_RELOC\_LM32\_RELATIVE  
 Lattice Mico32 relocations.

BFD\_RELOC\_MACH\_O\_SECTDIFF  
 Difference between two section addresses. Must be followed by a  
 BFD\_RELOC\_MACH\_O\_PAIR.

BFD\_RELOC\_MACH\_O\_PAIR  
 Pair of relocation. Contains the first symbol.

BFD\_RELOC\_MACH\_O\_X86\_64\_BRANCH32  
 BFD\_RELOC\_MACH\_O\_X86\_64\_BRANCH8  
 PCREL relocations. They are marked as branch to create PLT entry if required.

BFD\_RELOC\_MACH\_O\_X86\_64\_GOT  
 Used when referencing a GOT entry.

BFD\_RELOC\_MACH\_O\_X86\_64\_GOT\_LOAD  
 Used when loading a GOT entry with movq. It is specially marked so that the linker  
 could optimize the movq to a leaq if possible.

BFD\_RELOC\_MACH\_O\_X86\_64\_SUBTRACTOR32  
 Symbol will be subtracted. Must be followed by a BFD\_RELOC\_64.

BFD\_RELOC\_MACH\_O\_X86\_64\_SUBTRACTOR64  
 Symbol will be subtracted. Must be followed by a BFD\_RELOC\_64.

**BFD\_RELOC\_MACH\_O\_X86\_64\_PCREL32\_1**

Same as BFD\_RELOC\_32\_PCREL but with an implicit -1 addend.

**BFD\_RELOC\_MACH\_O\_X86\_64\_PCREL32\_2**

Same as BFD\_RELOC\_32\_PCREL but with an implicit -2 addend.

**BFD\_RELOC\_MACH\_O\_X86\_64\_PCREL32\_4**

Same as BFD\_RELOC\_32\_PCREL but with an implicit -4 addend.

**BFD\_RELOC\_MICROBLAZE\_32\_LO**

This is a 32 bit reloc for the microblaze that stores the low 16 bits of a value

**BFD\_RELOC\_MICROBLAZE\_32\_LO\_PCREL**

This is a 32 bit pc-relative reloc for the microblaze that stores the low 16 bits of a value

**BFD\_RELOC\_MICROBLAZE\_32\_ROSDA**

This is a 32 bit reloc for the microblaze that stores a value relative to the read-only small data area anchor

**BFD\_RELOC\_MICROBLAZE\_32\_RWSDA**

This is a 32 bit reloc for the microblaze that stores a value relative to the read-write small data area anchor

**BFD\_RELOC\_MICROBLAZE\_32\_SYM\_OP\_SYM**

This is a 32 bit reloc for the microblaze to handle expressions of the form "Symbol Op Symbol"

**BFD\_RELOC\_MICROBLAZE\_64\_NONE**

This is a 64 bit reloc that stores the 32 bit pc relative value in two words (with an imm instruction). No relocation is done here - only used for relaxing

**BFD\_RELOC\_MICROBLAZE\_64\_GOTPC**

This is a 64 bit reloc that stores the 32 bit pc relative value in two words (with an imm instruction). The relocation is PC-relative GOT offset

**BFD\_RELOC\_MICROBLAZE\_64\_GOT**

This is a 64 bit reloc that stores the 32 bit pc relative value in two words (with an imm instruction). The relocation is GOT offset

**BFD\_RELOC\_MICROBLAZE\_64\_PLT**

This is a 64 bit reloc that stores the 32 bit pc relative value in two words (with an imm instruction). The relocation is PC-relative offset into PLT

**BFD\_RELOC\_MICROBLAZE\_64\_GOTOFF**

This is a 64 bit reloc that stores the 32 bit GOT relative value in two words (with an imm instruction). The relocation is relative offset from \_GLOBAL\_OFFSET\_TABLE\_

**BFD\_RELOC\_MICROBLAZE\_32\_GOTOFF**

This is a 32 bit reloc that stores the 32 bit GOT relative value in a word. The relocation is relative offset from



**BFD\_RELOC\_MICROBLAZE\_COPY**

This is used to tell the dynamic linker to copy the value out of the dynamic object into the runtime process image.

```
typedef enum bfd_reloc_code_real bfd_reloc_code_real_type;
```

**2.10.2.2 bfd\_reloc\_type\_lookup****Synopsis**

```
reloc_howto_type *bfd_reloc_type_lookup
(bfd *abfd, bfd_reloc_code_real_type code);
reloc_howto_type *bfd_reloc_name_lookup
(bfd *abfd, const char *reloc_name);
```

**Description**

Return a pointer to a howto structure which, when invoked, will perform the relocation *code* on data from the architecture noted.

**2.10.2.3 bfd\_default\_reloc\_type\_lookup****Synopsis**

```
reloc_howto_type *bfd_default_reloc_type_lookup
(bfd *abfd, bfd_reloc_code_real_type code);
```

**Description**

Provides a default relocation lookup routine for any architecture.

**2.10.2.4 bfd\_get\_reloc\_code\_name****Synopsis**

```
const char *bfd_get_reloc_code_name (bfd_reloc_code_real_type code);
```

**Description**

Provides a printable name for the supplied relocation code. Useful mainly for printing error messages.

**2.10.2.5 bfd\_generic\_relax\_section****Synopsis**

```
bfd_boolean bfd_generic_relax_section
(bfd *abfd,
 asection *section,
 struct bfd_link_info *,
 bfd_boolean *);
```

**Description**

Provides default handling for relaxing for back ends which don't do relaxing.

**2.10.2.6 bfd\_generic\_gc\_sections****Synopsis**

```
bfd_boolean bfd_generic_gc_sections
(bfd *, struct bfd_link_info *);
```

**Description**

Provides default handling for relaxing for back ends which don't do section gc – i.e., does nothing.

**2.10.2.7 bfd\_generic\_merge\_sections****Synopsis**

```
bfd_boolean bfd_generic_merge_sections
(bfd *, struct bfd_link_info *);
```

**Description**

Provides default handling for SEC\_MERGE section merging for back ends which don't have SEC\_MERGE support – i.e., does nothing.

**2.10.2.8 bfd\_generic\_get\_relocated\_section\_contents****Synopsis**

```
bfd_byte *bfd_generic_get_relocated_section_contents
(bfd *abfd,
 struct bfd_link_info *link_info,
 struct bfd_link_order *link_order,
 bfd_byte *data,
 bfd_boolean relocatable,
 asymbol **symbols);
```

**Description**

Provides default handling of relocation effort for back ends which can't be bothered to do it efficiently.

**2.11 Core files****2.11.1 Core file functions****Description**

These are functions pertaining to core files.

**2.11.1.1 bfd\_core\_file\_failing\_command****Synopsis**

```
const char *bfd_core_file_failing_command (bfd *abfd);
```

**Description**

Return a read-only string explaining which program was running when it failed and produced the core file *abfd*.

**2.11.1.2 bfd\_core\_file\_failing\_signal****Synopsis**

```
int bfd_core_file_failing_signal (bfd *abfd);
```

**Description**

Returns the signal number which caused the core dump which generated the file the BFD *abfd* is attached to.

### 2.11.1.3 bfd\_core\_file\_pid

#### Synopsis

```
int bfd_core_file_pid (bfd *abfd);
```

#### Description

Returns the PID of the process the core dump the BFD *abfd* is attached to was generated from.

### 2.11.1.4 core\_file\_matches\_executable\_p

#### Synopsis

```
bfd_boolean core_file_matches_executable_p
(bfd *core_bfd, bfd *exec_bfd);
```

#### Description

Return TRUE if the core file attached to *core\_bfd* was generated by a run of the executable file attached to *exec\_bfd*, FALSE otherwise.

### 2.11.1.5 generic\_core\_file\_matches\_executable\_p

#### Synopsis

```
bfd_boolean generic_core_file_matches_executable_p
(bfd *core_bfd, bfd *exec_bfd);
```

#### Description

Return TRUE if the core file attached to *core\_bfd* was generated by a run of the executable file attached to *exec\_bfd*. The match is based on executable basenames only.

Note: When not able to determine the core file failing command or the executable name, we still return TRUE even though we're not sure that core file and executable match. This is to avoid generating a false warning in situations where we really don't know whether they match or not.

## 2.12 Targets

### Description

Each port of BFD to a different machine requires the creation of a target back end. All the back end provides to the root part of BFD is a structure containing pointers to functions which perform certain low level operations on files. BFD translates the applications's requests through a pointer into calls to the back end routines.

When a file is opened with `bfd_openr`, its format and target are unknown. BFD uses various mechanisms to determine how to interpret the file. The operations performed are:

- Create a BFD by calling the internal routine `_bfd_new_bfd`, then call `bfd_find_target` with the target string supplied to `bfd_openr` and the new BFD pointer.
- If a null target string was provided to `bfd_find_target`, look up the environment variable `GNUTARGET` and use that as the target string.
- If the target string is still NULL, or the target string is `default`, then use the first item in the target vector as the target type, and set `target_defaulted` in the BFD to cause `bfd_check_format` to loop through all the targets. See [Section 2.12.1 \[bfd\\_target\]](#), [page 108](#). See [Section 2.9 \[Formats\]](#), [page 49](#).

- Otherwise, inspect the elements in the target vector one by one, until a match on target name is found. When found, use it.
- Otherwise return the error `bfd_error_invalid_target` to `bfd_openr`.
- `bfd_openr` attempts to open the file using `bfd_open_file`, and returns the BFD.

Once the BFD has been opened and the target selected, the file format may be determined. This is done by calling `bfd_check_format` on the BFD with a suggested format. If `target_defaulted` has been set, each possible target type is tried to see if it recognizes the specified format. `bfd_check_format` returns TRUE when the caller guesses right.

### 2.12.1 bfd\_target

#### Description

This structure contains everything that BFD knows about a target. It includes things like its byte order, name, and which routines to call to do various operations.

Every BFD points to a target structure with its `xvec` member.

The macros below are used to dispatch to functions through the `bfd_target` vector. They are used in a number of macros further down in 'bfd.h', and are also used when calling various routines by hand inside the BFD implementation. The *arglist* argument must be parenthesized; it contains all the arguments to the called function.

They make the documentation (more) unpleasant to read, so if someone wants to fix this and not break the above, please do.

```
#define BFD_SEND(bfd, message, arglist) \
 ((*((bfd)->xvec->message)) arglist)

#ifdef DEBUG_BFD_SEND
#undef BFD_SEND
#define BFD_SEND(bfd, message, arglist) \
 (((bfd) && (bfd)->xvec && (bfd)->xvec->message) ? \
 ((*((bfd)->xvec->message)) arglist) : \
 (bfd_assert (__FILE__, __LINE__), NULL))
#endif
```

For operations which index on the BFD format:

```
#define BFD_SEND_FMT(bfd, message, arglist) \
 (((bfd)->xvec->message[(int) ((bfd)->format)]) arglist)

#ifdef DEBUG_BFD_SEND
#undef BFD_SEND_FMT
#define BFD_SEND_FMT(bfd, message, arglist) \
 (((bfd) && (bfd)->xvec && (bfd)->xvec->message) ? \
 (((bfd)->xvec->message[(int) ((bfd)->format)]) arglist) : \
 (bfd_assert (__FILE__, __LINE__), NULL))
#endif
```

This is the structure which defines the type of BFD this is. The `xvec` member of the struct `bfd` itself points here. Each module that implements access to a different target under BFD, defines one of these.

FIXME, these names should be rationalised with the names of the entry points which call them. Too bad we can't have one macro to define them both!

```
enum bfd_flavour
{
 bfd_target_unknown_flavour,
 bfd_target_aout_flavour,
 bfd_target_coff_flavour,
 bfd_target_ecoff_flavour,
 bfd_target_xcoff_flavour,
 bfd_target_elf_flavour,
 bfd_target_ieee_flavour,
 bfd_target_nlm_flavour,
 bfd_target_oasys_flavour,
 bfd_target_tekhhex_flavour,
 bfd_target_srec_flavour,
 bfd_target_verilog_flavour,
 bfd_target_ihex_flavour,
 bfd_target_som_flavour,
 bfd_target_os9k_flavour,
 bfd_target_versados_flavour,
 bfd_target_msdos_flavour,
 bfd_target_ovax_flavour,
 bfd_target_evax_flavour,
 bfd_target_mmo_flavour,
 bfd_target_mach_o_flavour,
 bfd_target_pef_flavour,
 bfd_target_pef_xlib_flavour,
 bfd_target_sym_flavour
};

enum bfd_endian { BFD_ENDIAN_BIG, BFD_ENDIAN_LITTLE, BFD_ENDIAN_UNKNOWN };■

/* Forward declaration. */
typedef struct bfd_link_info _bfd_link_info;

typedef struct bfd_target
{
 /* Identifies the kind of target, e.g., SunOS4, Ultrix, etc. */
 char *name;

 /* The "flavour" of a back end is a general indication about
 the contents of a file. */
 enum bfd_flavour flavour;

 /* The order of bytes within the data area of a file. */
 enum bfd_endian byteorder;
};
```

```

/* The order of bytes within the header parts of a file. */
enum bfd_endian header_byteorder;

/* A mask of all the flags which an executable may have set -
 from the set BFD_NO_FLAGS, HAS_RELOC, ...D_PAGED. */
flagword object_flags;

/* A mask of all the flags which a section may have set - from
 the set SEC_NO_FLAGS, SEC_ALLOC, ...SET_NEVER_LOAD. */
flagword section_flags;

/* The character normally found at the front of a symbol.
 (if any), perhaps '_'. */
char symbol_leading_char;

/* The pad character for file names within an archive header. */
char ar_pad_char;

/* The maximum number of characters in an archive header. */
unsigned short ar_max_namelen;

/* Entries for byte swapping for data. These are different from the
 other entry points, since they don't take a BFD as the first argument.
 Certain other handlers could do the same. */
bfd_uint64_t (*bfd_getx64) (const void *);
bfd_int64_t (*bfd_getx_signed_64) (const void *);
void (*bfd_putx64) (bfd_uint64_t, void *);
bfd_vma (*bfd_getx32) (const void *);
bfd_signed_vma (*bfd_getx_signed_32) (const void *);
void (*bfd_putx32) (bfd_vma, void *);
bfd_vma (*bfd_getx16) (const void *);
bfd_signed_vma (*bfd_getx_signed_16) (const void *);
void (*bfd_putx16) (bfd_vma, void *);

/* Byte swapping for the headers. */
bfd_uint64_t (*bfd_h_getx64) (const void *);
bfd_int64_t (*bfd_h_getx_signed_64) (const void *);
void (*bfd_h_putx64) (bfd_uint64_t, void *);
bfd_vma (*bfd_h_getx32) (const void *);
bfd_signed_vma (*bfd_h_getx_signed_32) (const void *);
void (*bfd_h_putx32) (bfd_vma, void *);
bfd_vma (*bfd_h_getx16) (const void *);
bfd_signed_vma (*bfd_h_getx_signed_16) (const void *);
void (*bfd_h_putx16) (bfd_vma, void *);

/* Format dependent routines: these are vectors of entry points

```

```

 within the target vector structure, one for each format to check. */

/* Check the format of a file being read. Return a bfd_target * or zero. */
const struct bfd_target *(*_bfd_check_format[bfd_type_end]) (bfd *);

/* Set the format of a file being written. */
bfd_boolean (*_bfd_set_format[bfd_type_end]) (bfd *);

/* Write cached information into a file being written, at bfd_close. */
bfd_boolean (*_bfd_write_contents[bfd_type_end]) (bfd *);

```

The general target vector. These vectors are initialized using the BFD\_JUMP\_TABLE macros.

```

/* Generic entry points. */
#define BFD_JUMP_TABLE_GENERIC(NAME) \
 NAME##_close_and_cleanup, \
 NAME##_bfd_free_cached_info, \
 NAME##_new_section_hook, \
 NAME##_get_section_contents, \
 NAME##_get_section_contents_in_window

/* Called when the BFD is being closed to do any necessary cleanup. */
bfd_boolean (*_close_and_cleanup) (bfd *);
/* Ask the BFD to free all cached information. */
bfd_boolean (*_bfd_free_cached_info) (bfd *);
/* Called when a new section is created. */
bfd_boolean (*_new_section_hook) (bfd *, sec_ptr);
/* Read the contents of a section. */
bfd_boolean (*_bfd_get_section_contents)
 (bfd *, sec_ptr, void *, file_ptr, bfd_size_type);
bfd_boolean (*_bfd_get_section_contents_in_window)
 (bfd *, sec_ptr, bfd_window *, file_ptr, bfd_size_type);

/* Entry points to copy private data. */
#define BFD_JUMP_TABLE_COPY(NAME) \
 NAME##_bfd_copy_private_bfd_data, \
 NAME##_bfd_merge_private_bfd_data, \
 _bfd_generic_init_private_section_data, \
 NAME##_bfd_copy_private_section_data, \
 NAME##_bfd_copy_private_symbol_data, \
 NAME##_bfd_copy_private_header_data, \
 NAME##_bfd_set_private_flags, \
 NAME##_bfd_print_private_bfd_data

/* Called to copy BFD general private data from one object file

```

```

 to another. */
bfd_boolean (*_bfd_copy_private_bfd_data) (bfd *, bfd *);
/* Called to merge BFD general private data from one object file
 to a common output file when linking. */
bfd_boolean (*_bfd_merge_private_bfd_data) (bfd *, bfd *);
/* Called to initialize BFD private section data from one object file
 to another. */
#define bfd_init_private_section_data(ibfd, isec, obfd, osec, link_info) \
 BFD_SEND (obfd, _bfd_init_private_section_data, (ibfd, isec, obfd, osec, link_info))
bfd_boolean (*_bfd_init_private_section_data)
 (bfd *, sec_ptr, bfd *, sec_ptr, struct bfd_link_info *);
/* Called to copy BFD private section data from one object file
 to another. */
bfd_boolean (*_bfd_copy_private_section_data)
 (bfd *, sec_ptr, bfd *, sec_ptr);
/* Called to copy BFD private symbol data from one symbol
 to another. */
bfd_boolean (*_bfd_copy_private_symbol_data)
 (bfd *, asymbol *, bfd *, asymbol *);
/* Called to copy BFD private header data from one object file
 to another. */
bfd_boolean (*_bfd_copy_private_header_data)
 (bfd *, bfd *);
/* Called to set private backend flags. */
bfd_boolean (*_bfd_set_private_flags) (bfd *, flagword);

/* Called to print private BFD data. */
bfd_boolean (*_bfd_print_private_bfd_data) (bfd *, void *);

/* Core file entry points. */
#define BFD_JUMP_TABLE_CORE(NAME) \
 NAME##_core_file_failing_command, \
 NAME##_core_file_failing_signal, \
 NAME##_core_file_matches_executable_p, \
 NAME##_core_file_pid

char * (*_core_file_failing_command) (bfd *);
int (*_core_file_failing_signal) (bfd *);
bfd_boolean (*_core_file_matches_executable_p) (bfd *, bfd *);
int (*_core_file_pid) (bfd *);

/* Archive entry points. */
#define BFD_JUMP_TABLE_ARCHIVE(NAME) \
 NAME##_slurp_armap, \
 NAME##_slurp_extended_name_table, \
 NAME##_construct_extended_name_table, \
 NAME##_truncate_arname, \

```



```

NAME##_write_armap, \
NAME##_read_ar_hdr, \
NAME##_write_ar_hdr, \
NAME##_openr_next_archived_file, \
NAME##_get_elt_at_index, \
NAME##_generic_stat_arch_elt, \
NAME##_update_armap_timestamp

bfd_boolean (*bfd_slurp_armap) (bfd *);
bfd_boolean (*bfd_slurp_extended_name_table) (bfd *);
bfd_boolean (*bfd_construct_extended_name_table)
 (bfd *, char **, bfd_size_type *, const char **);
void (*bfd_truncate_arname) (bfd *, const char *, char *);
bfd_boolean (*write_armap)
 (bfd *, unsigned int, struct orl *, unsigned int, int);
void * (*bfd_read_ar_hdr_fn) (bfd *);
bfd_boolean (*bfd_write_ar_hdr_fn) (bfd *, bfd *);
bfd * (*openr_next_archived_file) (bfd *, bfd *);
#define bfd_get_elt_at_index(b,i) BFD_SEND (b, _bfd_get_elt_at_index, (b,i))
bfd * (*bfd_get_elt_at_index) (bfd *, symindex);
int (*bfd_stat_arch_elt) (bfd *, struct stat *);
bfd_boolean (*bfd_update_armap_timestamp) (bfd *);

/* Entry points used for symbols. */
#define BFD_JUMP_TABLE_SYMBOLS(NAME) \
NAME##_get_symtab_upper_bound, \
NAME##_canonicalize_symtab, \
NAME##_make_empty_symbol, \
NAME##_print_symbol, \
NAME##_get_symbol_info, \
NAME##_bfd_is_local_label_name, \
NAME##_bfd_is_target_special_symbol, \
NAME##_get_lineno, \
NAME##_find_nearest_line, \
_bfd_generic_find_line, \
NAME##_find_inliner_info, \
NAME##_bfd_make_debug_symbol, \
NAME##_read_minisymbols, \
NAME##_minisymbol_to_symbol

long (*bfd_get_symtab_upper_bound) (bfd *);
long (*bfd_canonicalize_symtab)
 (bfd *, struct bfd_symbol **);
struct bfd_symbol *
 (*bfd_make_empty_symbol) (bfd *);
void (*bfd_print_symbol)
 (bfd *, void *, struct bfd_symbol *, bfd_print_symbol_type);

```

```

#define bfd_print_symbol(b,p,s,e) BFD_SEND (b, _bfd_print_symbol, (b,p,s,e))
void (*_bfd_get_symbol_info)
 (bfd *, struct bfd_symbol *, symbol_info *);
#define bfd_get_symbol_info(b,p,e) BFD_SEND (b, _bfd_get_symbol_info, (b,p,e))
bfd_boolean (*_bfd_is_local_label_name) (bfd *, const char *);
bfd_boolean (*_bfd_is_target_special_symbol) (bfd *, asymbol *);
arelent * (*_get_lineno) (bfd *, struct bfd_symbol *);
bfd_boolean (*_bfd_find_nearest_line)
 (bfd *, struct bfd_section *, struct bfd_symbol **, bfd_vma,
 const char **, const char **, unsigned int *);
bfd_boolean (*_bfd_find_line)
 (bfd *, struct bfd_symbol **, struct bfd_symbol *,
 const char **, unsigned int *);
bfd_boolean (*_bfd_find_inliner_info)
 (bfd *, const char **, const char **, unsigned int *);
/* Back-door to allow format-aware applications to create debug symbols
 while using BFD for everything else. Currently used by the assembler
 when creating COFF files. */
asymbol * (*_bfd_make_debug_symbol)
 (bfd *, void *, unsigned long size);
#define bfd_read_minisymbols(b, d, m, s) \
 BFD_SEND (b, _read_minisymbols, (b, d, m, s))
long (*_read_minisymbols)
 (bfd *, bfd_boolean, void **, unsigned int *);
#define bfd_minisymbol_to_symbol(b, d, m, f) \
 BFD_SEND (b, _minisymbol_to_symbol, (b, d, m, f))
asymbol * (*_minisymbol_to_symbol)
 (bfd *, bfd_boolean, const void *, asymbol *);

/* Routines for relocs. */
#define BFD_JUMP_TABLE_RELOCS(NAME) \
 NAME##_get_reloc_upper_bound, \
 NAME##_canonicalize_reloc, \
 NAME##_bfd_reloc_type_lookup, \
 NAME##_bfd_reloc_name_lookup

long (*_get_reloc_upper_bound) (bfd *, sec_ptr);
long (*_bfd_canonicalize_reloc)
 (bfd *, sec_ptr, arelent **, struct bfd_symbol **);
/* See documentation on reloc types. */
reloc_howto_type *
 (*reloc_type_lookup) (bfd *, bfd_reloc_code_real_type);
reloc_howto_type *
 (*reloc_name_lookup) (bfd *, const char *);

/* Routines used when writing an object file. */

```

```

#define BFD_JUMP_TABLE_WRITE(NAME) \
 NAME##_set_arch_mach, \
 NAME##_set_section_contents

bfd_boolean (*_bfd_set_arch_mach)
 (bfd *, enum bfd_architecture, unsigned long);
bfd_boolean (*_bfd_set_section_contents)
 (bfd *, sec_ptr, const void *, file_ptr, bfd_size_type);

/* Routines used by the linker. */
#define BFD_JUMP_TABLE_LINK(NAME) \
 NAME##_sizeof_headers, \
 NAME##_bfd_get_relocated_section_contents, \
 NAME##_bfd_relax_section, \
 NAME##_bfd_link_hash_table_create, \
 NAME##_bfd_link_hash_table_free, \
 NAME##_bfd_link_add_symbols, \
 NAME##_bfd_link_just_syms, \
 NAME##_bfd_copy_link_hash_symbol_type, \
 NAME##_bfd_final_link, \
 NAME##_bfd_link_split_section, \
 NAME##_bfd_gc_sections, \
 NAME##_bfd_merge_sections, \
 NAME##_bfd_is_group_section, \
 NAME##_bfd_discard_group, \
 NAME##_section_already_linked, \
 NAME##_bfd_define_common_symbol

int (*_bfd_sizeof_headers) (bfd *, struct bfd_link_info *);
bfd_byte * (*_bfd_get_relocated_section_contents)
 (bfd *, struct bfd_link_info *, struct bfd_link_order *,
 bfd_byte *, bfd_boolean, struct bfd_symbol **);

bfd_boolean (*_bfd_relax_section)
 (bfd *, struct bfd_section *, struct bfd_link_info *, bfd_boolean *);

/* Create a hash table for the linker. Different backends store
 different information in this table. */
struct bfd_link_hash_table *
 (*_bfd_link_hash_table_create) (bfd *);

/* Release the memory associated with the linker hash table. */
void (*_bfd_link_hash_table_free) (struct bfd_link_hash_table *);

/* Add symbols from this object file into the hash table. */
bfd_boolean (*_bfd_link_add_symbols) (bfd *, struct bfd_link_info *);

```

```

/* Indicate that we are only retrieving symbol values from this section. */
void (*_bfd_link_just_syms) (asection *, struct bfd_link_info *);

/* Copy the symbol type of a linker hash table entry. */
#define bfd_copy_link_hash_symbol_type(b, t, f) \
 BFD_SEND (b, _bfd_copy_link_hash_symbol_type, (b, t, f))
void (*_bfd_copy_link_hash_symbol_type)
 (bfd *, struct bfd_link_hash_entry *, struct bfd_link_hash_entry *);

/* Do a link based on the link_order structures attached to each
 section of the BFD. */
bfd_boolean (*_bfd_final_link) (bfd *, struct bfd_link_info *);

/* Should this section be split up into smaller pieces during linking. */
bfd_boolean (*_bfd_link_split_section) (bfd *, struct bfd_section *);

/* Remove sections that are not referenced from the output. */
bfd_boolean (*_bfd_gc_sections) (bfd *, struct bfd_link_info *);

/* Attempt to merge SEC_MERGE sections. */
bfd_boolean (*_bfd_merge_sections) (bfd *, struct bfd_link_info *);

/* Is this section a member of a group? */
bfd_boolean (*_bfd_is_group_section) (bfd *, const struct bfd_section *);

/* Discard members of a group. */
bfd_boolean (*_bfd_discard_group) (bfd *, struct bfd_section *);

/* Check if SEC has been already linked during a reloceatable or
 final link. */
void (*_section_already_linked) (bfd *, struct bfd_section *,
 struct bfd_link_info *);

/* Define a common symbol. */
bfd_boolean (*_bfd_define_common_symbol) (bfd *, struct bfd_link_info *,
 struct bfd_link_hash_entry *);

/* Routines to handle dynamic symbols and relocs. */
#define BFD_JUMP_TABLE_DYNAMIC(NAME) \
 NAME##_get_dynamic_symtab_upper_bound, \
 NAME##_canonicalize_dynamic_symtab, \
 NAME##_get_synthetic_symtab, \
 NAME##_get_dynamic_reloc_upper_bound, \
 NAME##_canonicalize_dynamic_reloc

/* Get the amount of memory required to hold the dynamic symbols. */
long (*_bfd_get_dynamic_symtab_upper_bound) (bfd *);

```

```

/* Read in the dynamic symbols. */
long (*_bfd_canonicalize_dynamic_symtab)
 (bfd *, struct bfd_symbol **);
/* Create synthesized symbols. */
long (*_bfd_get_synthetic_symtab)
 (bfd *, long, struct bfd_symbol **, long, struct bfd_symbol **,
 struct bfd_symbol **);
/* Get the amount of memory required to hold the dynamic relocs. */
long (*_bfd_get_dynamic_reloc_upper_bound) (bfd *);
/* Read in the dynamic relocs. */
long (*_bfd_canonicalize_dynamic_reloc)
 (bfd *, arelent **, struct bfd_symbol **);

```

A pointer to an alternative `bfd_target` in case the current one is not satisfactory. This can happen when the target cpu supports both big and little endian code, and target chosen by the linker has the wrong endianness. The function `open_output()` in `ld/ldlang.c` uses this field to find an alternative output format that is suitable.

```

/* Opposite endian version of this target. */
const struct bfd_target * alternative_target;

/* Data for use by back-end routines, which isn't
 generic enough to belong in this structure. */
const void *backend_data;

} bfd_target;

```

### 2.12.1.1 `bfd_set_default_target`

#### Synopsis

```
bfd_boolean bfd_set_default_target (const char *name);
```

#### Description

Set the default target vector to use when recognizing a BFD. This takes the name of the target, which may be a BFD target name or a configuration triplet.

### 2.12.1.2 `bfd_find_target`

#### Synopsis

```
const bfd_target *bfd_find_target (const char *target_name, bfd *abfd);
```

#### Description

Return a pointer to the transfer vector for the object target named *target\_name*. If *target\_name* is NULL, choose the one in the environment variable `GNUTARGET`; if that is null or not defined, then choose the first entry in the target list. Passing in the string "default" or setting the environment variable to "default" will cause the first entry in the target list to be returned, and "target-defaulted" will be set in the BFD if *abfd* isn't NULL. This causes `bfd_check_format` to loop over all the targets to find the one that matches the file being read.

### 2.12.1.3 bfd\_get\_target\_info

#### Synopsis

```
const bfd_target *bfd_get_target_info (const char *target_name,
 bfd *abfd,
 bfd_boolean *is_bigendian,
 int *underscoring,
 const char **def_target_arch);
```

#### Description

Return a pointer to the transfer vector for the object target named *target\_name*. If *target\_name* is NULL, choose the one in the environment variable GNUTARGET; if that is null or not defined, then choose the first entry in the target list. Passing in the string "default" or setting the environment variable to "default" will cause the first entry in the target list to be returned, and "target\_defaulted" will be set in the BFD if *abfd* isn't NULL. This causes *bfd\_check\_format* to loop over all the targets to find the one that matches the file being read. If *is\_bigendian* is not NULL, then set this value to target's endian mode. True for big-endian, FALSE for little-endian or for invalid target. If *underscoring* is not NULL, then set this value to target's underscoring mode. Zero for none-underscoring, -1 for invalid target, else the value of target vector's symbol underscoring. If *def\_target\_arch* is not NULL, then set it to the architecture string specified by the *target\_name*.

### 2.12.1.4 bfd\_target\_list

#### Synopsis

```
const char ** bfd_target_list (void);
```

#### Description

Return a freshly malloced NULL-terminated vector of the names of all the valid BFD targets. Do not modify the names.

### 2.12.1.5 bfd\_seach\_for\_target

#### Synopsis

```
const bfd_target *bfd_search_for_target
(int (*search_func) (const bfd_target *, void *),
 void *);
```

#### Description

Return a pointer to the first transfer vector in the list of transfer vectors maintained by BFD that produces a non-zero result when passed to the function *search\_func*. The parameter *data* is passed, unexamined, to the search function.

## 2.13 Architectures

BFD keeps one atom in a BFD describing the architecture of the data attached to the BFD: a pointer to a *bfd\_arch\_info\_type*.

Pointers to structures can be requested independently of a BFD so that an architecture's information can be interrogated without access to an open BFD.

The architecture information is provided by each architecture package. The set of default architectures is selected by the macro *SELECT\_ARCHITECTURES*. This is normally set up

in the ‘`config/target.mt`’ file of your choice. If the name is not defined, then all the architectures supported are included.

When BFD starts up, all the architectures are called with an initialize method. It is up to the architecture back end to insert as many items into the list of architectures as it wants to; generally this would be one for each machine and one for the default case (an item with a machine field of 0).

BFD’s idea of an architecture is implemented in ‘`archures.c`’.

### 2.13.1 bfd\_architecture

#### Description

This enum gives the object file’s CPU architecture, in a global sense—i.e., what processor family does it belong to? Another field indicates which processor within the family is in use. The machine gives a number which distinguishes different versions of the architecture, containing, for example, 2 and 3 for Intel i960 KA and i960 KB, and 68020 and 68030 for Motorola 68020 and 68030.

```
enum bfd_architecture
{
 bfd_arch_unknown, /* File arch not known. */
 bfd_arch_obscure, /* Arch known, not one of these. */
 bfd_arch_m68k, /* Motorola 68xxx */
#define bfd_mach_m68000 1
#define bfd_mach_m68008 2
#define bfd_mach_m68010 3
#define bfd_mach_m68020 4
#define bfd_mach_m68030 5
#define bfd_mach_m68040 6
#define bfd_mach_m68060 7
#define bfd_mach_cpu32 8
#define bfd_mach_fido 9
#define bfd_mach_mcf_isa_a_nodiv 10
#define bfd_mach_mcf_isa_a 11
#define bfd_mach_mcf_isa_a_mac 12
#define bfd_mach_mcf_isa_a_emac 13
#define bfd_mach_mcf_isa_a_plus 14
#define bfd_mach_mcf_isa_a_plus_mac 15
#define bfd_mach_mcf_isa_a_plus_emac 16
#define bfd_mach_mcf_isa_b_nousp 17
#define bfd_mach_mcf_isa_b_nousp_mac 18
#define bfd_mach_mcf_isa_b_nousp_emac 19
#define bfd_mach_mcf_isa_b 20
#define bfd_mach_mcf_isa_b_mac 21
#define bfd_mach_mcf_isa_b_emac 22
#define bfd_mach_mcf_isa_b_float 23
#define bfd_mach_mcf_isa_b_float_mac 24
#define bfd_mach_mcf_isa_b_float_emac 25
#define bfd_mach_mcf_isa_c 26
```

```

#define bfd_mach_mcf_isa_c_mac 27
#define bfd_mach_mcf_isa_c_emac 28
#define bfd_mach_mcf_isa_c_nodiv 29
#define bfd_mach_mcf_isa_c_nodiv_mac 30
#define bfd_mach_mcf_isa_c_nodiv_emac 31
 bfd_arch_vax, /* DEC Vax */
 bfd_arch_i960, /* Intel 960 */
 /* The order of the following is important.
 lower number indicates a machine type that
 only accepts a subset of the instructions
 available to machines with higher numbers.
 The exception is the "ca", which is
 incompatible with all other machines except
 "core". */

#define bfd_mach_i960_core 1
#define bfd_mach_i960_ka_sa 2
#define bfd_mach_i960_kb_sb 3
#define bfd_mach_i960_mc 4
#define bfd_mach_i960_xa 5
#define bfd_mach_i960_ca 6
#define bfd_mach_i960_jx 7
#define bfd_mach_i960_hx 8

 bfd_arch_or32, /* OpenRISC 32 */

 bfd_arch_sparc, /* SPARC */
#define bfd_mach_sparc 1
/* The difference between v8plus and v9 is that v9 is a true 64 bit env. */
#define bfd_mach_sparc_sparclet 2
#define bfd_mach_sparc_sparclite 3
#define bfd_mach_sparc_v8plus 4
#define bfd_mach_sparc_v8plusa 5 /* with ultrasparc add'ns. */
#define bfd_mach_sparc_sparclite_le 6
#define bfd_mach_sparc_v9 7
#define bfd_mach_sparc_v9a 8 /* with ultrasparc add'ns. */
#define bfd_mach_sparc_v8plusb 9 /* with cheetah add'ns. */
#define bfd_mach_sparc_v9b 10 /* with cheetah add'ns. */
/* Nonzero if MACH has the v9 instruction set. */
#define bfd_mach_sparc_v9_p(mach) \
 ((mach) >= bfd_mach_sparc_v8plus && (mach) <= bfd_mach_sparc_v9b \
 && (mach) != bfd_mach_sparc_sparclite_le)
/* Nonzero if MACH is a 64 bit sparc architecture. */
#define bfd_mach_sparc_64bit_p(mach) \
 ((mach) >= bfd_mach_sparc_v9 && (mach) != bfd_mach_sparc_v8plusb)
 bfd_arch_spu, /* PowerPC SPU */
#define bfd_mach_spu 256

```



```

 bfd_arch_mips, /* MIPS Rxxxx */
#define bfd_mach_mips3000 3000
#define bfd_mach_mips3900 3900
#define bfd_mach_mips4000 4000
#define bfd_mach_mips4010 4010
#define bfd_mach_mips4100 4100
#define bfd_mach_mips4111 4111
#define bfd_mach_mips4120 4120
#define bfd_mach_mips4300 4300
#define bfd_mach_mips4400 4400
#define bfd_mach_mips4600 4600
#define bfd_mach_mips4650 4650
#define bfd_mach_mips5000 5000
#define bfd_mach_mips5400 5400
#define bfd_mach_mips5500 5500
#define bfd_mach_mips6000 6000
#define bfd_mach_mips7000 7000
#define bfd_mach_mips8000 8000
#define bfd_mach_mips9000 9000
#define bfd_mach_mips10000 10000
#define bfd_mach_mips12000 12000
#define bfd_mach_mips14000 14000
#define bfd_mach_mips16000 16000
#define bfd_mach_mips16 16
#define bfd_mach_mips5 5
#define bfd_mach_mips_loongson_2e 3001
#define bfd_mach_mips_loongson_2f 3002
#define bfd_mach_mips_loongson_3a 3003
#define bfd_mach_mips_sb1 12310201 /* octal 'SB', 01 */
#define bfd_mach_mips_octeon 6501
#define bfd_mach_mips_xlr 887682 /* decimal 'XLR' */
#define bfd_mach_mipsisa32 32
#define bfd_mach_mipsisa32r2 33
#define bfd_mach_mipsisa64 64
#define bfd_mach_mipsisa64r2 65
 bfd_arch_i386, /* Intel 386 */
#define bfd_mach_i386_i386 1
#define bfd_mach_i386_i8086 2
#define bfd_mach_i386_i386_intel_syntax 3
#define bfd_mach_x64_32 32
#define bfd_mach_x64_32_intel_syntax 33
#define bfd_mach_x86_64 64
#define bfd_mach_x86_64_intel_syntax 65
 bfd_arch_l10m, /* Intel L10M */
#define bfd_mach_l10m 66
#define bfd_mach_l10m_intel_syntax 67
 bfd_arch_we32k, /* AT&T WE32xxx */

```

```

bfd_arch_tahoe, /* CCI/Harris Tahoe */
bfd_arch_i860, /* Intel 860 */
bfd_arch_i370, /* IBM 360/370 Mainframes */
bfd_arch_romp, /* IBM ROMP PC/RT */
bfd_arch_convex, /* Convex */
bfd_arch_m88k, /* Motorola 88xxx */
bfd_arch_m98k, /* Motorola 98xxx */
bfd_arch_pyramid, /* Pyramid Technology */
bfd_arch_h8300, /* Renesas H8/300 (formerly Hitachi H8/300) */
#define bfd_mach_h8300 1
#define bfd_mach_h8300h 2
#define bfd_mach_h8300s 3
#define bfd_mach_h8300hn 4
#define bfd_mach_h8300sn 5
#define bfd_mach_h8300sx 6
#define bfd_mach_h8300sxn 7
bfd_arch_pdp11, /* DEC PDP-11 */
bfd_arch_plugin,
bfd_arch_powerpc, /* PowerPC */
#define bfd_mach_ppc 32
#define bfd_mach_ppc64 64
#define bfd_mach_ppc_403 403
#define bfd_mach_ppc_403gc 4030
#define bfd_mach_ppc_405 405
#define bfd_mach_ppc_505 505
#define bfd_mach_ppc_601 601
#define bfd_mach_ppc_602 602
#define bfd_mach_ppc_603 603
#define bfd_mach_ppc_ec603e 6031
#define bfd_mach_ppc_604 604
#define bfd_mach_ppc_620 620
#define bfd_mach_ppc_630 630
#define bfd_mach_ppc_750 750
#define bfd_mach_ppc_860 860
#define bfd_mach_ppc_a35 35
#define bfd_mach_ppc_rs64ii 642
#define bfd_mach_ppc_rs64iii 643
#define bfd_mach_ppc_7400 7400
#define bfd_mach_ppc_e500 500
#define bfd_mach_ppc_e500mc 5001
#define bfd_mach_ppc_e500mc64 5005
#define bfd_mach_ppc_titan 83
bfd_arch_rs6000, /* IBM RS/6000 */
#define bfd_mach_rs6k 6000
#define bfd_mach_rs6k_rs1 6001
#define bfd_mach_rs6k_rsc 6003
#define bfd_mach_rs6k_rs2 6002

```

```

 bfd_arch_hppa, /* HP PA RISC */
#define bfd_mach_hppa10 10
#define bfd_mach_hppa11 11
#define bfd_mach_hppa20 20
#define bfd_mach_hppa20w 25
 bfd_arch_d10v, /* Mitsubishi D10V */
#define bfd_mach_d10v 1
#define bfd_mach_d10v_ts2 2
#define bfd_mach_d10v_ts3 3
 bfd_arch_d30v, /* Mitsubishi D30V */
 bfd_arch_dlx, /* DLX */
 bfd_arch_m68hc11, /* Motorola 68HC11 */
 bfd_arch_m68hc12, /* Motorola 68HC12 */
#define bfd_mach_m6812_default 0
#define bfd_mach_m6812 1
#define bfd_mach_m6812s 2
 bfd_arch_z8k, /* Zilog Z8000 */
#define bfd_mach_z8001 1
#define bfd_mach_z8002 2
 bfd_arch_h8500, /* Renesas H8/500 (formerly Hitachi H8/500) */
 bfd_arch_sh, /* Renesas / SuperH SH (formerly Hitachi SH) */
#define bfd_mach_sh 1
#define bfd_mach_sh2 0x20
#define bfd_mach_sh_dsp 0x2d
#define bfd_mach_sh2a 0x2a
#define bfd_mach_sh2a_nofpu 0x2b
#define bfd_mach_sh2a_nofpu_or_sh4_nommu_nofpu 0x2a1
#define bfd_mach_sh2a_nofpu_or_sh3_nommu 0x2a2
#define bfd_mach_sh2a_or_sh4 0x2a3
#define bfd_mach_sh2a_or_sh3e 0x2a4
#define bfd_mach_sh2e 0x2e
#define bfd_mach_sh3 0x30
#define bfd_mach_sh3_nommu 0x31
#define bfd_mach_sh3_dsp 0x3d
#define bfd_mach_sh3e 0x3e
#define bfd_mach_sh4 0x40
#define bfd_mach_sh4_nofpu 0x41
#define bfd_mach_sh4_nommu_nofpu 0x42
#define bfd_mach_sh4a 0x4a
#define bfd_mach_sh4a_nofpu 0x4b
#define bfd_mach_sh4a1_dsp 0x4d
#define bfd_mach_sh5 0x50
 bfd_arch_alpha, /* Dec Alpha */
#define bfd_mach_alpha_ev4 0x10
#define bfd_mach_alpha_ev5 0x20
#define bfd_mach_alpha_ev6 0x30
 bfd_arch_arm, /* Advanced Risc Machines ARM. */

```

```

#define bfd_mach_arm_unknown 0
#define bfd_mach_arm_2 1
#define bfd_mach_arm_2a 2
#define bfd_mach_arm_3 3
#define bfd_mach_arm_3M 4
#define bfd_mach_arm_4 5
#define bfd_mach_arm_4T 6
#define bfd_mach_arm_5 7
#define bfd_mach_arm_5T 8
#define bfd_mach_arm_5TE 9
#define bfd_mach_arm_XScale 10
#define bfd_mach_arm_ep9312 11
#define bfd_mach_arm_iWMMXt 12
#define bfd_mach_arm_iWMMXt2 13
 bfd_arch_ns32k, /* National Semiconductors ns32000 */
 bfd_arch_w65, /* WDC 65816 */
 bfd_arch_tic30, /* Texas Instruments TMS320C30 */
 bfd_arch_tic4x, /* Texas Instruments TMS320C3X/4X */
#define bfd_mach_tic3x 30
#define bfd_mach_tic4x 40
 bfd_arch_tic54x, /* Texas Instruments TMS320C54X */
 bfd_arch_tic6x, /* Texas Instruments TMS320C6X */
 bfd_arch_tic80, /* TI TMS320c80 (MVP) */
 bfd_arch_v850, /* NEC V850 */
#define bfd_mach_v850 1
#define bfd_mach_v850e 'E'
#define bfd_mach_v850e1 '1'
#define bfd_mach_v850e2 0x4532
#define bfd_mach_v850e2v3 0x45325633
 bfd_arch_arc, /* ARC Cores */
#define bfd_mach_arc_5 5
#define bfd_mach_arc_6 6
#define bfd_mach_arc_7 7
#define bfd_mach_arc_8 8
 bfd_arch_m32c, /* Renesas M16C/M32C. */
#define bfd_mach_m16c 0x75
#define bfd_mach_m32c 0x78
 bfd_arch_m32r, /* Renesas M32R (formerly Mitsubishi M32R/D) */
#define bfd_mach_m32r 1 /* For backwards compatibility. */
#define bfd_mach_m32rx 'x'
#define bfd_mach_m32r2 '2'
 bfd_arch_mn10200, /* Matsushita MN10200 */
 bfd_arch_mn10300, /* Matsushita MN10300 */
#define bfd_mach_mn10300 300
#define bfd_mach_am33 330
#define bfd_mach_am33_2 332
 bfd_arch_fr30,

```

```

#define bfd_mach_fr30 0x46523330
 bfd_arch_frv,
#define bfd_mach_frv 1
#define bfd_mach_frvsimple 2
#define bfd_mach_fr300 300
#define bfd_mach_fr400 400
#define bfd_mach_fr450 450
#define bfd_mach_frvtomcat 499 /* fr500 prototype */
#define bfd_mach_fr500 500
#define bfd_mach_fr550 550
 bfd_arch_moxie, /* The moxie processor */
#define bfd_mach_moxie 1
 bfd_arch_mcore,
 bfd_arch_mep,
#define bfd_mach_mep 1
#define bfd_mach_mep_h1 0x6831
#define bfd_mach_mep_c5 0x6335
 bfd_arch_ia64, /* HP/Intel ia64 */
#define bfd_mach_ia64_elf64 64
#define bfd_mach_ia64_elf32 32
 bfd_arch_ip2k, /* Ubicom IP2K microcontrollers. */
#define bfd_mach_ip2022 1
#define bfd_mach_ip2022ext 2
 bfd_arch_iq2000, /* Vitesse IQ2000. */
#define bfd_mach_iq2000 1
#define bfd_mach_iq10 2
 bfd_arch_mt,
#define bfd_mach_ms1 1
#define bfd_mach_mrisc2 2
#define bfd_mach_ms2 3
 bfd_arch_pj,
 bfd_arch_avr, /* Atmel AVR microcontrollers. */
#define bfd_mach_avr1 1
#define bfd_mach_avr2 2
#define bfd_mach_avr25 25
#define bfd_mach_avr3 3
#define bfd_mach_avr31 31
#define bfd_mach_avr35 35
#define bfd_mach_avr4 4
#define bfd_mach_avr5 5
#define bfd_mach_avr51 51
#define bfd_mach_avr6 6
#define bfd_mach_avrxmega1 101
#define bfd_mach_avrxmega2 102
#define bfd_mach_avrxmega3 103
#define bfd_mach_avrxmega4 104
#define bfd_mach_avrxmega5 105

```

```

#define bfd_mach_avrxmega6 106
#define bfd_mach_avrxmega7 107
 bfd_arch_bfin, /* ADI Blackfin */
#define bfd_mach_bfin 1
 bfd_arch_cr16, /* National Semiconductor CompactRISC (ie CR16). */
#define bfd_mach_cr16 1
 bfd_arch_cr16c, /* National Semiconductor CompactRISC. */
#define bfd_mach_cr16c 1
 bfd_arch_crx, /* National Semiconductor CRX. */
#define bfd_mach_crx 1
 bfd_arch_cris, /* Axis CRIS */
#define bfd_mach_cris_v0_v10 255
#define bfd_mach_cris_v32 32
#define bfd_mach_cris_v10_v32 1032
 bfd_arch_rx, /* Renesas RX. */
#define bfd_mach_rx 0x75
 bfd_arch_s390, /* IBM s390 */
#define bfd_mach_s390_31 31
#define bfd_mach_s390_64 64
 bfd_arch_score, /* Sunplus score */
#define bfd_mach_score3 3
#define bfd_mach_score7 7
 bfd_arch_openrisc, /* OpenRISC */
 bfd_arch_mmix, /* Donald Knuth's educational processor. */
 bfd_arch_xstormy16,
#define bfd_mach_xstormy16 1
 bfd_arch_msp430, /* Texas Instruments MSP430 architecture. */
#define bfd_mach_msp11 11
#define bfd_mach_msp110 110
#define bfd_mach_msp12 12
#define bfd_mach_msp13 13
#define bfd_mach_msp14 14
#define bfd_mach_msp15 15
#define bfd_mach_msp16 16
#define bfd_mach_msp21 21
#define bfd_mach_msp31 31
#define bfd_mach_msp32 32
#define bfd_mach_msp33 33
#define bfd_mach_msp41 41
#define bfd_mach_msp42 42
#define bfd_mach_msp43 43
#define bfd_mach_msp44 44
 bfd_arch_xc16x, /* Infineon's XC16X Series. */
#define bfd_mach_xc16x 1
#define bfd_mach_xc16xl 2
#define bfd_mach_xc16xs 3
 bfd_arch_xtensa, /* Tensilica's Xtensa cores. */

```

```

#define bfd_mach_xtensa 1
 bfd_arch_z80,
#define bfd_mach_z80strict 1 /* No undocumented opcodes. */
#define bfd_mach_z80 3 /* With ixl, ixh, iyl, and iyh. */
#define bfd_mach_z80full 7 /* All undocumented instructions. */
#define bfd_mach_r800 11 /* R800: successor with multiplication. */
 bfd_arch_lm32, /* Lattice Mico32 */
#define bfd_mach_lm32 1
 bfd_arch_microblaze, /* Xilinx MicroBlaze. */
 bfd_arch_last
};

```

### 2.13.2 bfd\_arch\_info

#### Description

This structure contains information on architectures for use within BFD.

```

typedef struct bfd_arch_info
{
 int bits_per_word;
 int bits_per_address;
 int bits_per_byte;
 enum bfd_architecture arch;
 unsigned long mach;
 const char *arch_name;
 const char *printable_name;
 unsigned int section_align_power;
 /* TRUE if this is the default machine for the architecture.
 The default arch should be the first entry for an arch so that
 all the entries for that arch can be accessed via next. */
 bfd_boolean the_default;
 const struct bfd_arch_info * (*compatible)
 (const struct bfd_arch_info *a, const struct bfd_arch_info *b);

 bfd_boolean (*scan) (const struct bfd_arch_info *, const char *);

 const struct bfd_arch_info *next;
}
bfd_arch_info_type;

```

#### 2.13.2.1 bfd\_printable\_name

##### Synopsis

```
const char *bfd_printable_name (bfd *abfd);
```

##### Description

Return a printable string representing the architecture and machine from the pointer to the architecture info structure.

### 2.13.2.2 bfd\_scan\_arch

#### Synopsis

```
const bfd_arch_info_type *bfd_scan_arch (const char *string);
```

#### Description

Figure out if BFD supports any cpu which could be described with the name *string*. Return a pointer to an `arch_info` structure if a machine is found, otherwise NULL.

### 2.13.2.3 bfd\_arch\_list

#### Synopsis

```
const char **bfd_arch_list (void);
```

#### Description

Return a freshly malloced NULL-terminated vector of the names of all the valid BFD architectures. Do not modify the names.

### 2.13.2.4 bfd\_arch\_get\_compatible

#### Synopsis

```
const bfd_arch_info_type *bfd_arch_get_compatible
(const bfd *abfd, const bfd *bbfd, bfd_boolean accept_unknowns);
```

#### Description

Determine whether two BFDs' architectures and machine types are compatible. Calculates the lowest common denominator between the two architectures and machine types implied by the BFDs and returns a pointer to an `arch_info` structure describing the compatible machine.

### 2.13.2.5 bfd\_default\_arch\_struct

#### Description

The `bfd_default_arch_struct` is an item of `bfd_arch_info_type` which has been initialized to a fairly generic state. A BFD starts life by pointing to this structure, until the correct back end has determined the real architecture of the file.

```
extern const bfd_arch_info_type bfd_default_arch_struct;
```

### 2.13.2.6 bfd\_set\_arch\_info

#### Synopsis

```
void bfd_set_arch_info (bfd *abfd, const bfd_arch_info_type *arg);
```

#### Description

Set the architecture info of *abfd* to *arg*.

### 2.13.2.7 bfd\_default\_set\_arch\_mach

#### Synopsis

```
bfd_boolean bfd_default_set_arch_mach
(bfd *abfd, enum bfd_architecture arch, unsigned long mach);
```

#### Description

Set the architecture and machine type in BFD *abfd* to *arch* and *mach*. Find the correct pointer to a structure and insert it into the `arch_info` pointer.



### 2.13.2.8 bfd\_get\_arch

**Synopsis**

```
enum bfd_architecture bfd_get_arch (bfd *abfd);
```

**Description**

Return the enumerated type which describes the BFD *abfd*'s architecture.

### 2.13.2.9 bfd\_get\_mach

**Synopsis**

```
unsigned long bfd_get_mach (bfd *abfd);
```

**Description**

Return the long type which describes the BFD *abfd*'s machine.

### 2.13.2.10 bfd\_arch\_bits\_per\_byte

**Synopsis**

```
unsigned int bfd_arch_bits_per_byte (bfd *abfd);
```

**Description**

Return the number of bits in one of the BFD *abfd*'s architecture's bytes.

### 2.13.2.11 bfd\_arch\_bits\_per\_address

**Synopsis**

```
unsigned int bfd_arch_bits_per_address (bfd *abfd);
```

**Description**

Return the number of bits in one of the BFD *abfd*'s architecture's addresses.

### 2.13.2.12 bfd\_default\_compatible

**Synopsis**

```
const bfd_arch_info_type *bfd_default_compatible
(const bfd_arch_info_type *a, const bfd_arch_info_type *b);
```

**Description**

The default function for testing for compatibility.

### 2.13.2.13 bfd\_default\_scan

**Synopsis**

```
bfd_boolean bfd_default_scan
(const struct bfd_arch_info *info, const char *string);
```

**Description**

The default function for working out whether this is an architecture hit and a machine hit.

### 2.13.2.14 bfd\_get\_arch\_info

**Synopsis**

```
const bfd_arch_info_type *bfd_get_arch_info (bfd *abfd);
```

**Description**

Return the architecture info struct in *abfd*.

### 2.13.2.15 bfd\_lookup\_arch

#### Synopsis

```
const bfd_arch_info_type *bfd_lookup_arch
 (enum bfd_architecture arch, unsigned long machine);
```

#### Description

Look for the architecture info structure which matches the arguments *arch* and *machine*. A machine of 0 matches the machine/architecture structure which marks itself as the default.

### 2.13.2.16 bfd\_printable\_arch\_mach

#### Synopsis

```
const char *bfd_printable_arch_mach
 (enum bfd_architecture arch, unsigned long machine);
```

#### Description

Return a printable string representing the architecture and machine type.

This routine is depreciated.

### 2.13.2.17 bfd\_octets\_per\_byte

#### Synopsis

```
unsigned int bfd_octets_per_byte (bfd *abfd);
```

#### Description

Return the number of octets (8-bit quantities) per target byte (minimum addressable unit). In most cases, this will be one, but some DSP targets have 16, 32, or even 48 bits per byte.

### 2.13.2.18 bfd\_arch\_mach\_octets\_per\_byte

#### Synopsis

```
unsigned int bfd_arch_mach_octets_per_byte
 (enum bfd_architecture arch, unsigned long machine);
```

#### Description

See `bfd_octets_per_byte`.

This routine is provided for those cases where a `bfd *` is not available

```
/* Set to N to open the next N BFDs using an alternate id space. */
extern unsigned int bfd_use_reserved_id;
```

## 2.14 Opening and closing BFDs

### 2.14.1 Functions for opening and closing

#### 2.14.1.1 bfd\_fopen

#### Synopsis

```
bfd *bfd_fopen (const char *filename, const char *target,
 const char *mode, int fd);
```

#### Description

Open the file *filename* with the target *target*. Return a pointer to the created BFD. If *fd*

is not -1, then `fdopen` is used to open the file; otherwise, `fopen` is used. *mode* is passed directly to `fopen` or `fdopen`.

Calls `bfd_find_target`, so *target* is interpreted as by that function.

The new BFD is marked as cacheable iff *fd* is -1.

If NULL is returned then an error has occurred. Possible errors are `bfd_error_no_memory`, `bfd_error_invalid_target` or `system_call` error.

### 2.14.1.2 `bfd_openr`

#### Synopsis

```
bfd *bfd_openr (const char *filename, const char *target);
```

#### Description

Open the file *filename* (using `fopen`) with the target *target*. Return a pointer to the created BFD.

Calls `bfd_find_target`, so *target* is interpreted as by that function.

If NULL is returned then an error has occurred. Possible errors are `bfd_error_no_memory`, `bfd_error_invalid_target` or `system_call` error.

### 2.14.1.3 `bfd_fdopenr`

#### Synopsis

```
bfd *bfd_fdopenr (const char *filename, const char *target, int fd);
```

#### Description

`bfd_fdopenr` is to `bfd_fopenr` much like `fdopen` is to `fopen`. It opens a BFD on a file already described by the *fd* supplied.

When the file is later `bfd_closed`, the file descriptor will be closed. If the caller desires that this file descriptor be cached by BFD (opened as needed, closed as needed to free descriptors for other opens), with the supplied *fd* used as an initial file descriptor (but subject to closure at any time), call `bfd_set_cacheable(bfd, 1)` on the returned BFD. The default is to assume no caching; the file descriptor will remain open until `bfd_close`, and will not be affected by BFD operations on other files.

Possible errors are `bfd_error_no_memory`, `bfd_error_invalid_target` and `bfd_error_system_call`.

### 2.14.1.4 `bfd_openstreamr`

#### Synopsis

```
bfd *bfd_openstreamr (const char *, const char *, void *);
```

#### Description

Open a BFD for read access on an existing stdio stream. When the BFD is passed to `bfd_close`, the stream will be closed.

### 2.14.1.5 `bfd_openr_iovec`

#### Synopsis

```
bfd *bfd_openr_iovec (const char *filename, const char *target,
void *(*open_func) (struct bfd *nbfd,
```

```

void *open_closure),
void *open_closure,
file_ptr (*pread_func) (struct bfd *nbfd,
void *stream,
void *buf,
file_ptr nbytes,
file_ptr offset),
int (*close_func) (struct bfd *nbfd,
void *stream),
int (*stat_func) (struct bfd *abfd,
void *stream,
struct stat *sb));

```

### Description

Create and return a BFD backed by a read-only *stream*. The *stream* is created using *open\_func*, accessed using *pread\_func* and destroyed using *close\_func*.

Calls *bfd\_find\_target*, so *target* is interpreted as by that function.

Calls *open\_func* (which can call *bfd\_zalloc* and *bfd\_get\_filename*) to obtain the read-only stream backing the BFD. *open\_func* either succeeds returning the non-NULL *stream*, or fails returning NULL (setting *bfd\_error*).

Calls *pread\_func* to request *nbytes* of data from *stream* starting at *offset* (e.g., via a call to *bfd\_read*). *pread\_func* either succeeds returning the number of bytes read (which can be less than *nbytes* when end-of-file), or fails returning -1 (setting *bfd\_error*).

Calls *close\_func* when the BFD is later closed using *bfd\_close*. *close\_func* either succeeds returning 0, or fails returning -1 (setting *bfd\_error*).

Calls *stat\_func* to fill in a stat structure for *bfd\_stat*, *bfd\_get\_size*, and *bfd\_get\_mtime* calls. *stat\_func* returns 0 on success, or returns -1 on failure (setting *bfd\_error*).

If *bfd\_openr\_iovec* returns NULL then an error has occurred. Possible errors are *bfd\_error\_no\_memory*, *bfd\_error\_invalid\_target* and *bfd\_error\_system\_call*.

### 2.14.1.6 bfd\_openw

#### Synopsis

```
bfd *bfd_openw (const char *filename, const char *target);
```

#### Description

Create a BFD, associated with file *filename*, using the file format *target*, and return a pointer to it.

Possible errors are *bfd\_error\_system\_call*, *bfd\_error\_no\_memory*, *bfd\_error\_invalid\_target*.

### 2.14.1.7 bfd\_close

#### Synopsis

```
bfd_boolean bfd_close (bfd *abfd);
```

#### Description

Close a BFD. If the BFD was open for writing, then pending operations are completed and

the file written out and closed. If the created file is executable, then `chmod` is called to mark it as such.

All memory attached to the BFD is released.

The file descriptor associated with the BFD is closed (even if it was passed in to BFD by `bfd_fdopenr`).

**Returns**

TRUE is returned if all is ok, otherwise FALSE.

### 2.14.1.8 `bfd_close_all_done`

**Synopsis**

```
bfd_boolean bfd_close_all_done (bfd *);
```

**Description**

Close a BFD. Differs from `bfd_close` since it does not complete any pending operations. This routine would be used if the application had just used BFD for swapping and didn't want to use any of the writing code.

If the created file is executable, then `chmod` is called to mark it as such.

All memory attached to the BFD is released.

**Returns**

TRUE is returned if all is ok, otherwise FALSE.

### 2.14.1.9 `bfd_create`

**Synopsis**

```
bfd *bfd_create (const char *filename, bfd *templ);
```

**Description**

Create a new BFD in the manner of `bfd_openw`, but without opening a file. The new BFD takes the target from the target used by *templ*. The format is always set to `bfd_object`.

### 2.14.1.10 `bfd_make_writable`

**Synopsis**

```
bfd_boolean bfd_make_writable (bfd *abfd);
```

**Description**

Takes a BFD as created by `bfd_create` and converts it into one like as returned by `bfd_openw`. It does this by converting the BFD to `BFD_IN_MEMORY`. It's assumed that you will call `bfd_make_readable` on this bfd later.

**Returns**

TRUE is returned if all is ok, otherwise FALSE.

### 2.14.1.11 `bfd_make_readable`

**Synopsis**

```
bfd_boolean bfd_make_readable (bfd *abfd);
```

**Description**

Takes a BFD as created by `bfd_create` and `bfd_make_writable` and converts it into one

like as returned by `bfd_openr`. It does this by writing the contents out to the memory buffer, then reversing the direction.

**Returns**

TRUE is returned if all is ok, otherwise FALSE.

**2.14.1.12 bfd\_alloc****Synopsis**

```
void *bfd_alloc (bfd *abfd, bfd_size_type wanted);
```

**Description**

Allocate a block of *wanted* bytes of memory attached to `abfd` and return a pointer to it.

**2.14.1.13 bfd\_alloc2****Synopsis**

```
void *bfd_alloc2 (bfd *abfd, bfd_size_type nmemb, bfd_size_type size);
```

**Description**

Allocate a block of *nmemb* elements of *size* bytes each of memory attached to `abfd` and return a pointer to it.

**2.14.1.14 bfd\_zalloc****Synopsis**

```
void *bfd_zalloc (bfd *abfd, bfd_size_type wanted);
```

**Description**

Allocate a block of *wanted* bytes of zeroed memory attached to `abfd` and return a pointer to it.

**2.14.1.15 bfd\_zalloc2****Synopsis**

```
void *bfd_zalloc2 (bfd *abfd, bfd_size_type nmemb, bfd_size_type size);
```

**Description**

Allocate a block of *nmemb* elements of *size* bytes each of zeroed memory attached to `abfd` and return a pointer to it.

**2.14.1.16 bfd\_calc\_gnu\_debuglink\_crc32****Synopsis**

```
unsigned long bfd_calc_gnu_debuglink_crc32
(unsigned long crc, const unsigned char *buf, bfd_size_type len);
```

**Description**

Computes a CRC value as used in the `.gnu.debuglink` section. Advances the previously computed *crc* value by computing and adding in the *crc32* for *len* bytes of *buf*.

**Returns**

Return the updated CRC32 value.

### 2.14.1.17 get\_debug\_link\_info

#### Synopsis

```
char *get_debug_link_info (bfd *abfd, unsigned long *crc32_out);
```

#### Description

fetch the filename and CRC32 value for any separate debuginfo associated with *abfd*. Return NULL if no such info found, otherwise return filename and update *crc32\_out*.

### 2.14.1.18 separate\_debug\_file\_exists

#### Synopsis

```
bfd_boolean separate_debug_file_exists
(char *name, unsigned long crc32);
```

#### Description

Checks to see if *name* is a file and if its contents match *crc32*.

### 2.14.1.19 find\_separate\_debug\_file

#### Synopsis

```
char *find_separate_debug_file (bfd *abfd);
```

#### Description

Searches *abfd* for a reference to separate debugging information, scans various locations in the filesystem, including the file tree rooted at *debug-file-directory*, and returns a filename of such debugging information if the file is found and has matching CRC32. Returns NULL if no reference to debugging file exists, or file cannot be found.

### 2.14.1.20 bfd\_follow\_gnu\_debuglink

#### Synopsis

```
char *bfd_follow_gnu_debuglink (bfd *abfd, const char *dir);
```

#### Description

Takes a BFD and searches it for a *.gnu.debuglink* section. If this section is found, it examines the section for the name and checksum of a *.debug* file containing auxiliary debugging information. It then searches the filesystem for this *.debug* file in some standard locations, including the directory tree rooted at *dir*, and if found returns the full filename. If *dir* is NULL, it will search a default path configured into libbfd at build time. [XXX this feature is not currently implemented].

#### Returns

NULL on any errors or failure to locate the *.debug* file, otherwise a pointer to a heap-allocated string containing the filename. The caller is responsible for freeing this string.

### 2.14.1.21 bfd\_create\_gnu\_debuglink\_section

#### Synopsis

```
struct bfd_section *bfd_create_gnu_debuglink_section
(bfd *abfd, const char *filename);
```

#### Description

Takes a *BFD* and adds a *.gnu.debuglink* section to it. The section is sized to be big enough to contain a link to the specified *filename*.

**Returns**

A pointer to the new section is returned if all is ok. Otherwise NULL is returned and `bfd_error` is set.

**2.14.1.22 bfd\_fill\_in\_gnu\_debuglink\_section****Synopsis**

```
bfd_boolean bfd_fill_in_gnu_debuglink_section
 (bfd *abfd, struct bfd_section *sect, const char *filename);
```

**Description**

Takes a *BFD* and containing a `.gnu_debuglink` section *SECT* and fills in the contents of the section to contain a link to the specified *filename*. The filename should be relative to the current directory.

**Returns**

TRUE is returned if all is ok. Otherwise FALSE is returned and `bfd_error` is set.

**2.15 Implementation details****2.15.1 Internal functions****Description**

These routines are used within BFD. They are not intended for export, but are documented here for completeness.

**2.15.1.1 bfd\_write\_bigendian\_4byte\_int****Synopsis**

```
bfd_boolean bfd_write_bigendian_4byte_int (bfd *, unsigned int);
```

**Description**

Write a 4 byte integer *i* to the output BFD *abfd*, in big endian order regardless of what else is going on. This is useful in archives.

**2.15.1.2 bfd\_put\_size****2.15.1.3 bfd\_get\_size****Description**

These macros as used for reading and writing raw data in sections; each access (except for bytes) is vectored through the target format of the BFD and mangled accordingly. The mangling performs any necessary endian translations and removes alignment restrictions. Note that types accepted and returned by these macros are identical so they can be swapped around in macros—for example, `libaout.h` defines `GET_WORD` to either `bfd_get_32` or `bfd_get_64`.

In the put routines, *val* must be a `bfd_vma`. If we are on a system without prototypes, the caller is responsible for making sure that is true, with a cast if necessary. We don't cast them in the macro definitions because that would prevent `lint` or `gcc -Wall` from detecting sins such as passing a pointer. To detect calling these with less than a `bfd_vma`, use `gcc -Wconversion` on a host with 64 bit `bfd_vma`'s.



```

/* Byte swapping macros for user section data. */

#define bfd_put_8(abfd, val, ptr) \
 ((void) (*((unsigned char *) (ptr)) = (val) & 0xff))
#define bfd_put_signed_8 \
 bfd_put_8
#define bfd_get_8(abfd, ptr) \
 (*((unsigned char *) (ptr) & 0xff)
#define bfd_get_signed_8(abfd, ptr) \
 (((*(unsigned char *) (ptr) & 0xff) ^ 0x80) - 0x80)

#define bfd_put_16(abfd, val, ptr) \
 BFD_SEND (abfd, bfd_putx16, ((val),(ptr)))
#define bfd_put_signed_16 \
 bfd_put_16
#define bfd_get_16(abfd, ptr) \
 BFD_SEND (abfd, bfd_getx16, (ptr))
#define bfd_get_signed_16(abfd, ptr) \
 BFD_SEND (abfd, bfd_getx_signed_16, (ptr))

#define bfd_put_32(abfd, val, ptr) \
 BFD_SEND (abfd, bfd_putx32, ((val),(ptr)))
#define bfd_put_signed_32 \
 bfd_put_32
#define bfd_get_32(abfd, ptr) \
 BFD_SEND (abfd, bfd_getx32, (ptr))
#define bfd_get_signed_32(abfd, ptr) \
 BFD_SEND (abfd, bfd_getx_signed_32, (ptr))

#define bfd_put_64(abfd, val, ptr) \
 BFD_SEND (abfd, bfd_putx64, ((val), (ptr)))
#define bfd_put_signed_64 \
 bfd_put_64
#define bfd_get_64(abfd, ptr) \
 BFD_SEND (abfd, bfd_getx64, (ptr))
#define bfd_get_signed_64(abfd, ptr) \
 BFD_SEND (abfd, bfd_getx_signed_64, (ptr))

#define bfd_get(bits, abfd, ptr) \
 ((bits) == 8 ? (bfd_vma) bfd_get_8 (abfd, ptr) \
 : (bits) == 16 ? bfd_get_16 (abfd, ptr) \
 : (bits) == 32 ? bfd_get_32 (abfd, ptr) \
 : (bits) == 64 ? bfd_get_64 (abfd, ptr) \
 : (abort (), (bfd_vma) - 1))

#define bfd_put(bits, abfd, val, ptr) \

```

```

((bits) == 8 ? bfd_put_8 (abfd, val, ptr) \
: (bits) == 16 ? bfd_put_16 (abfd, val, ptr) \
: (bits) == 32 ? bfd_put_32 (abfd, val, ptr) \
: (bits) == 64 ? bfd_put_64 (abfd, val, ptr) \
: (abort (), (void) 0))

```

#### 2.15.1.4 bfd\_h\_put\_size

##### Description

These macros have the same function as their `bfd_get_x` brethren, except that they are used for removing information for the header records of object files. Believe it or not, some object files keep their header records in big endian order and their data in little endian order.

```

/* Byte swapping macros for file header data. */

#define bfd_h_put_8(abfd, val, ptr) \
 bfd_put_8 (abfd, val, ptr)
#define bfd_h_put_signed_8(abfd, val, ptr) \
 bfd_put_8 (abfd, val, ptr)
#define bfd_h_get_8(abfd, ptr) \
 bfd_get_8 (abfd, ptr)
#define bfd_h_get_signed_8(abfd, ptr) \
 bfd_get_signed_8 (abfd, ptr)

#define bfd_h_put_16(abfd, val, ptr) \
 BFD_SEND (abfd, bfd_h_putx16, (val, ptr))
#define bfd_h_put_signed_16 \
 bfd_h_put_16
#define bfd_h_get_16(abfd, ptr) \
 BFD_SEND (abfd, bfd_h_getx16, (ptr))
#define bfd_h_get_signed_16(abfd, ptr) \
 BFD_SEND (abfd, bfd_h_getx_signed_16, (ptr))

#define bfd_h_put_32(abfd, val, ptr) \
 BFD_SEND (abfd, bfd_h_putx32, (val, ptr))
#define bfd_h_put_signed_32 \
 bfd_h_put_32
#define bfd_h_get_32(abfd, ptr) \
 BFD_SEND (abfd, bfd_h_getx32, (ptr))
#define bfd_h_get_signed_32(abfd, ptr) \
 BFD_SEND (abfd, bfd_h_getx_signed_32, (ptr))

#define bfd_h_put_64(abfd, val, ptr) \
 BFD_SEND (abfd, bfd_h_putx64, (val, ptr))
#define bfd_h_put_signed_64 \
 bfd_h_put_64

```

```

 bfd_h_put_64
#define bfd_h_get_64(abfd, ptr) \
 BFD_SEND (abfd, bfd_h_getx64, (ptr))
#define bfd_h_get_signed_64(abfd, ptr) \
 BFD_SEND (abfd, bfd_h_getx_signed_64, (ptr))

/* Aliases for the above, which should eventually go away. */

#define H_PUT_64 bfd_h_put_64
#define H_PUT_32 bfd_h_put_32
#define H_PUT_16 bfd_h_put_16
#define H_PUT_8 bfd_h_put_8
#define H_PUT_S64 bfd_h_put_signed_64
#define H_PUT_S32 bfd_h_put_signed_32
#define H_PUT_S16 bfd_h_put_signed_16
#define H_PUT_S8 bfd_h_put_signed_8
#define H_GET_64 bfd_h_get_64
#define H_GET_32 bfd_h_get_32
#define H_GET_16 bfd_h_get_16
#define H_GET_8 bfd_h_get_8
#define H_GET_S64 bfd_h_get_signed_64
#define H_GET_S32 bfd_h_get_signed_32
#define H_GET_S16 bfd_h_get_signed_16
#define H_GET_S8 bfd_h_get_signed_8

```

### 2.15.1.5 bfd\_log2

#### Synopsis

```
unsigned int bfd_log2 (bfd_vma x);
```

#### Description

Return the log base 2 of the value supplied, rounded up. E.g., an `x` of 1025 returns 11. A `x` of 0 returns 0.

## 2.16 File caching

The file caching mechanism is embedded within BFD and allows the application to open as many BFDs as it wants without regard to the underlying operating system's file descriptor limit (often as low as 20 open files). The module in `cache.c` maintains a least recently used list of `BFD_CACHE_MAX_OPEN` files, and exports the name `bfd_cache_lookup`, which runs around and makes sure that the required BFD is open. If not, then it chooses a file to close, closes it and opens the one wanted, returning its file handle.

### 2.16.1 Caching functions

#### 2.16.1.1 bfd\_cache\_init

#### Synopsis

```
bfd_boolean bfd_cache_init (bfd *abfd);
```

**Description**

Add a newly opened BFD to the cache.

**2.16.1.2 bfd\_cache\_close****Synopsis**

```
bfd_boolean bfd_cache_close (bfd *abfd);
```

**Description**

Remove the BFD *abfd* from the cache. If the attached file is open, then close it too.

**Returns**

FALSE is returned if closing the file fails, TRUE is returned if all is well.

**2.16.1.3 bfd\_cache\_close\_all****Synopsis**

```
bfd_boolean bfd_cache_close_all (void);
```

**Description**

Remove all BFDs from the cache. If the attached file is open, then close it too.

**Returns**

FALSE is returned if closing one of the file fails, TRUE is returned if all is well.

**2.16.1.4 bfd\_open\_file****Synopsis**

```
FILE* bfd_open_file (bfd *abfd);
```

**Description**

Call the OS to open a file for *abfd*. Return the `FILE *` (possibly NULL) that results from this operation. Set up the BFD so that future accesses know the file is open. If the `FILE *` returned is NULL, then it won't have been put in the cache, so it won't have to be removed from it.

**2.17 Linker Functions**

The linker uses three special entry points in the BFD target vector. It is not necessary to write special routines for these entry points when creating a new BFD back end, since generic versions are provided. However, writing them can speed up linking and make it use significantly less runtime memory.

The first routine creates a hash table used by the other routines. The second routine adds the symbols from an object file to the hash table. The third routine takes all the object files and links them together to create the output file. These routines are designed so that the linker proper does not need to know anything about the symbols in the object files that it is linking. The linker merely arranges the sections as directed by the linker script and lets BFD handle the details of symbols and relocs.

The second routine and third routines are passed a pointer to a `struct bfd_link_info` structure (defined in `bfdlink.h`) which holds information relevant to the link, including the linker hash table (which was created by the first routine) and a set of callback functions to the linker proper.

The generic linker routines are in `linker.c`, and use the header file `genlink.h`. As of this writing, the only back ends which have implemented versions of these routines are a.out (in `aoutx.h`) and ECOFF (in `ecoff.c`). The a.out routines are used as examples throughout this section.

### 2.17.1 Creating a linker hash table

The linker routines must create a hash table, which must be derived from `struct bfd_link_hash_table` described in `bfdlink.c`. See [Section 2.18 \[Hash Tables\]](#), page 146, for information on how to create a derived hash table. This entry point is called using the target vector of the linker output file.

The `_bfd_link_hash_table_create` entry point must allocate and initialize an instance of the desired hash table. If the back end does not require any additional information to be stored with the entries in the hash table, the entry point may simply create a `struct bfd_link_hash_table`. Most likely, however, some additional information will be needed.

For example, with each entry in the hash table the a.out linker keeps the index the symbol has in the final output file (this index number is used so that when doing a relocatable link the symbol index used in the output file can be quickly filled in when copying over a reloc). The a.out linker code defines the required structures and functions for a hash table derived from `struct bfd_link_hash_table`. The a.out linker hash table is created by the function `NAME(aout,link_hash_table_create)`; it simply allocates space for the hash table, initializes it, and returns a pointer to it.

When writing the linker routines for a new back end, you will generally not know exactly which fields will be required until you have finished. You should simply create a new hash table which defines no additional fields, and then simply add fields as they become necessary.

### 2.17.2 Adding symbols to the hash table

The linker proper will call the `_bfd_link_add_symbols` entry point for each object file or archive which is to be linked (typically these are the files named on the command line, but some may also come from the linker script). The entry point is responsible for examining the file. For an object file, BFD must add any relevant symbol information to the hash table. For an archive, BFD must determine which elements of the archive should be used and adding them to the link.

The a.out version of this entry point is `NAME(aout,link_add_symbols)`.

#### 2.17.2.1 Differing file formats

Normally all the files involved in a link will be of the same format, but it is also possible to link together different format object files, and the back end must support that. The `_bfd_link_add_symbols` entry point is called via the target vector of the file to be added. This has an important consequence: the function may not assume that the hash table is the type created by the corresponding `_bfd_link_hash_table_create` vector. All the `_bfd_link_add_symbols` function can assume about the hash table is that it is derived from `struct bfd_link_hash_table`.

Sometimes the `_bfd_link_add_symbols` function must store some information in the hash table entry to be used by the `_bfd_final_link` function. In such a case the output bfd xvec must be checked to make sure that the hash table was created by an object file of the same format.

The `_bfd_final_link` routine must be prepared to handle a hash entry without any extra information added by the `_bfd_link_add_symbols` function. A hash entry without extra information will also occur when the linker script directs the linker to create a symbol. Note that, regardless of how a hash table entry is added, all the fields will be initialized to some sort of null value by the hash table entry initialization function.

See `ecoff_link_add_externals` for an example of how to check the output bfd before saving information (in this case, the ECOFF external symbol debugging information) in a hash table entry.

### 2.17.2.2 Adding symbols from an object file

When the `_bfd_link_add_symbols` routine is passed an object file, it must add all externally visible symbols in that object file to the hash table. The actual work of adding the symbol to the hash table is normally handled by the function `_bfd_generic_link_add_one_symbol`. The `_bfd_link_add_symbols` routine is responsible for reading all the symbols from the object file and passing the correct information to `_bfd_generic_link_add_one_symbol`.

The `_bfd_link_add_symbols` routine should not use `bfd_canonicalize_symbtab` to read the symbols. The point of providing this routine is to avoid the overhead of converting the symbols into generic `asymbol` structures.

`_bfd_generic_link_add_one_symbol` handles the details of combining common symbols, warning about multiple definitions, and so forth. It takes arguments which describe the symbol to add, notably symbol flags, a section, and an offset. The symbol flags include such things as `BSF_WEAK` or `BSF_INDIRECT`. The section is a section in the object file, or something like `bfd_und_section_ptr` for an undefined symbol or `bfd_com_section_ptr` for a common symbol.

If the `_bfd_final_link` routine is also going to need to read the symbol information, the `_bfd_link_add_symbols` routine should save it somewhere attached to the object file BFD. However, the information should only be saved if the `keep_memory` field of the `info` argument is `TRUE`, so that the `-no-keep-memory` linker switch is effective.

The `a.out` function which adds symbols from an object file is `aout_link_add_object_symbols`, and most of the interesting work is in `aout_link_add_symbols`. The latter saves pointers to the hash tables entries created by `_bfd_generic_link_add_one_symbol` indexed by symbol number, so that the `_bfd_final_link` routine does not have to call the hash table lookup routine to locate the entry.

### 2.17.2.3 Adding symbols from an archive

When the `_bfd_link_add_symbols` routine is passed an archive, it must look through the symbols defined by the archive and decide which elements of the archive should be included in the link. For each such element it must call the `add_archive_element` linker callback, and it must add the symbols from the object file to the linker hash table. (The callback may in fact indicate that a replacement BFD should be used, in which case the symbols from that BFD should be added to the linker hash table instead.)

In most cases the work of looking through the symbols in the archive should be done by the `_bfd_generic_link_add_archive_symbols` function. This function builds a hash table from the archive symbol table and looks through the list of undefined symbols to see

which elements should be included. `_bfd_generic_link_add_archive_symbols` is passed a function to call to make the final decision about adding an archive element to the link and to do the actual work of adding the symbols to the linker hash table.

The function passed to `_bfd_generic_link_add_archive_symbols` must read the symbols of the archive element and decide whether the archive element should be included in the link. If the element is to be included, the `add_archive_element` linker callback routine must be called with the element as an argument, and the element's symbols must be added to the linker hash table just as though the element had itself been passed to the `_bfd_link_add_symbols` function. The `add_archive_element` callback has the option to indicate that it would like to replace the element archive with a substitute BFD, in which case it is the symbols of that substitute BFD that must be added to the linker hash table instead.

When the a.out `_bfd_link_add_symbols` function receives an archive, it calls `_bfd_generic_link_add_archive_symbols` passing `aout_link_check_archive_element` as the function argument. `aout_link_check_archive_element` calls `aout_link_check_ar_symbols`. If the latter decides to add the element (an element is only added if it provides a real, non-common, definition for a previously undefined or common symbol) it calls the `add_archive_element` callback and then `aout_link_check_archive_element` calls `aout_link_add_symbols` to actually add the symbols to the linker hash table - possibly those of a substitute BFD, if the `add_archive_element` callback avails itself of that option.

The ECOFF back end is unusual in that it does not normally call `_bfd_generic_link_add_archive_symbols`, because ECOFF archives already contain a hash table of symbols. The ECOFF back end searches the archive itself to avoid the overhead of creating a new hash table.

### 2.17.3 Performing the final link

When all the input files have been processed, the linker calls the `_bfd_final_link` entry point of the output BFD. This routine is responsible for producing the final output file, which has several aspects. It must relocate the contents of the input sections and copy the data into the output sections. It must build an output symbol table including any local symbols from the input files and the global symbols from the hash table. When producing relocatable output, it must modify the input relocations and write them into the output file. There may also be object format dependent work to be done.

The linker will also call the `write_object_contents` entry point when the BFD is closed. The two entry points must work together in order to produce the correct output file.

The details of how this works are inevitably dependent upon the specific object file format. The a.out `_bfd_final_link` routine is `NAME(aout,final_link)`.

#### 2.17.3.1 Information provided by the linker

Before the linker calls the `_bfd_final_link` entry point, it sets up some data structures for the function to use.

The `input_bfds` field of the `bfd_link_info` structure will point to a list of all the input files included in the link. These files are linked through the `link_next` field of the `bfd` structure.



Each section in the output file will have a list of `link_order` structures attached to the `map_head.link_order` field (the `link_order` structure is defined in `bfdlink.h`). These structures describe how to create the contents of the output section in terms of the contents of various input sections, fill constants, and, eventually, other types of information. They also describe relocs that must be created by the BFD backend, but do not correspond to any input file; this is used to support `-Ur`, which builds constructors while generating a relocatable object file.

### 2.17.3.2 Relocating the section contents

The `_bfd_final_link` function should look through the `link_order` structures attached to each section of the output file. Each `link_order` structure should either be handled specially, or it should be passed to the function `_bfd_default_link_order` which will do the right thing (`_bfd_default_link_order` is defined in `linker.c`).

For efficiency, a `link_order` of type `bfd_indirect_link_order` whose associated section belongs to a BFD of the same format as the output BFD must be handled specially. This type of `link_order` describes part of an output section in terms of a section belonging to one of the input files. The `_bfd_final_link` function should read the contents of the section and any associated relocs, apply the relocs to the section contents, and write out the modified section contents. If performing a relocatable link, the relocs themselves must also be modified and written out.

The functions `_bfd_relocate_contents` and `_bfd_final_link_relocate` provide some general support for performing the actual relocations, notably overflow checking. Their arguments include information about the symbol the relocation is against and a `reloc_howto_type` argument which describes the relocation to perform. These functions are defined in `reloc.c`.

The `a.out` function which handles reading, relocating, and writing section contents is `aout_link_input_section`. The actual relocation is done in `aout_link_input_section_std` and `aout_link_input_section_ext`.

### 2.17.3.3 Writing the symbol table

The `_bfd_final_link` function must gather all the symbols in the input files and write them out. It must also write out all the symbols in the global hash table. This must be controlled by the `strip` and `discard` fields of the `bfd_link_info` structure.

The local symbols of the input files will not have been entered into the linker hash table. The `_bfd_final_link` routine must consider each input file and include the symbols in the output file. It may be convenient to do this when looking through the `link_order` structures, or it may be done by stepping through the `input_bfds` list.

The `_bfd_final_link` routine must also traverse the global hash table to gather all the externally visible symbols. It is possible that most of the externally visible symbols may be written out when considering the symbols of each input file, but it is still necessary to traverse the hash table since the linker script may have defined some symbols that are not in any of the input files.

The `strip` field of the `bfd_link_info` structure controls which symbols are written out. The possible values are listed in `bfdlink.h`. If the value is `strip_some`, then the `keep_hash` field of the `bfd_link_info` structure is a hash table of symbols to keep; each symbol should



be looked up in this hash table, and only symbols which are present should be included in the output file.

If the `strip` field of the `bfd_link_info` structure permits local symbols to be written out, the `discard` field is used to further controls which local symbols are included in the output file. If the value is `discard_1`, then all local symbols which begin with a certain prefix are discarded; this is controlled by the `bfd_is_local_label_name` entry point.

The a.out backend handles symbols by calling `aout_link_write_symbols` on each input BFD and then traversing the global hash table with the function `aout_link_write_other_symbol`. It builds a string table while writing out the symbols, which is written to the output file at the end of `NAME(aout,final_link)`.

### 2.17.3.4 bfd\_link\_split\_section

#### Synopsis

```
bfd_boolean bfd_link_split_section (bfd *abfd, asection *sec);
```

#### Description

Return nonzero if `sec` should be split during a reloceatable or final link.

```
#define bfd_link_split_section(abfd, sec) \
 BFD_SEND (abfd, _bfd_link_split_section, (abfd, sec))
```

### 2.17.3.5 bfd\_section\_already\_linked

#### Synopsis

```
void bfd_section_already_linked (bfd *abfd, asection *sec,
 struct bfd_link_info *info);
```

#### Description

Check if `sec` has been already linked during a reloceatable or final link.

```
#define bfd_section_already_linked(abfd, sec, info) \
 BFD_SEND (abfd, _section_already_linked, (abfd, sec, info))
```

### 2.17.3.6 bfd\_generic\_define\_common\_symbol

#### Synopsis

```
bfd_boolean bfd_generic_define_common_symbol
 (bfd *output_bfd, struct bfd_link_info *info,
 struct bfd_link_hash_entry *h);
```

#### Description

Convert common symbol `h` into a defined symbol. Return TRUE on success and FALSE on failure.

```
#define bfd_define_common_symbol(output_bfd, info, h) \
 BFD_SEND (output_bfd, _bfd_define_common_symbol, (output_bfd, info, h))■
```

### 2.17.3.7 bfd\_find\_version\_for\_sym

#### Synopsis

```
struct bfd_elf_version_tree * bfd_find_version_for_sym
(struct bfd_elf_version_tree *verdefs,
 const char *sym_name, bfd_boolean *hide);
```

#### Description

Search an elf version script tree for symbol versioning info and export / don't-export status for a given symbol. Return non-NULL on success and NULL on failure; also sets the output 'hide' boolean parameter.

## 2.18 Hash Tables

BFD provides a simple set of hash table functions. Routines are provided to initialize a hash table, to free a hash table, to look up a string in a hash table and optionally create an entry for it, and to traverse a hash table. There is currently no routine to delete a string from a hash table.

The basic hash table does not permit any data to be stored with a string. However, a hash table is designed to present a base class from which other types of hash tables may be derived. These derived types may store additional information with the string. Hash tables were implemented in this way, rather than simply providing a data pointer in a hash table entry, because they were designed for use by the linker back ends. The linker may create thousands of hash table entries, and the overhead of allocating private data and storing and following pointers becomes noticeable.

The basic hash table code is in `hash.c`.

### 2.18.1 Creating and freeing a hash table

To create a hash table, create an instance of a `struct bfd_hash_table` (defined in `bfd.h`) and call `bfd_hash_table_init` (if you know approximately how many entries you will need, the function `bfd_hash_table_init_n`, which takes a *size* argument, may be used). `bfd_hash_table_init` returns `FALSE` if some sort of error occurs.

The function `bfd_hash_table_init` take as an argument a function to use to create new entries. For a basic hash table, use the function `bfd_hash_newfunc`. See [Section 2.18.4 \[Deriving a New Hash Table Type\]](#), page 147, for why you would want to use a different value for this argument.

`bfd_hash_table_init` will create an `objalloc` which will be used to allocate new entries. You may allocate memory on this `objalloc` using `bfd_hash_allocate`.

Use `bfd_hash_table_free` to free up all the memory that has been allocated for a hash table. This will not free up the `struct bfd_hash_table` itself, which you must provide.

Use `bfd_hash_set_default_size` to set the default size of hash table to use.

### 2.18.2 Looking up or entering a string

The function `bfd_hash_lookup` is used both to look up a string in the hash table and to create a new entry.

If the *create* argument is `FALSE`, `bfd_hash_lookup` will look up a string. If the string is found, it will returns a pointer to a `struct bfd_hash_entry`. If the string is not found in

the table `bfd_hash_lookup` will return `NULL`. You should not modify any of the fields in the returns `struct bfd_hash_entry`.

If the *create* argument is `TRUE`, the string will be entered into the hash table if it is not already there. Either way a pointer to a `struct bfd_hash_entry` will be returned, either to the existing structure or to a newly created one. In this case, a `NULL` return means that an error occurred.

If the *create* argument is `TRUE`, and a new entry is created, the *copy* argument is used to decide whether to copy the string onto the hash table objalloc or not. If *copy* is passed as `FALSE`, you must be careful not to deallocate or modify the string as long as the hash table exists.

### 2.18.3 Traversing a hash table

The function `bfd_hash_traverse` may be used to traverse a hash table, calling a function on each element. The traversal is done in a random order.

`bfd_hash_traverse` takes as arguments a function and a generic `void *` pointer. The function is called with a hash table entry (a `struct bfd_hash_entry *`) and the generic pointer passed to `bfd_hash_traverse`. The function must return a `boolean` value, which indicates whether to continue traversing the hash table. If the function returns `FALSE`, `bfd_hash_traverse` will stop the traversal and return immediately.

### 2.18.4 Deriving a new hash table type

Many uses of hash tables want to store additional information which each entry in the hash table. Some also find it convenient to store additional information with the hash table itself. This may be done using a derived hash table.

Since C is not an object oriented language, creating a derived hash table requires sticking together some boilerplate routines with a few differences specific to the type of hash table you want to create.

An example of a derived hash table is the linker hash table. The structures for this are defined in `bfdlink.h`. The functions are in `linker.c`.

You may also derive a hash table from an already derived hash table. For example, the a.out linker backend code uses a hash table derived from the linker hash table.

#### 2.18.4.1 Define the derived structures

You must define a structure for an entry in the hash table, and a structure for the hash table itself.

The first field in the structure for an entry in the hash table must be of the type used for an entry in the hash table you are deriving from. If you are deriving from a basic hash table this is `struct bfd_hash_entry`, which is defined in `bfd.h`. The first field in the structure for the hash table itself must be of the type of the hash table you are deriving from itself. If you are deriving from a basic hash table, this is `struct bfd_hash_table`.

For example, the linker hash table defines `struct bfd_link_hash_entry` (in `bfdlink.h`). The first field, `root`, is of type `struct bfd_hash_entry`. Similarly, the first field in `struct bfd_link_hash_table`, `table`, is of type `struct bfd_hash_table`.

### 2.18.4.2 Write the derived creation routine

You must write a routine which will create and initialize an entry in the hash table. This routine is passed as the function argument to `bfd_hash_table_init`.

In order to permit other hash tables to be derived from the hash table you are creating, this routine must be written in a standard way.

The first argument to the creation routine is a pointer to a hash table entry. This may be `NULL`, in which case the routine should allocate the right amount of space. Otherwise the space has already been allocated by a hash table type derived from this one.

After allocating space, the creation routine must call the creation routine of the hash table type it is derived from, passing in a pointer to the space it just allocated. This will initialize any fields used by the base hash table.

Finally the creation routine must initialize any local fields for the new hash table type.

Here is a boilerplate example of a creation routine. *function\_name* is the name of the routine. *entry\_type* is the type of an entry in the hash table you are creating. *base\_newfunc* is the name of the creation routine of the hash table type your hash table is derived from.

```
struct bfd_hash_entry *
function_name (struct bfd_hash_entry *entry,
 struct bfd_hash_table *table,
 const char *string)
{
 struct entry_type *ret = (entry_type *) entry;

 /* Allocate the structure if it has not already been allocated by a
 derived class. */
 if (ret == NULL)
 {
 ret = bfd_hash_allocate (table, sizeof (* ret));
 if (ret == NULL)
 return NULL;
 }

 /* Call the allocation method of the base class. */
 ret = ((entry_type *)
 base_newfunc ((struct bfd_hash_entry *) ret, table, string));

 /* Initialize the local fields here. */

 return (struct bfd_hash_entry *) ret;
}
```

#### Description

The creation routine for the linker hash table, which is in `linker.c`, looks just like this example. *function\_name* is `_bfd_link_hash_newfunc`. *entry\_type* is `struct bfd_link_hash_entry`. *base\_newfunc* is `bfd_hash_newfunc`, the creation routine for a basic hash table.

`_bfd_link_hash_newfunc` also initializes the local fields in a linker hash table entry: `type`, `written` and `next`.

### 2.18.4.3 Write other derived routines

You will want to write other routines for your new hash table, as well.

You will want an initialization routine which calls the initialization routine of the hash table you are deriving from and initializes any other local fields. For the linker hash table, this is `_bfd_link_hash_table_init` in `linker.c`.

You will want a lookup routine which calls the lookup routine of the hash table you are deriving from and casts the result. The linker hash table uses `bfd_link_hash_lookup` in `linker.c` (this actually takes an additional argument which it uses to decide how to return the looked up value).

You may want a traversal routine. This should just call the traversal routine of the hash table you are deriving from with appropriate casts. The linker hash table uses `bfd_link_hash_traverse` in `linker.c`.

These routines may simply be defined as macros. For example, the a.out backend linker hash table, which is derived from the linker hash table, uses macros for the lookup and traversal routines. These are `aout_link_hash_lookup` and `aout_link_hash_traverse` in `aoutx.h`.

## 3 BFD back ends

### 3.1 What to Put Where

All of BFD lives in one directory.

### 3.2 a.out backends

#### Description

BFD supports a number of different flavours of a.out format, though the major differences are only the sizes of the structures on disk, and the shape of the relocation information.

The support is split into a basic support file ‘aoutx.h’ and other files which derive functions from the base. One derivation file is ‘aoutf1.h’ (for a.out flavour 1), and adds to the basic a.out functions support for sun3, sun4, 386 and 29k a.out files, to create a target jump vector for a specific target.

This information is further split out into more specific files for each machine, including ‘sunos.c’ for sun3 and sun4, ‘newsos3.c’ for the Sony NEWS, and ‘demo64.c’ for a demonstration of a 64 bit a.out format.

The base file ‘aoutx.h’ defines general mechanisms for reading and writing records to and from disk and various other methods which BFD requires. It is included by ‘aout32.c’ and ‘aout64.c’ to form the names aout\_32\_swap\_exec\_header\_in, aout\_64\_swap\_exec\_header\_in, etc.

As an example, this is what goes on to make the back end for a sun4, from ‘aout32.c’:

```
#define ARCH_SIZE 32
#include "aoutx.h"
```

Which exports names:

```
...
aout_32_canonicalize_reloc
aout_32_find_nearest_line
aout_32_get_lineno
aout_32_get_reloc_upper_bound
...
```

from ‘sunos.c’:

```
#define TARGET_NAME "a.out-sunos-big"
#define VECNAME sunos_big_vec
#include "aoutf1.h"
```

requires all the names from ‘aout32.c’, and produces the jump vector

```
sunos_big_vec
```

The file ‘host-aout.c’ is a special case. It is for a large set of hosts that use “more or less standard” a.out files, and for which cross-debugging is not interesting. It uses the standard 32-bit a.out support routines, but determines the file offsets and addresses of the text, data, and BSS sections, the machine architecture and machine type, and the entry point address, in a host-dependent manner. Once these values have been determined, generic code is used to handle the object file.

When porting it to run on a new system, you must supply:

```

HOST_PAGE_SIZE
HOST_SEGMENT_SIZE
HOST_MACHINE_ARCH (optional)
HOST_MACHINE_MACHINE (optional)
HOST_TEXT_START_ADDR
HOST_STACK_END_ADDR

```

in the file ‘`../include/sys/h-XXX.h`’ (for your host). These values, plus the structures and macros defined in ‘`a.out.h`’ on your host system, will produce a BFD target that will access ordinary a.out files on your host. To configure a new machine to use ‘`host-aout.c`’, specify:

```

TDEFAULTS = -DDEFAULT_VECTOR=host_aout_big_vec
TDEPFILES= host-aout.o trad-core.o

```

in the ‘`config/XXX.mt`’ file, and modify ‘`configure.in`’ to use the ‘`XXX.mt`’ file (by setting “`bfd_target=XXX`”) when your configuration is selected.

### 3.2.1 Relocations

#### Description

The file ‘`aoutx.h`’ provides for both the *standard* and *extended* forms of a.out relocation records.

The standard records contain only an address, a symbol index, and a type field. The extended records (used on 29ks and sparcs) also have a full integer for an addend.

### 3.2.2 Internal entry points

#### Description

‘`aoutx.h`’ exports several routines for accessing the contents of an a.out file, which are gathered and exported in turn by various format specific files (eg `sunos.c`).

#### 3.2.2.1 `aout_size_swap_exec_header_in`

##### Synopsis

```

void aout_size_swap_exec_header_in,
 (bfd *abfd,
 struct external_exec *bytes,
 struct internal_exec *execp);

```

#### Description

Swap the information in an executable header *raw\_bytes* taken from a raw byte stream memory image into the internal exec header structure *execp*.

#### 3.2.2.2 `aout_size_swap_exec_header_out`

##### Synopsis

```

void aout_size_swap_exec_header_out
 (bfd *abfd,
 struct internal_exec *execp,
 struct external_exec *raw_bytes);

```

**Description**

Swap the information in an internal exec header structure *execp* into the buffer *raw\_bytes* ready for writing to disk.

**3.2.2.3 aout\_size\_some\_aout\_object\_p****Synopsis**

```
const bfd_target *aout_size_some_aout_object_p
(bfd *abfd,
 struct internal_exec *execp,
 const bfd_target *(*callback_to_real_object_p) (bfd *));
```

**Description**

Some a.out variant thinks that the file open in *abfd* checking is an a.out file. Do some more checking, and set up for access if it really is. Call back to the calling environment's "finish up" function just before returning, to handle any last-minute setup.

**3.2.2.4 aout\_size\_mkobject****Synopsis**

```
bfd_boolean aout_size_mkobject, (bfd *abfd);
```

**Description**

Initialize BFD *abfd* for use with a.out files.

**3.2.2.5 aout\_size\_machine\_type****Synopsis**

```
enum machine_type aout_size_machine_type
(enum bfd_architecture arch,
 unsigned long machine,
 bfd_boolean *unknown);
```

**Description**

Keep track of machine architecture and machine type for a.out's. Return the *machine\_type* for a particular architecture and machine, or M\_UNKNOWN if that exact architecture and machine can't be represented in a.out format.

If the architecture is understood, machine type 0 (default) is always understood.

**3.2.2.6 aout\_size\_set\_arch\_mach****Synopsis**

```
bfd_boolean aout_size_set_arch_mach,
(bfd *,
 enum bfd_architecture arch,
 unsigned long machine);
```

**Description**

Set the architecture and the machine of the BFD *abfd* to the values *arch* and *machine*. Verify that *abfd*'s format can support the architecture required.



### 3.2.2.7 aout\_size\_new\_section\_hook

#### Synopsis

```
bfd_boolean aout_size_new_section_hook,
 (bfd *abfd,
 asection *newsect);
```

#### Description

Called by the BFD in response to a `bfd_make_section` request.

## 3.3 coff backends

BFD supports a number of different flavours of coff format. The major differences between formats are the sizes and alignments of fields in structures on disk, and the occasional extra field.

Coff in all its varieties is implemented with a few common files and a number of implementation specific files. For example, The 88k bcs coff format is implemented in the file `'coff-m88k.c'`. This file `#includes` `'coff/m88k.h'` which defines the external structure of the coff format for the 88k, and `'coff/internal.h'` which defines the internal structure. `'coff-m88k.c'` also defines the relocations used by the 88k format See [Section 2.10 \[Relocations\]](#), page 51.

The Intel i960 processor version of coff is implemented in `'coff-i960.c'`. This file has the same structure as `'coff-m88k.c'`, except that it includes `'coff/i960.h'` rather than `'coff-m88k.h'`.

### 3.3.1 Porting to a new version of coff

The recommended method is to select from the existing implementations the version of coff which is most like the one you want to use. For example, we'll say that i386 coff is the one you select, and that your coff flavour is called foo. Copy `'i386coff.c'` to `'foocoff.c'`, copy `'../include/coff/i386.h'` to `'../include/coff/foo.h'`, and add the lines to `'targets.c'` and `'Makefile.in'` so that your new back end is used. Alter the shapes of the structures in `'../include/coff/foo.h'` so that they match what you need. You will probably also have to add `#ifdefs` to the code in `'coff/internal.h'` and `'coffcode.h'` if your version of coff is too wild.

You can verify that your new BFD backend works quite simply by building `'objdump'` from the `'binutils'` directory, and making sure that its version of what's going on and your host system's idea (assuming it has the pretty standard coff dump utility, usually called `att-dump` or just `dump`) are the same. Then clean up your code, and send what you've done to Cygnus. Then your stuff will be in the next release, and you won't have to keep integrating it.

### 3.3.2 How the coff backend works

#### 3.3.2.1 File layout

The Coff backend is split into generic routines that are applicable to any Coff target and routines that are specific to a particular target. The target-specific routines are further split into ones which are basically the same for all Coff targets except that they use the external symbol format or use different values for certain constants.

The generic routines are in `'coffgen.c'`. These routines work for any Coff target. They use some hooks into the target specific code; the hooks are in a `bfd_coff_backend_data` structure, one of which exists for each target.

The essentially similar target-specific routines are in `'coffcode.h'`. This header file includes executable C code. The various Coff targets first include the appropriate Coff header file, make any special defines that are needed, and then include `'coffcode.h'`.

Some of the Coff targets then also have additional routines in the target source file itself.

For example, `'coff-i960.c'` includes `'coff/internal.h'` and `'coff/i960.h'`. It then defines a few constants, such as `I960`, and includes `'coffcode.h'`. Since the i960 has complex relocation types, `'coff-i960.c'` also includes some code to manipulate the i960 relocs. This code is not in `'coffcode.h'` because it would not be used by any other target.

### 3.3.2.2 Coff long section names

In the standard Coff object format, section names are limited to the eight bytes available in the `s_name` field of the `SCNHDR` section header structure. The format requires the field to be NUL-padded, but not necessarily NUL-terminated, so the longest section names permitted are a full eight characters.

The Microsoft PE variants of the Coff object file format add an extension to support the use of long section names. This extension is defined in section 4 of the Microsoft PE/COFF specification (rev 8.1). If a section name is too long to fit into the section header's `s_name` field, it is instead placed into the string table, and the `s_name` field is filled with a slash ("/") followed by the ASCII decimal representation of the offset of the full name relative to the string table base.

Note that this implies that the extension can only be used in object files, as executables do not contain a string table. The standard specifies that long section names from objects emitted into executable images are to be truncated.

However, as a GNU extension, BFD can generate executable images that contain a string table and long section names. This would appear to be technically valid, as the standard only says that Coff debugging information is deprecated, not forbidden, and in practice it works, although some tools that parse PE files expecting the MS standard format may become confused; `'PEview'` is one known example.

The functionality is supported in BFD by code implemented under the control of the macro `COFF_LONG_SECTION_NAMES`. If not defined, the format does not support long section names in any way. If defined, it is used to initialise a flag, `_bfd_coff_long_section_names`, and a hook function pointer, `_bfd_coff_set_long_section_names`, in the Coff backend data structure. The flag controls the generation of long section names in output BFDs at runtime; if it is false, as it will be by default when generating an executable image, long section names are truncated; if true, the long section names extension is employed. The hook points to a function that allows the value of the flag to be altered at runtime, on formats that support long section names at all; on other formats it points to a stub that returns an error indication. With input BFDs, the flag is set according to whether any long section names are detected while reading the section headers. For a completely new BFD, the flag is set to the default for the target format. This information can be used by a client of the BFD library when deciding what output format to generate, and means that a

BFD that is opened for read and subsequently converted to a writeable BFD and modified in-place will retain whatever format it had on input.

If `COFF_LONG_SECTION_NAMES` is simply defined (blank), or is defined to the value "1", then long section names are enabled by default; if it is defined to the value zero, they are disabled by default (but still accepted in input BFDs). The header '`coffcode.h`' defines a macro, `COFF_DEFAULT_LONG_SECTION_NAMES`, which is used in the backends to initialise the backend data structure fields appropriately; see the comments for further detail.

### 3.3.2.3 Bit twiddling

Each flavour of coff supported in BFD has its own header file describing the external layout of the structures. There is also an internal description of the coff layout, in '`coff/internal.h`'. A major function of the coff backend is swapping the bytes and twiddling the bits to translate the external form of the structures into the normal internal form. This is all performed in the `bfd_swap_thing_direction` routines. Some elements are different sizes between different versions of coff; it is the duty of the coff version specific include file to override the definitions of various packing routines in '`coffcode.h`'. E.g., the size of line number entry in coff is sometimes 16 bits, and sometimes 32 bits. `#define`ing `PUT_LNSZ_LNNO` and `GET_LNSZ_LNNO` will select the correct one. No doubt, some day someone will find a version of coff which has a varying field size not catered to at the moment. To port BFD, that person will have to add more `#defines`. Three of the bit twiddling routines are exported to `gdb`; `coff_swap_aux_in`, `coff_swap_sym_in` and `coff_swap_lineno_in`. `GDB` reads the symbol table on its own, but uses BFD to fix things up. More of the bit twiddlers are exported for `gas`; `coff_swap_aux_out`, `coff_swap_sym_out`, `coff_swap_lineno_out`, `coff_swap_reloc_out`, `coff_swap_filehdr_out`, `coff_swap_aouthdr_out`, `coff_swap_scnhdr_out`. `Gas` currently keeps track of all the symbol table and reloc drudgery itself, thereby saving the internal BFD overhead, but uses BFD to swap things on the way out, making cross ports much safer. Doing so also allows BFD (and thus the linker) to use the same header files as `gas`, which makes one avenue to disaster disappear.

### 3.3.2.4 Symbol reading

The simple canonical form for symbols used by BFD is not rich enough to keep all the information available in a coff symbol table. The back end gets around this problem by keeping the original symbol table around, "behind the scenes".

When a symbol table is requested (through a call to `bfd_canonicalize_syntab`), a request gets through to `coff_get_normalized_syntab`. This reads the symbol table from the coff file and swaps all the structures inside into the internal form. It also fixes up all the pointers in the table (represented in the file by offsets from the first symbol in the table) into physical pointers to elements in the new internal table. This involves some work since the meanings of fields change depending upon context: a field that is a pointer to another structure in the symbol table at one moment may be the size in bytes of a structure at the next. Another pass is made over the table. All symbols which mark file names (`C_FILE` symbols) are modified so that the internal string points to the value in the auxent (the real filename) rather than the normal text associated with the symbol ("`.file`").

At this time the symbol names are moved around. Coff stores all symbols less than nine characters long physically within the symbol table; longer strings are kept at the end of the

file in the string table. This pass moves all strings into memory and replaces them with pointers to the strings.

The symbol table is massaged once again, this time to create the canonical table used by the BFD application. Each symbol is inspected in turn, and a decision made (using the `sclass` field) about the various flags to set in the `asymbol`. See [Section 2.7 \[Symbols\]](#), page 40. The generated canonical table shares strings with the hidden internal symbol table.

Any linenumbers are read from the coff file too, and attached to the symbols which own the functions the linenumbers belong to.

### 3.3.2.5 Symbol writing

Writing a symbol to a coff file which didn't come from a coff file will lose any debugging information. The `asymbol` structure remembers the BFD from which the symbol was taken, and on output the back end makes sure that the same destination target as source target is present.

When the symbols have come from a coff file then all the debugging information is preserved. Symbol tables are provided for writing to the back end in a vector of pointers to pointers. This allows applications like the linker to accumulate and output large symbol tables without having to do too much byte copying.

This function runs through the provided symbol table and patches each symbol marked as a file place holder (`C_FILE`) to point to the next file place holder in the list. It also marks each `offset` field in the list with the offset from the first symbol of the current symbol.

Another function of this procedure is to turn the canonical value form of BFD into the form used by coff. Internally, BFD expects symbol values to be offsets from a section base; so a symbol physically at 0x120, but in a section starting at 0x100, would have the value 0x20. Coff expects symbols to contain their final value, so symbols have their values changed at this point to reflect their sum with their owning section. This transformation uses the `output_section` field of the `asymbol`'s `asection`. See [Section 2.6 \[Sections\]](#), page 22.

- `coff_mangle_symbols`

This routine runs through the provided symbol table and uses the offsets generated by the previous pass and the pointers generated when the symbol table was read in to create the structured hierarchy required by coff. It changes each pointer to a symbol into the index into the symbol table of the `asymbol`.

- `coff_write_symbols`

This routine runs through the symbol table and patches up the symbols from their internal form into the coff way, calls the bit twiddlers, and writes out the table to the file.

### 3.3.2.6 `coff_symbol_type`

#### Description

The hidden information for an `asymbol` is described in a `combined_entry_type`:

```
typedef struct coff_ptr_struct
{
 /* Remembers the offset from the first symbol in the file for
 this symbol. Generated by coff_renumber_symbols. */

```

```

 unsigned int offset;

 /* Should the value of this symbol be renumbered. Used for
 XCOFF C_BSTAT symbols. Set by coff_slurp_symbol_table. */
 unsigned int fix_value : 1;

 /* Should the tag field of this symbol be renumbered.
 Created by coff_pointerize_aux. */
 unsigned int fix_tag : 1;

 /* Should the endidx field of this symbol be renumbered.
 Created by coff_pointerize_aux. */
 unsigned int fix_end : 1;

 /* Should the x_csect.x_scnlen field be renumbered.
 Created by coff_pointerize_aux. */
 unsigned int fix_scnlen : 1;

 /* Fix up an XCOFF C_BINCL/C_EINCL symbol. The value is the
 index into the line number entries. Set by coff_slurp_symbol_table. */
 unsigned int fix_line : 1;

 /* The container for the symbol structure as read and translated
 from the file. */
 union
 {
 union internal_auxent auxent;
 struct internal_syment syment;
 } u;
} combined_entry_type;

/* Each canonical asymbol really looks like this: */

typedef struct coff_symbol_struct
{
 /* The actual symbol which the rest of BFD works with */
 asymbol symbol;

 /* A pointer to the hidden information for this symbol */
 combined_entry_type *native;

 /* A pointer to the lineno information for this symbol */
 struct lineno_cache_entry *lineno;

 /* Have the line numbers been relocated yet ? */
 bfd_boolean done_lineno;

```

```
 } coff_symbol_type;
```

### 3.3.2.7 bfd\_coff\_backend\_data

```
/* COFF symbol classifications. */

enum coff_symbol_classification
{
 /* Global symbol. */
 COFF_SYMBOL_GLOBAL,
 /* Common symbol. */
 COFF_SYMBOL_COMMON,
 /* Undefined symbol. */
 COFF_SYMBOL_UNDEFINED,
 /* Local symbol. */
 COFF_SYMBOL_LOCAL,
 /* PE section symbol. */
 COFF_SYMBOL_PE_SECTION
};
```

Special entry points for gdb to swap in coff symbol table parts:

```
typedef struct
{
 void (*_bfd_coff_swap_aux_in)
 (bfd *, void *, int, int, int, int, void *);

 void (*_bfd_coff_swap_sym_in)
 (bfd *, void *, void *);

 void (*_bfd_coff_swap_lineno_in)
 (bfd *, void *, void *);

 unsigned int (*_bfd_coff_swap_aux_out)
 (bfd *, void *, int, int, int, int, void *);

 unsigned int (*_bfd_coff_swap_sym_out)
 (bfd *, void *, void *);

 unsigned int (*_bfd_coff_swap_lineno_out)
 (bfd *, void *, void *);

 unsigned int (*_bfd_coff_swap_reloc_out)
 (bfd *, void *, void *);

 unsigned int (*_bfd_coff_swap_filehdr_out)
 (bfd *, void *, void *);
```

```
unsigned int (*_bfd_coff_swap_aouthdr_out)
 (bfd *, void *, void *);

unsigned int (*_bfd_coff_swap_scnhdr_out)
 (bfd *, void *, void *);

unsigned int _bfd_filhsz;
unsigned int _bfd_aoutsz;
unsigned int _bfd_scnhsz;
unsigned int _bfd_symesz;
unsigned int _bfd_auxesz;
unsigned int _bfd_relsz;
unsigned int _bfd_linesz;
unsigned int _bfd_filnmlen;
bfd_boolean _bfd_coff_long_filenames;

bfd_boolean _bfd_coff_long_section_names;
bfd_boolean (*_bfd_coff_set_long_section_names)
 (bfd *, int);

unsigned int _bfd_coff_default_section_alignment_power;
bfd_boolean _bfd_coff_force_symnames_in_strings;
unsigned int _bfd_coff_debug_string_prefix_length;

void (*_bfd_coff_swap_filehdr_in)
 (bfd *, void *, void *);

void (*_bfd_coff_swap_aouthdr_in)
 (bfd *, void *, void *);

void (*_bfd_coff_swap_scnhdr_in)
 (bfd *, void *, void *);

void (*_bfd_coff_swap_reloc_in)
 (bfd *abfd, void *, void *);

bfd_boolean (*_bfd_coff_bad_format_hook)
 (bfd *, void *);

bfd_boolean (*_bfd_coff_set_arch_mach_hook)
 (bfd *, void *);

void * (*_bfd_coff_mkobject_hook)
 (bfd *, void *, void *);

bfd_boolean (*_bfd_styp_to_sec_flags_hook)
 (bfd *, void *, const char *, asection *, flagword *);
```

```

void (*_bfd_set_alignment_hook)
 (bfd *, asection *, void *);

bfd_boolean (*_bfd_coff_slurp_symbol_table)
 (bfd *);

bfd_boolean (*_bfd_coff_symname_in_debug)
 (bfd *, struct internal_syment *);

bfd_boolean (*_bfd_coff_pointerize_aux_hook)
 (bfd *, combined_entry_type *, combined_entry_type *,
 unsigned int, combined_entry_type *);

bfd_boolean (*_bfd_coff_print_aux)
 (bfd *, FILE *, combined_entry_type *, combined_entry_type *,
 combined_entry_type *, unsigned int);

void (*_bfd_coff_reloc16_extra_cases)
 (bfd *, struct bfd_link_info *, struct bfd_link_order *, arelent *,
 bfd_byte *, unsigned int *, unsigned int *);

int (*_bfd_coff_reloc16_estimate)
 (bfd *, asection *, arelent *, unsigned int,
 struct bfd_link_info *);

enum coff_symbol_classification (*_bfd_coff_classify_symbol)
 (bfd *, struct internal_syment *);

bfd_boolean (*_bfd_coff_compute_section_file_positions)
 (bfd *);

bfd_boolean (*_bfd_coff_start_final_link)
 (bfd *, struct bfd_link_info *);

bfd_boolean (*_bfd_coff_relocate_section)
 (bfd *, struct bfd_link_info *, bfd *, asection *, bfd_byte *,
 struct internal_reloc *, struct internal_syment *, asection **);

reloc_howto_type (*_bfd_coff_rtype_to_howto)
 (bfd *, asection *, struct internal_reloc *,
 struct coff_link_hash_entry *, struct internal_syment *,
 bfd_vma *);

bfd_boolean (*_bfd_coff_adjust_symndx)
 (bfd *, struct bfd_link_info *, bfd *, asection *,
 struct internal_reloc *, bfd_boolean *);

```



```

bfd_boolean (*_bfd_coff_link_add_one_symbol)
 (struct bfd_link_info *, bfd *, const char *, flagword,
 asection *, bfd_vma, const char *, bfd_boolean, bfd_boolean,
 struct bfd_link_hash_entry **);

bfd_boolean (*_bfd_coff_link_output_has_begun)
 (bfd *, struct coff_final_link_info *);

bfd_boolean (*_bfd_coff_final_link_postscript)
 (bfd *, struct coff_final_link_info *);

bfd_boolean (*_bfd_coff_print_pdata)
 (bfd *, void *);

} bfd_coff_backend_data;

#define coff_backend_info(abfd) \
 ((bfd_coff_backend_data *) (abfd)->xvec->backend_data)

#define bfd_coff_swap_aux_in(a,e,t,c,ind,num,i) \
 ((coff_backend_info (a)->_bfd_coff_swap_aux_in) (a,e,t,c,ind,num,i))

#define bfd_coff_swap_sym_in(a,e,i) \
 ((coff_backend_info (a)->_bfd_coff_swap_sym_in) (a,e,i))

#define bfd_coff_swap_lineno_in(a,e,i) \
 ((coff_backend_info (a)->_bfd_coff_swap_lineno_in) (a,e,i))

#define bfd_coff_swap_reloc_out(abfd, i, o) \
 ((coff_backend_info (abfd)->_bfd_coff_swap_reloc_out) (abfd, i, o))

#define bfd_coff_swap_lineno_out(abfd, i, o) \
 ((coff_backend_info (abfd)->_bfd_coff_swap_lineno_out) (abfd, i, o))

#define bfd_coff_swap_aux_out(a,i,t,c,ind,num,o) \
 ((coff_backend_info (a)->_bfd_coff_swap_aux_out) (a,i,t,c,ind,num,o))

#define bfd_coff_swap_sym_out(abfd, i,o) \
 ((coff_backend_info (abfd)->_bfd_coff_swap_sym_out) (abfd, i, o))

#define bfd_coff_swap_scnhdr_out(abfd, i,o) \
 ((coff_backend_info (abfd)->_bfd_coff_swap_scnhdr_out) (abfd, i, o))

#define bfd_coff_swap_filehdr_out(abfd, i,o) \
 ((coff_backend_info (abfd)->_bfd_coff_swap_filehdr_out) (abfd, i, o))

```

```

#define bfd_coff_swap_aouthdr_out(abfd, i,o) \
 ((coff_backend_info (abfd)->_bfd_coff_swap_aouthdr_out) (abfd, i, o))

#define bfd_coff_filhsz(abfd) (coff_backend_info (abfd)->_bfd_filhsz)
#define bfd_coff_aoutsz(abfd) (coff_backend_info (abfd)->_bfd_aoutsz)
#define bfd_coff_scnhsz(abfd) (coff_backend_info (abfd)->_bfd_scnhsz)
#define bfd_coff_symesz(abfd) (coff_backend_info (abfd)->_bfd_symesz)
#define bfd_coff_auxesz(abfd) (coff_backend_info (abfd)->_bfd_auxesz)
#define bfd_coff_relsz(abfd) (coff_backend_info (abfd)->_bfd_relsz)
#define bfd_coff_linesz(abfd) (coff_backend_info (abfd)->_bfd_linesz)
#define bfd_coff_filnmlen(abfd) (coff_backend_info (abfd)->_bfd_filnmlen)
#define bfd_coff_long_filenames(abfd) \
 (coff_backend_info (abfd)->_bfd_coff_long_filenames)
#define bfd_coff_long_section_names(abfd) \
 (coff_backend_info (abfd)->_bfd_coff_long_section_names)
#define bfd_coff_set_long_section_names(abfd, enable) \
 ((coff_backend_info (abfd)->_bfd_coff_set_long_section_names) (abfd, enable))
#define bfd_coff_default_section_alignment_power(abfd) \
 (coff_backend_info (abfd)->_bfd_coff_default_section_alignment_power)
#define bfd_coff_swap_filehdr_in(abfd, i,o) \
 ((coff_backend_info (abfd)->_bfd_coff_swap_filehdr_in) (abfd, i, o))

#define bfd_coff_swap_aouthdr_in(abfd, i,o) \
 ((coff_backend_info (abfd)->_bfd_coff_swap_aouthdr_in) (abfd, i, o))

#define bfd_coff_swap_scnhdr_in(abfd, i,o) \
 ((coff_backend_info (abfd)->_bfd_coff_swap_scnhdr_in) (abfd, i, o))

#define bfd_coff_swap_reloc_in(abfd, i, o) \
 ((coff_backend_info (abfd)->_bfd_coff_swap_reloc_in) (abfd, i, o))

#define bfd_coff_bad_format_hook(abfd, filehdr) \
 ((coff_backend_info (abfd)->_bfd_coff_bad_format_hook) (abfd, filehdr))

#define bfd_coff_set_arch_mach_hook(abfd, filehdr)\
 ((coff_backend_info (abfd)->_bfd_coff_set_arch_mach_hook) (abfd, filehdr))
#define bfd_coff_mkobject_hook(abfd, filehdr, aouthdr)\
 ((coff_backend_info (abfd)->_bfd_coff_mkobject_hook)\
 (abfd, filehdr, aouthdr))

#define bfd_coff_styp_to_sec_flags_hook(abfd, scnhdr, name, section, flags_ptr)\
 ((coff_backend_info (abfd)->_bfd_styp_to_sec_flags_hook)\
 (abfd, scnhdr, name, section, flags_ptr))

#define bfd_coff_set_alignment_hook(abfd, sec, scnhdr)\
 ((coff_backend_info (abfd)->_bfd_set_alignment_hook) (abfd, sec, scnhdr))

```



```

((coff_backend_info (abfd)->_bfd_coff_link_add_one_symbol)\
 (info, abfd, name, flags, section, value, string, cp, coll, hashp))

#define bfd_coff_link_output_has_begun(a,p) \
 ((coff_backend_info (a)->_bfd_coff_link_output_has_begun) (a, p))
#define bfd_coff_final_link_postscript(a,p) \
 ((coff_backend_info (a)->_bfd_coff_final_link_postscript) (a, p))

#define bfd_coff_have_print_pdata(a) \
 (coff_backend_info (a)->_bfd_coff_print_pdata)
#define bfd_coff_print_pdata(a,p) \
 ((coff_backend_info (a)->_bfd_coff_print_pdata) (a, p))

/* Macro: Returns true if the bfd is a PE executable as opposed to a
 PE object file. */
#define bfd_pei_p(abfd) \
 (CONST_STRNEQ ((abfd)->xvec->name, "pei-"))

```

### 3.3.2.8 Writing relocations

To write relocations, the back end steps through the canonical relocation table and create an `internal_reloc`. The symbol index to use is removed from the `offset` field in the symbol table supplied. The address comes directly from the sum of the section base address and the relocation offset; the type is dug directly from the `howto` field. Then the `internal_reloc` is swapped into the shape of an `external_reloc` and written out to disk.

### 3.3.2.9 Reading linenumbers

Creating the linenumber table is done by reading in the entire coff linenumber table, and creating another table for internal use.

A coff linenumber table is structured so that each function is marked as having a line number of 0. Each line within the function is an offset from the first line in the function. The base of the line number information for the table is stored in the symbol associated with the function.

Note: The PE format uses line number 0 for a flag indicating a new source file.

The information is copied from the external to the internal table, and each symbol which marks a function is marked by pointing its...

How does this work ?

### 3.3.2.10 Reading relocations

Coff relocations are easily transformed into the internal BFD form (`arelent`).

Reading a coff relocation table is done in the following stages:

- Read the entire coff relocation table into memory.
- Process each relocation in turn; first swap it from the external to the internal form.
- Turn the symbol referenced in the relocation's symbol index into a pointer into the canonical symbol table. This table is the same as the one returned by a call to `bfd_`

`canonicalize_symtab`. The back end will call that routine and save the result if a canonicalization hasn't been done.

- The reloc index is turned into a pointer to a howto structure, in a back end specific way. For instance, the 386 and 960 use the `r_type` to directly produce an index into a howto table vector; the 88k subtracts a number from the `r_type` field and creates an addend field.

### 3.4 ELF backends

BFD support for ELF formats is being worked on. Currently, the best supported back ends are for sparc and i386 (running svr4 or Solaris 2).

Documentation of the internals of the support code still needs to be written. The code is changing quickly enough that we haven't bothered yet.

### 3.5 mmo backend

The mmo object format is used exclusively together with Professor Donald E. Knuth's educational 64-bit processor MMIX. The simulator `mmix` which is available at <http://www-cs-faculty.stanford.edu/~knuth/programs/mmix.tar.gz> understands this format. That package also includes a combined assembler and linker called `mmixal`. The mmo format has no advantages feature-wise compared to e.g. ELF. It is a simple non-relocatable object format with no support for archives or debugging information, except for symbol value information and line numbers (which is not yet implemented in BFD). See <http://www-cs-faculty.stanford.edu/~knuth/mmix.html> for more information about MMIX. The ELF format is used for intermediate object files in the BFD implementation.

#### 3.5.1 File layout

The mmo file contents is not partitioned into named sections as with e.g. ELF. Memory areas is formed by specifying the location of the data that follows. Only the memory area '0x0000...00' to '0x01ff...ff' is executable, so it is used for code (and constants) and the area '0x2000...00' to '0x20ff...ff' is used for writable data. See [Section 3.5.3 \[mmo section mapping\]](#), page 169.

There is provision for specifying "special data" of 65536 different types. We use type 80 (decimal), arbitrarily chosen the same as the ELF `e_machine` number for MMIX, filling it with section information normally found in ELF objects. See [Section 3.5.3 \[mmo section mapping\]](#), page 169.

Contents is entered as 32-bit words, xored over previous contents, always zero-initialized. A word that starts with the byte '0x98' forms a command called a 'lopcode', where the next byte distinguished between the thirteen lopcodes. The two remaining bytes, called the 'Y' and 'Z' fields, or the 'YZ' field (a 16-bit big-endian number), are used for various purposes different for each lopcode. As documented in <http://www-cs-faculty.stanford.edu/~knuth/mmixal-intro.ps.gz>, the lopcodes are:

`lop_quote`

0x98000001. The next word is contents, regardless of whether it starts with 0x98 or not.

|                        |                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                            |
|------------------------|----------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|
| <code>lop_loc</code>   | 0x9801YYZZ, where ‘Z’ is 1 or 2. This is a location directive, setting the location for the next data to the next 32-bit word (for $Z = 1$ ) or 64-bit word (for $Z = 2$ ), plus $Y * 2^56$ . Normally ‘Y’ is 0 for the text segment and 2 for the data segment.                                                                                                                                                                                                                                                                                                                                                                                                           |
| <code>lop_skip</code>  | 0x9802YYZZ. Increase the current location by ‘YZ’ bytes.                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                   |
| <code>lop_fixo</code>  | 0x9803YYZZ, where ‘Z’ is 1 or 2. Store the current location as 64 bits into the location pointed to by the next 32-bit ( $Z = 1$ ) or 64-bit ( $Z = 2$ ) word, plus $Y * 2^56$ .                                                                                                                                                                                                                                                                                                                                                                                                                                                                                           |
| <code>lop_fixr</code>  | 0x9804YYZZ. ‘YZ’ is stored into the current location plus $2 - 4 * YZ$ .                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                   |
| <code>lop_fixrx</code> | 0x980500ZZ. ‘Z’ is 16 or 24. A value ‘L’ derived from the following 32-bit word are used in a manner similar to ‘YZ’ in <code>lop_fixr</code> : it is xor:ed into the current location minus $4 * L$ . The first byte of the word is 0 or 1. If it is 1, then $L = (\text{lowest24bitsofword}) - 2^Z$ , if 0, then $L = (\text{lowest24bitsofword})$ .                                                                                                                                                                                                                                                                                                                     |
| <code>lop_file</code>  | 0x9806YYZZ. ‘Y’ is the file number, ‘Z’ is count of 32-bit words. Set the file number to ‘Y’ and the line counter to 0. The next $Z * 4$ bytes contain the file name, padded with zeros if the count is not a multiple of four. The same ‘Y’ may occur multiple times, but ‘Z’ must be 0 for all but the first occurrence.                                                                                                                                                                                                                                                                                                                                                 |
| <code>lop_line</code>  | 0x9807YYZZ. ‘YZ’ is the line number. Together with <code>lop_file</code> , it forms the source location for the next 32-bit word. Note that for each non-lopcode 32-bit word, line numbers are assumed incremented by one.                                                                                                                                                                                                                                                                                                                                                                                                                                                 |
| <code>lop_spec</code>  | 0x9808YYZZ. ‘YZ’ is the type number. Data until the next lopcode other than <code>lop_quote</code> forms special data of type ‘YZ’. See <a href="#">Section 3.5.3 [mmo section mapping]</a> , page 169.<br><br>Other types than 80, (or type 80 with a content that does not parse) is stored in sections named <code>.MMIX.spec_data.n</code> where $n$ is the ‘YZ’-type. The flags for such a sections say not to allocate or load the data. The vma is 0. Contents of multiple occurrences of special data $n$ is concatenated to the data of the previous <code>lop_spec ns</code> . The location in data or code at which the <code>lop_spec</code> occurred is lost. |
| <code>lop_pre</code>   | 0x980901ZZ. The first lopcode in a file. The ‘Z’ field forms the length of header information in 32-bit words, where the first word tells the time in seconds since ‘00:00:00 GMT Jan 1 1970’.                                                                                                                                                                                                                                                                                                                                                                                                                                                                             |
| <code>lop_post</code>  | 0x980a00ZZ. $Z > 32$ . This lopcode follows after all content-generating lpcodes in a program. The ‘Z’ field denotes the value of ‘rG’ at the beginning of the program. The following $256 - Z$ big-endian 64-bit words are loaded into global registers ‘\$G’ ... ‘\$255’.                                                                                                                                                                                                                                                                                                                                                                                                |
| <code>lop_stab</code>  | 0x980b0000. The next-to-last lopcode in a program. Must follow immediately after the <code>lop_post</code> lopcode and its data. After this lopcode follows all symbols in a compressed format (see <a href="#">Section 3.5.2 [Symbol-table]</a> , page 167).                                                                                                                                                                                                                                                                                                                                                                                                              |
| <code>lop_end</code>   | 0x980cYYZZ. The last lopcode in a program. It must follow the <code>lop_stab</code> lopcode and its data. The ‘YZ’ field contains the number of 32-bit words of symbol table information after the preceding <code>lop_stab</code> lopcode.                                                                                                                                                                                                                                                                                                                                                                                                                                |

Note that the opcode "fixups"; `lop_fixr`, `lop_fixrx` and `lop_fixo` are not generated by BFD, but are handled. They are generated by `mmixal`.

This trivial one-label, one-instruction file:

```
:Main TRAP 1,2,3
```

can be represented this way in mmo:

```
0x98090101 - lop_pre, one 32-bit word with timestamp.
<timestamp>
0x98010002 - lop_loc, text segment, using a 64-bit address.
 Note that mmixal does not emit this for the file above.
0x00000000 - Address, high 32 bits.
0x00000000 - Address, low 32 bits.
0x98060002 - lop_file, 2 32-bit words for file-name.
0x74657374 - "test"
0x2e730000 - ".s\0\0"
0x98070001 - lop_line, line 1.
0x00010203 - TRAP 1,2,3
0x980a00ff - lop_post, setting $255 to 0.
0x00000000
0x00000000
0x980b0000 - lop_stab for ":Main" = 0, serial 1.
0x203a4040 See Section 3.5.2 [Symbol-table], page 167.
0x10404020
0x4d206120
0x69016e00
0x81000000
0x980c0005 - lop_end; symbol table contained five 32-bit words.
```

### 3.5.2 Symbol table format

From `mmixal.w` (or really, the generated `mmixal.tex`) in <http://www-cs-faculty.stanford.edu/~knuth/programmmixal.w>:  
 “Symbols are stored and retrieved by means of a ‘ternary search trie’, following ideas of Bentley and Sedgewick. (See ACM–SIAM Symp. on Discrete Algorithms ‘8’ (1997), 360–369; R.Sedgewick, ‘Algorithms in C’ (Reading, Mass. Addison–Wesley, 1998), ‘15.4’.) Each trie node stores a character, and there are branches to subtries for the cases where a given character is less than, equal to, or greater than the character in the trie. There also is a pointer to a symbol table entry if a symbol ends at the current node.”

So it’s a tree encoded as a stream of bytes. The stream of bytes acts on a single virtual global symbol, adding and removing characters and signalling complete symbol points. Here, we read the stream and create symbols at the completion points.

First, there’s a control byte `m`. If any of the listed bits in `m` is nonzero, we execute what stands at the right, in the listed order:

```
(MM03_LEFT)
0x40 - Traverse left trie.
 (Read a new command byte and recurse.)

(MM03_SYMBITS)
```

0x2f - Read the next byte as a character and store it in the current character position; increment character position. Test the bits of m:

(MM03\_WCHAR)

0x80 - The character is 16-bit (so read another byte, merge into current character.

(MM03\_TYPEBITS)

0xf - We have a complete symbol; parse the type, value and serial number and do what should be done with a symbol. The type and length information is in  $j = (m \& 0xf)$ .

(MM03\_REGQUAL\_BITS)

$j == 0xf$ : A register variable. The following byte tells which register.  
 $j \leq 8$ : An absolute symbol. Read  $j$  bytes as the big-endian number the symbol equals. A  $j = 2$  with two zero bytes denotes an unknown symbol.  
 $j > 8$ : As with  $j \leq 8$ , but add  $(0x20 \ll 56)$  to the value in the following  $j - 8$  bytes.

Then comes the serial number, as a variant of uleb128, but better named ubeb128:  
 Read bytes and shift the previous value left 7 (multiply by 128). Add in the new byte, repeat until a byte has bit 7 set. The serial number is the computed value minus 128.

(MM03\_MIDDLE)

0x20 - Traverse middle trie. (Read a new command byte and recurse.) Decrement character position.

(MM03\_RIGHT)

0x10 - Traverse right trie. (Read a new command byte and recurse.)

Let's look again at the `lop_stab` for the trivial file (see [Section 3.5.1 \[File layout\]](#), page 165).

```
0x980b0000 - lop_stab for ":Main" = 0, serial 1.
0x203a4040
0x10404020
0x4d206120
0x69016e00
0x81000000
```



This forms the trivial trie (note that the path between “:” and “M” is redundant):

```

203a ":"
40 /
40 /
10 \
40 /
40 /
204d "M"
2061 "a"
2069 "i"
016e "n" is the last character in a full symbol, and
 with a value represented in one byte.
00 The value is 0.
81 The serial number is 1.
```

### 3.5.3 mmo section mapping

The implementation in BFD uses special data type 80 (decimal) to encapsulate and describe named sections, containing e.g. debug information. If needed, any datum in the encapsulation will be quoted using `lop_quote`. First comes a 32-bit word holding the number of 32-bit words containing the zero-terminated zero-padded segment name. After the name there's a 32-bit word holding flags describing the section type. Then comes a 64-bit big-endian word with the section length (in bytes), then another with the section start address. Depending on the type of section, the contents might follow, zero-padded to 32-bit boundary. For a loadable section (such as data or code), the contents might follow at some later point, not necessarily immediately, as a `lop_loc` with the same start address as in the section description, followed by the contents. This in effect forms a descriptor that must be emitted before the actual contents. Sections described this way must not overlap.

For areas that don't have such descriptors, synthetic sections are formed by BFD. Consecutive contents in the two memory areas '0x0000...00' to '0x01ff...ff' and '0x2000...00' to '0x20ff...ff' are entered in sections named `.text` and `.data` respectively. If an area is not otherwise described, but would together with a neighboring lower area be less than '0x40000000' bytes long, it is joined with the lower area and the gap is zero-filled. For other cases, a new section is formed, named `.MMIX.sec.n`. Here, *n* is a number, a running count through the mmo file, starting at 0.

A loadable section specified as:

```

.section secname,"ax"
TETRA 1,2,3,4,-1,-2009
BYTE 80
```

and linked to address '0x4', is represented by the sequence:

```

0x98080050 - lop_spec 80
0x00000002 - two 32-bit words for the section name
0x7365636e - "secn"
0x616d6500 - "ame\0"
0x00000033 - flags CODE, READONLY, LOAD, ALLOC
0x00000000 - high 32 bits of section length
```

```

0x0000001c - section length is 28 bytes; 6 * 4 + 1 + alignment to 32 bits
0x00000000 - high 32 bits of section address
0x00000004 - section address is 4
0x98010002 - 64 bits with address of following data
0x00000000 - high 32 bits of address
0x00000004 - low 32 bits: data starts at address 4
0x00000001 - 1
0x00000002 - 2
0x00000003 - 3
0x00000004 - 4
0xffffffff - -1
0xfffff827 - -2009
0x50000000 - 80 as a byte, padded with zeros.

```

Note that the `lop_spec` wrapping does not include the section contents. Compare this to a non-loaded section specified as:

```

.section thirdsec
TETRA 200001,100002
BYTE 38,40

```

This, when linked to address ‘0x200000000000001c’, is represented by:

```

0x98080050 - lop_spec 80
0x00000002 - two 32-bit words for the section name
0x7365636e - "thir"
0x616d6500 - "dsec"
0x00000010 - flag READONLY
0x00000000 - high 32 bits of section length
0x0000000c - section length is 12 bytes; 2 * 4 + 2 + alignment to 32 bits
0x20000000 - high 32 bits of address
0x0000001c - low 32 bits of address 0x200000000000001c
0x00030d41 - 200001
0x000186a2 - 100002
0x26280000 - 38, 40 as bytes, padded with zeros

```

For the latter example, the section contents must not be loaded in memory, and is therefore specified as part of the special data. The address is usually unimportant but might provide information for e.g. the DWARF 2 debugging format.

Version 1.3, 3 November 2008

Copyright © 2000, 2001, 2002, 2007, 2008 Free Software Foundation, Inc.

<http://fsf.org/>

Everyone is permitted to copy and distribute verbatim copies of this license document, but changing it is not allowed.

## 0. PREAMBLE

The purpose of this License is to make a manual, textbook, or other functional and useful document *free* in the sense of freedom: to assure everyone the effective freedom to copy and redistribute it, with or without modifying it, either commercially or non-commercially. Secondly, this License preserves for the author and publisher a way

to get credit for their work, while not being considered responsible for modifications made by others.

This License is a kind of “copyleft”, which means that derivative works of the document must themselves be free in the same sense. It complements the GNU General Public License, which is a copyleft license designed for free software.

We have designed this License in order to use it for manuals for free software, because free software needs free documentation: a free program should come with manuals providing the same freedoms that the software does. But this License is not limited to software manuals; it can be used for any textual work, regardless of subject matter or whether it is published as a printed book. We recommend this License principally for works whose purpose is instruction or reference.

## 1. APPLICABILITY AND DEFINITIONS

This License applies to any manual or other work, in any medium, that contains a notice placed by the copyright holder saying it can be distributed under the terms of this License. Such a notice grants a world-wide, royalty-free license, unlimited in duration, to use that work under the conditions stated herein. The “Document”, below, refers to any such manual or work. Any member of the public is a licensee, and is addressed as “you”. You accept the license if you copy, modify or distribute the work in a way requiring permission under copyright law.

A “Modified Version” of the Document means any work containing the Document or a portion of it, either copied verbatim, or with modifications and/or translated into another language.

A “Secondary Section” is a named appendix or a front-matter section of the Document that deals exclusively with the relationship of the publishers or authors of the Document to the Document’s overall subject (or to related matters) and contains nothing that could fall directly within that overall subject. (Thus, if the Document is in part a textbook of mathematics, a Secondary Section may not explain any mathematics.) The relationship could be a matter of historical connection with the subject or with related matters, or of legal, commercial, philosophical, ethical or political position regarding them.

The “Invariant Sections” are certain Secondary Sections whose titles are designated, as being those of Invariant Sections, in the notice that says that the Document is released under this License. If a section does not fit the above definition of Secondary then it is not allowed to be designated as Invariant. The Document may contain zero Invariant Sections. If the Document does not identify any Invariant Sections then there are none.

The “Cover Texts” are certain short passages of text that are listed, as Front-Cover Texts or Back-Cover Texts, in the notice that says that the Document is released under this License. A Front-Cover Text may be at most 5 words, and a Back-Cover Text may be at most 25 words.

A “Transparent” copy of the Document means a machine-readable copy, represented in a format whose specification is available to the general public, that is suitable for revising the document straightforwardly with generic text editors or (for images composed of pixels) generic paint programs or (for drawings) some widely available drawing editor, and that is suitable for input to text formatters or for automatic translation to a variety of formats suitable for input to text formatters. A copy made in an otherwise

Transparent file format whose markup, or absence of markup, has been arranged to thwart or discourage subsequent modification by readers is not Transparent. An image format is not Transparent if used for any substantial amount of text. A copy that is not “Transparent” is called “Opaque”.

Examples of suitable formats for Transparent copies include plain ASCII without markup, Texinfo input format, LaTeX input format, SGML or XML using a publicly available DTD, and standard-conforming simple HTML, PostScript or PDF designed for human modification. Examples of transparent image formats include PNG, XCF and JPG. Opaque formats include proprietary formats that can be read and edited only by proprietary word processors, SGML or XML for which the DTD and/or processing tools are not generally available, and the machine-generated HTML, PostScript or PDF produced by some word processors for output purposes only.

The “Title Page” means, for a printed book, the title page itself, plus such following pages as are needed to hold, legibly, the material this License requires to appear in the title page. For works in formats which do not have any title page as such, “Title Page” means the text near the most prominent appearance of the work’s title, preceding the beginning of the body of the text.

The “publisher” means any person or entity that distributes copies of the Document to the public.

A section “Entitled XYZ” means a named subunit of the Document whose title either is precisely XYZ or contains XYZ in parentheses following text that translates XYZ in another language. (Here XYZ stands for a specific section name mentioned below, such as “Acknowledgements”, “Dedications”, “Endorsements”, or “History”.) To “Preserve the Title” of such a section when you modify the Document means that it remains a section “Entitled XYZ” according to this definition.

The Document may include Warranty Disclaimers next to the notice which states that this License applies to the Document. These Warranty Disclaimers are considered to be included by reference in this License, but only as regards disclaiming warranties: any other implication that these Warranty Disclaimers may have is void and has no effect on the meaning of this License.

## 2. VERBATIM COPYING

You may copy and distribute the Document in any medium, either commercially or noncommercially, provided that this License, the copyright notices, and the license notice saying this License applies to the Document are reproduced in all copies, and that you add no other conditions whatsoever to those of this License. You may not use technical measures to obstruct or control the reading or further copying of the copies you make or distribute. However, you may accept compensation in exchange for copies. If you distribute a large enough number of copies you must also follow the conditions in section 3.

You may also lend copies, under the same conditions stated above, and you may publicly display copies.

## 3. COPYING IN QUANTITY

If you publish printed copies (or copies in media that commonly have printed covers) of the Document, numbering more than 100, and the Document’s license notice requires Cover Texts, you must enclose the copies in covers that carry, clearly and legibly, all

these Cover Texts: Front-Cover Texts on the front cover, and Back-Cover Texts on the back cover. Both covers must also clearly and legibly identify you as the publisher of these copies. The front cover must present the full title with all words of the title equally prominent and visible. You may add other material on the covers in addition. Copying with changes limited to the covers, as long as they preserve the title of the Document and satisfy these conditions, can be treated as verbatim copying in other respects.

If the required texts for either cover are too voluminous to fit legibly, you should put the first ones listed (as many as fit reasonably) on the actual cover, and continue the rest onto adjacent pages.

If you publish or distribute Opaque copies of the Document numbering more than 100, you must either include a machine-readable Transparent copy along with each Opaque copy, or state in or with each Opaque copy a computer-network location from which the general network-using public has access to download using public-standard network protocols a complete Transparent copy of the Document, free of added material. If you use the latter option, you must take reasonably prudent steps, when you begin distribution of Opaque copies in quantity, to ensure that this Transparent copy will remain thus accessible at the stated location until at least one year after the last time you distribute an Opaque copy (directly or through your agents or retailers) of that edition to the public.

It is requested, but not required, that you contact the authors of the Document well before redistributing any large number of copies, to give them a chance to provide you with an updated version of the Document.

#### 4. MODIFICATIONS

You may copy and distribute a Modified Version of the Document under the conditions of sections 2 and 3 above, provided that you release the Modified Version under precisely this License, with the Modified Version filling the role of the Document, thus licensing distribution and modification of the Modified Version to whoever possesses a copy of it. In addition, you must do these things in the Modified Version:

- A. Use in the Title Page (and on the covers, if any) a title distinct from that of the Document, and from those of previous versions (which should, if there were any, be listed in the History section of the Document). You may use the same title as a previous version if the original publisher of that version gives permission.
- B. List on the Title Page, as authors, one or more persons or entities responsible for authorship of the modifications in the Modified Version, together with at least five of the principal authors of the Document (all of its principal authors, if it has fewer than five), unless they release you from this requirement.
- C. State on the Title page the name of the publisher of the Modified Version, as the publisher.
- D. Preserve all the copyright notices of the Document.
- E. Add an appropriate copyright notice for your modifications adjacent to the other copyright notices.
- F. Include, immediately after the copyright notices, a license notice giving the public permission to use the Modified Version under the terms of this License, in the form shown in the Addendum below.

- G. Preserve in that license notice the full lists of Invariant Sections and required Cover Texts given in the Document’s license notice.
- H. Include an unaltered copy of this License.
- I. Preserve the section Entitled “History”, Preserve its Title, and add to it an item stating at least the title, year, new authors, and publisher of the Modified Version as given on the Title Page. If there is no section Entitled “History” in the Document, create one stating the title, year, authors, and publisher of the Document as given on its Title Page, then add an item describing the Modified Version as stated in the previous sentence.
- J. Preserve the network location, if any, given in the Document for public access to a Transparent copy of the Document, and likewise the network locations given in the Document for previous versions it was based on. These may be placed in the “History” section. You may omit a network location for a work that was published at least four years before the Document itself, or if the original publisher of the version it refers to gives permission.
- K. For any section Entitled “Acknowledgements” or “Dedications”, Preserve the Title of the section, and preserve in the section all the substance and tone of each of the contributor acknowledgements and/or dedications given therein.
- L. Preserve all the Invariant Sections of the Document, unaltered in their text and in their titles. Section numbers or the equivalent are not considered part of the section titles.
- M. Delete any section Entitled “Endorsements”. Such a section may not be included in the Modified Version.
- N. Do not retitle any existing section to be Entitled “Endorsements” or to conflict in title with any Invariant Section.
- O. Preserve any Warranty Disclaimers.

If the Modified Version includes new front-matter sections or appendices that qualify as Secondary Sections and contain no material copied from the Document, you may at your option designate some or all of these sections as invariant. To do this, add their titles to the list of Invariant Sections in the Modified Version’s license notice. These titles must be distinct from any other section titles.

You may add a section Entitled “Endorsements”, provided it contains nothing but endorsements of your Modified Version by various parties—for example, statements of peer review or that the text has been approved by an organization as the authoritative definition of a standard.

You may add a passage of up to five words as a Front-Cover Text, and a passage of up to 25 words as a Back-Cover Text, to the end of the list of Cover Texts in the Modified Version. Only one passage of Front-Cover Text and one of Back-Cover Text may be added by (or through arrangements made by) any one entity. If the Document already includes a cover text for the same cover, previously added by you or by arrangement made by the same entity you are acting on behalf of, you may not add another; but you may replace the old one, on explicit permission from the previous publisher that added the old one.

The author(s) and publisher(s) of the Document do not by this License give permission to use their names for publicity for or to assert or imply endorsement of any Modified Version.

## 5. COMBINING DOCUMENTS

You may combine the Document with other documents released under this License, under the terms defined in section 4 above for modified versions, provided that you include in the combination all of the Invariant Sections of all of the original documents, unmodified, and list them all as Invariant Sections of your combined work in its license notice, and that you preserve all their Warranty Disclaimers.

The combined work need only contain one copy of this License, and multiple identical Invariant Sections may be replaced with a single copy. If there are multiple Invariant Sections with the same name but different contents, make the title of each such section unique by adding at the end of it, in parentheses, the name of the original author or publisher of that section if known, or else a unique number. Make the same adjustment to the section titles in the list of Invariant Sections in the license notice of the combined work.

In the combination, you must combine any sections Entitled “History” in the various original documents, forming one section Entitled “History”; likewise combine any sections Entitled “Acknowledgements”, and any sections Entitled “Dedications”. You must delete all sections Entitled “Endorsements.”

## 6. COLLECTIONS OF DOCUMENTS

You may make a collection consisting of the Document and other documents released under this License, and replace the individual copies of this License in the various documents with a single copy that is included in the collection, provided that you follow the rules of this License for verbatim copying of each of the documents in all other respects.

You may extract a single document from such a collection, and distribute it individually under this License, provided you insert a copy of this License into the extracted document, and follow this License in all other respects regarding verbatim copying of that document.

## 7. AGGREGATION WITH INDEPENDENT WORKS

A compilation of the Document or its derivatives with other separate and independent documents or works, in or on a volume of a storage or distribution medium, is called an “aggregate” if the copyright resulting from the compilation is not used to limit the legal rights of the compilation’s users beyond what the individual works permit. When the Document is included in an aggregate, this License does not apply to the other works in the aggregate which are not themselves derivative works of the Document.

If the Cover Text requirement of section 3 is applicable to these copies of the Document, then if the Document is less than one half of the entire aggregate, the Document’s Cover Texts may be placed on covers that bracket the Document within the aggregate, or the electronic equivalent of covers if the Document is in electronic form. Otherwise they must appear on printed covers that bracket the whole aggregate.

## 8. TRANSLATION

Translation is considered a kind of modification, so you may distribute translations of the Document under the terms of section 4. Replacing Invariant Sections with



translations requires special permission from their copyright holders, but you may include translations of some or all Invariant Sections in addition to the original versions of these Invariant Sections. You may include a translation of this License, and all the license notices in the Document, and any Warranty Disclaimers, provided that you also include the original English version of this License and the original versions of those notices and disclaimers. In case of a disagreement between the translation and the original version of this License or a notice or disclaimer, the original version will prevail.

If a section in the Document is Entitled “Acknowledgements”, “Dedications”, or “History”, the requirement (section 4) to Preserve its Title (section 1) will typically require changing the actual title.

## 9. TERMINATION

You may not copy, modify, sublicense, or distribute the Document except as expressly provided under this License. Any attempt otherwise to copy, modify, sublicense, or distribute it is void, and will automatically terminate your rights under this License.

However, if you cease all violation of this License, then your license from a particular copyright holder is reinstated (a) provisionally, unless and until the copyright holder explicitly and finally terminates your license, and (b) permanently, if the copyright holder fails to notify you of the violation by some reasonable means prior to 60 days after the cessation.

Moreover, your license from a particular copyright holder is reinstated permanently if the copyright holder notifies you of the violation by some reasonable means, this is the first time you have received notice of violation of this License (for any work) from that copyright holder, and you cure the violation prior to 30 days after your receipt of the notice.

Termination of your rights under this section does not terminate the licenses of parties who have received copies or rights from you under this License. If your rights have been terminated and not permanently reinstated, receipt of a copy of some or all of the same material does not give you any rights to use it.

## 10. FUTURE REVISIONS OF THIS LICENSE

The Free Software Foundation may publish new, revised versions of the GNU Free Documentation License from time to time. Such new versions will be similar in spirit to the present version, but may differ in detail to address new problems or concerns. See <http://www.gnu.org/copyleft/>.

Each version of the License is given a distinguishing version number. If the Document specifies that a particular numbered version of this License “or any later version” applies to it, you have the option of following the terms and conditions either of that specified version or of any later version that has been published (not as a draft) by the Free Software Foundation. If the Document does not specify a version number of this License, you may choose any version ever published (not as a draft) by the Free Software Foundation. If the Document specifies that a proxy can decide which future versions of this License can be used, that proxy’s public statement of acceptance of a version permanently authorizes you to choose that version for the Document.

## 11. RELICENSING



“Massive Multiauthor Collaboration Site” (or “MMC Site”) means any World Wide Web server that publishes copyrightable works and also provides prominent facilities for anybody to edit those works. A public wiki that anybody can edit is an example of such a server. A “Massive Multiauthor Collaboration” (or “MMC”) contained in the site means any set of copyrightable works thus published on the MMC site.

“CC-BY-SA” means the Creative Commons Attribution-Share Alike 3.0 license published by Creative Commons Corporation, a not-for-profit corporation with a principal place of business in San Francisco, California, as well as future copyleft versions of that license published by that same organization.

“Incorporate” means to publish or republish a Document, in whole or in part, as part of another Document.

An MMC is “eligible for relicensing” if it is licensed under this License, and if all works that were first published under this License somewhere other than this MMC, and subsequently incorporated in whole or in part into the MMC, (1) had no cover texts or invariant sections, and (2) were thus incorporated prior to November 1, 2008.

The operator of an MMC Site may republish an MMC contained in the site under CC-BY-SA on the same site at any time before August 1, 2009, provided the MMC is eligible for relicensing.

## ADDENDUM: How to use this License for your documents

To use this License in a document you have written, include a copy of the License in the document and put the following copyright and license notices just after the title page:

```
Copyright (C) year your name.
Permission is granted to copy, distribute and/or modify this document
under the terms of the GNU Free Documentation License, Version 1.3
or any later version published by the Free Software Foundation;
with no Invariant Sections, no Front-Cover Texts, and no Back-Cover
Texts. A copy of the license is included in the section entitled ‘‘GNU
Free Documentation License’’.
```

If you have Invariant Sections, Front-Cover Texts and Back-Cover Texts, replace the “with...Texts.” line with this:

```
with the Invariant Sections being list their titles, with
the Front-Cover Texts being list, and with the Back-Cover Texts
being list.
```

If you have Invariant Sections without Cover Texts, or some other combination of the three, merge those two alternatives to suit the situation.

If your document contains nontrivial examples of program code, we recommend releasing these examples in parallel under your choice of free software license, such as the GNU General Public License, to permit their use in free software.

## BFD Index

(Index is nonexistent)

The body of this manual is set in  
cmr10 at 10.95pt,  
with headings in **cmb10 at 10.95pt**  
and examples in cmr10 at 10.95pt.  
*cmr10 at 10.95pt* and  
*cml10 at 10.95pt*  
are used for emphasis.